

142(8)
V
8, 1, 1, 1

TDMA

→ different frequencies at ^{same} time

→ different frequencies separated by time division.

— This pattern is called Fixed TDMA

→ fixed Tdm.

— Time slots were divided.

— The sender will send data in a particular time slot,
at the time, receiver will receive at the same time
slot.

— Again, MAC will help in dividing the time slot

total pattern will repeat after 10ms.

each time slot → 417 us.

uplink → 2nd 12 slots.

downlink 1st 12 slots.

disadvantages

— If no data is sent, though time slot is provided then

there is a waste of Bandwidth, & Frequency.

— while accessing some data, if it exceeds the division time

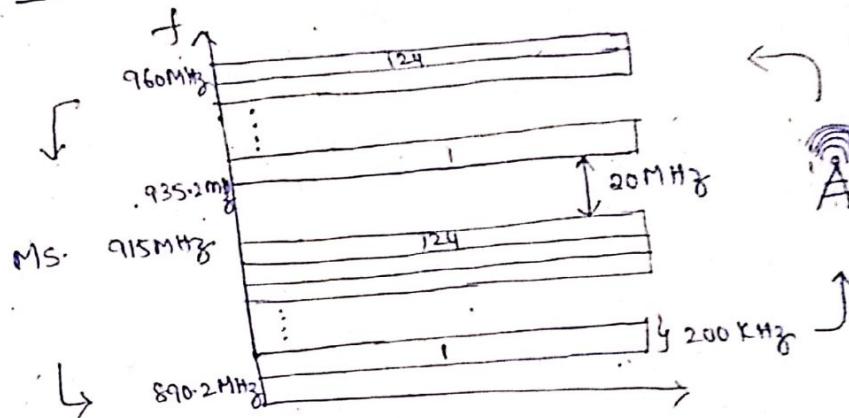
it has to wait till the next division comes to the station

classical Aloha

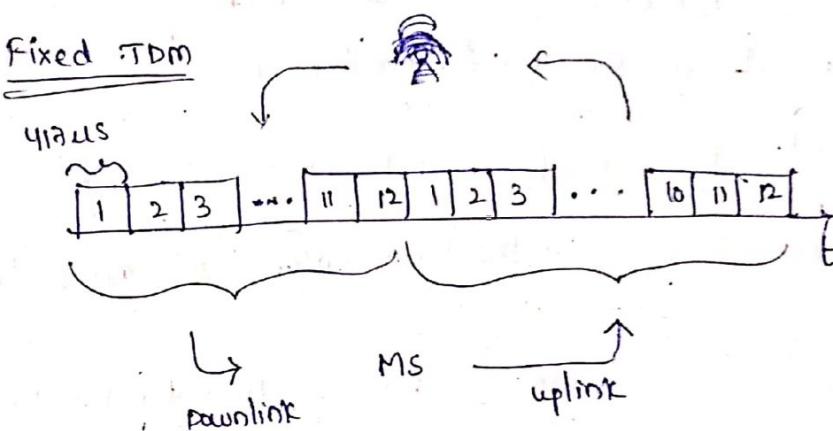
— ~~Aloha~~

— slotted Aloha.

FDMA

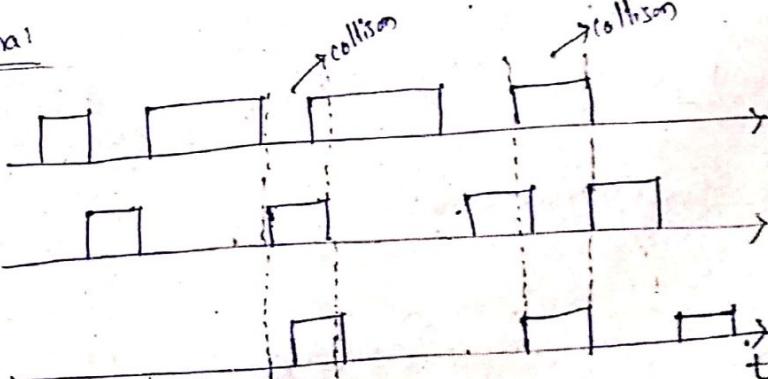


Fixed TDM



Classical Aloha

Sender A

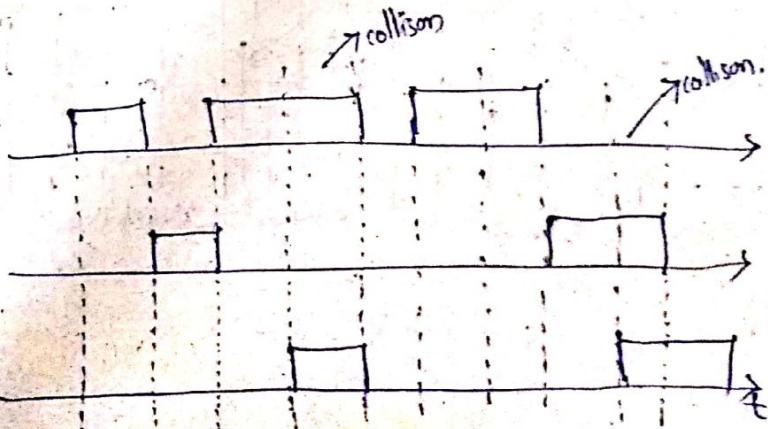


Slotted Aloha

Sender A:

Sender B:

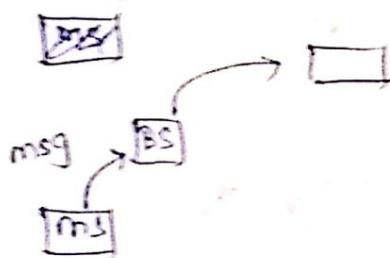
Sender C:



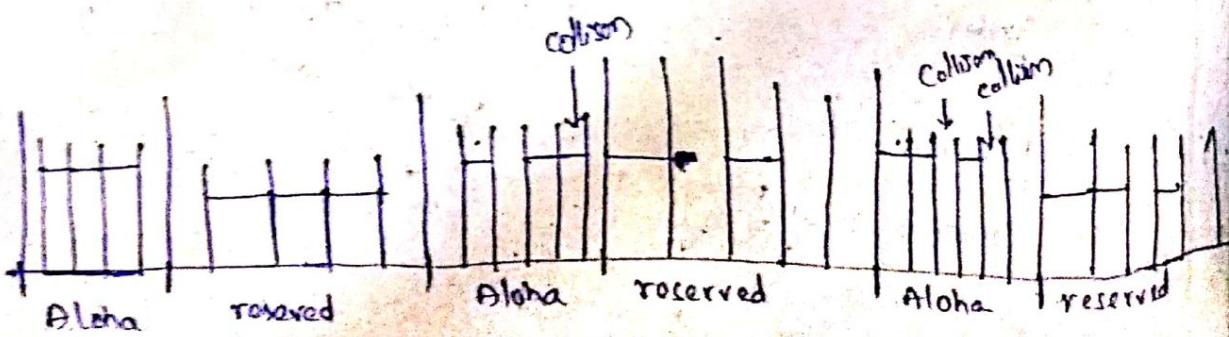
zglossis

DAMA : Demand Assigned multiple access

→ The station who ever want to send data, it will demand for collision free transmission

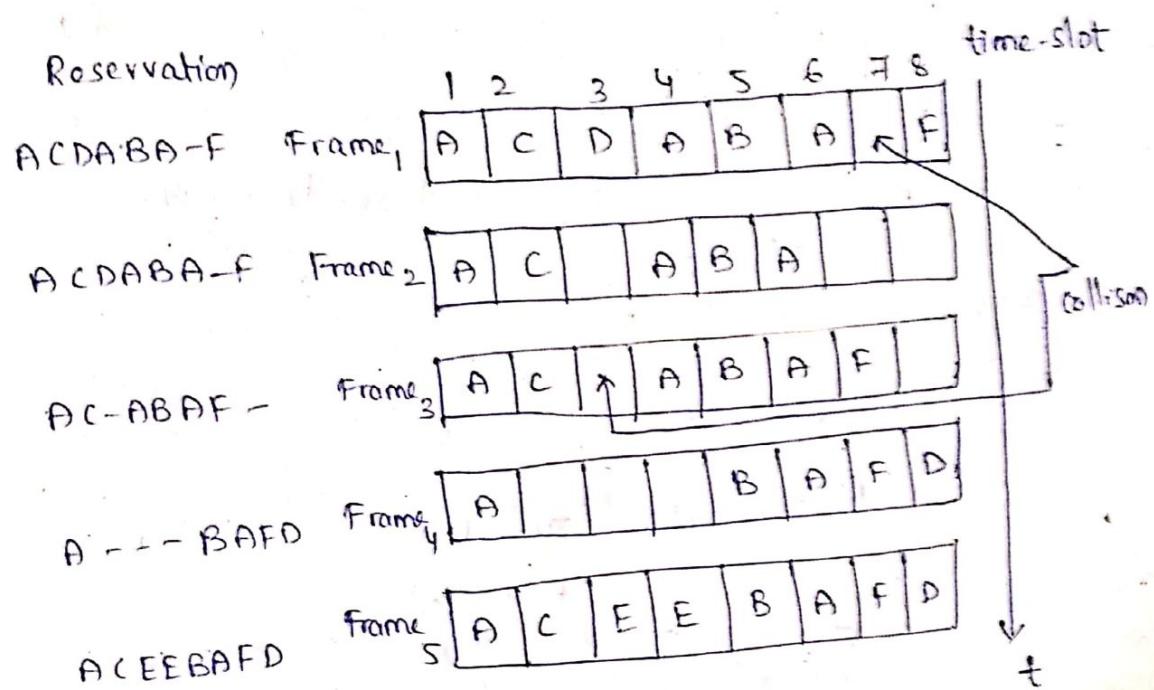


- It reserve time.
- at time of reservation collisions may take place
- but during transmission no collision taken place
- After reservation, satellite control will send back a list to all the base stations, all the base stations will accept the list
- If collision occurs during any time slot; then it will not be allocated to any station.
- This is called "Explicit Reservation" it is also called as "Reservation Aloha"



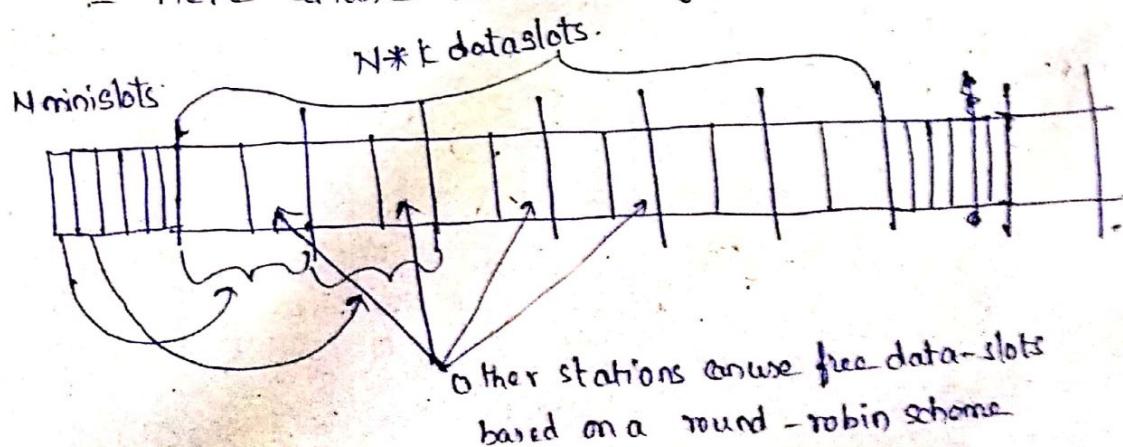
PRMA: Packet Reservation multiple Access

- In this 2 or more time slots will form a frame.
- here also reservation takes place
- According to message size, future slots are also be reserved.
- The request to satellite \rightarrow short request



Reservation TDMA:

- 6 slots will become 1 frame,
- each time slot divided into two
- here unused minislots be given to other stations.



Eg: N=6, k=2.

multiple access with collision avoidance (MAC):

→ It is a simple technique which avoid collision without use of TDMA / FDMA etc.

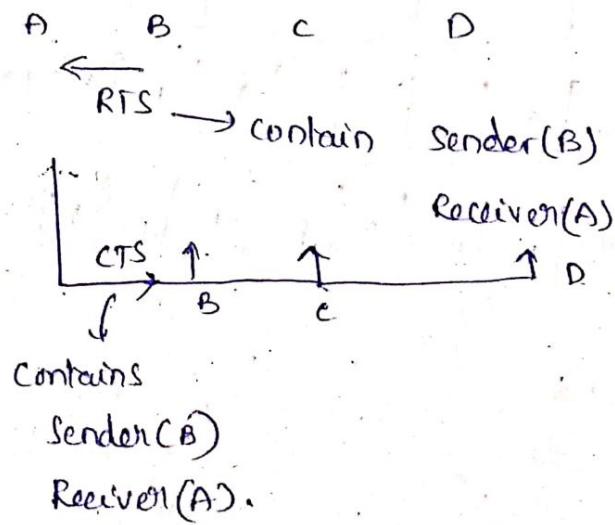
→ RTS - Request to send.

→ CTS - clear to send.

which station receives to send CTS, then only those stations are allowed to send.

→ This helps to avoid collisions in hidden terminals.

→ In Exposed terminal



29/08/19

21/03/19

Mobile IP

→ Mobile internet protocol

Goals:

prefix \downarrow 244.13.12.24

physical subnet of server

- until there is a link layer connection to mobileip.

It should provide disconnection less

data transfer.

Assumptions

* DHCP - Dynamic host configuration protocol.

- If any device moving from one place to another.

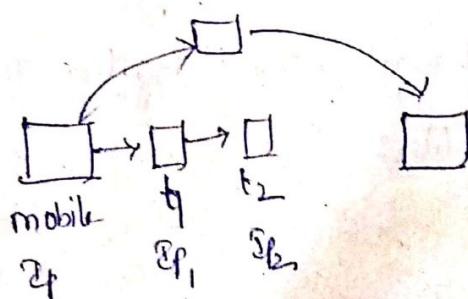
In order to get connected, allocate different mobile IP

by using DHCP.

* mobile there is a lot of mobile ip.

- problem here which will store all ip addresses and routing table

problem DNS takes some time to update the routing table



* Tcp - Transmission control protocol

Tcp will transfer data, but when ip address changes,

Tcp connection get disconnection.

Tcp - send based on socket pair .

(source, port).



→ Ip changing also gives packet losing,

when Ip is changed then Routing don't know where to route,
it sometimes drop so, data lost.

Requirements of new mobile Ip:

Assignment
①

- Compatibility
- Transparency
- Efficiency / Scalability
- Security

Compatibility: It should integrate with the previously designed services.

- They haven't changed the connectivity

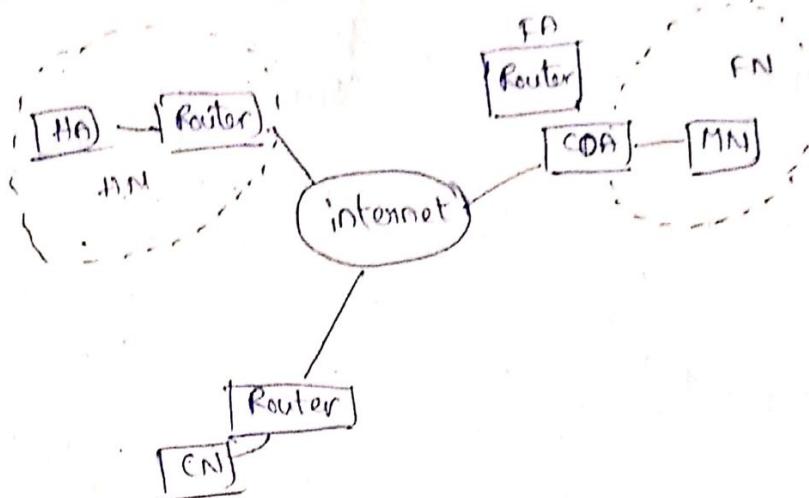
- It enhanced the efficiency of routers, switches without changing h/w or s/w.

Transparency: Internal transfer of data / Routing is not known to the end users. [no transparency] here.

Security: Security will be given by Encryption & decryption by the end routers

~~Log 9A~~ Entities & Terminology:

- mobile node (MN)



Foreign Agent

CoA - c/o address.

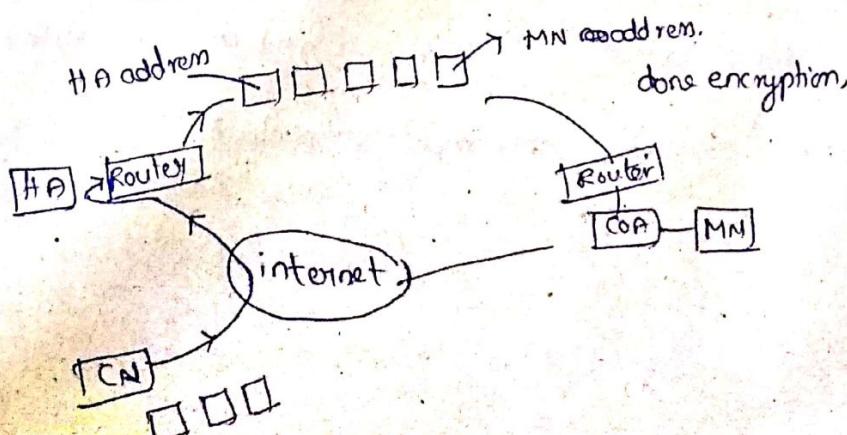
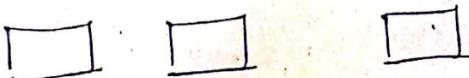
without passing through CoA, no packet will reach and get out from MN

→ 2 types

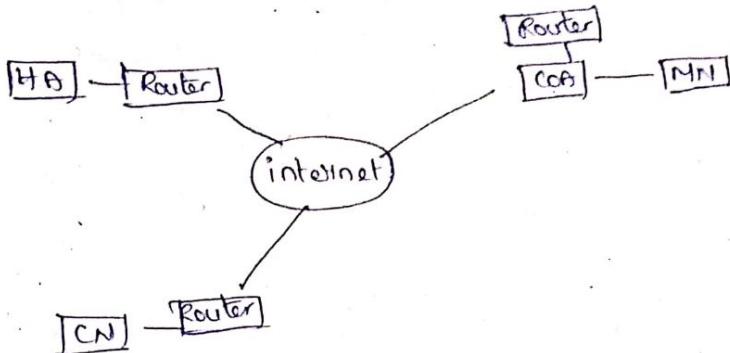
- foreign agent CoA
- co-located CoA

- correspondent-node (CN)

IP address of sender	Packet	Destination address
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- Foreign agent COA
- co-located agent COA



type - 9

code - m

checksum

address

lifetime

Route

Pre

mobile

type

length

sequ

Foreign agent COA → CN to HA
HA to MN

co-located agent COA → CN to internet
internet to MN using
DHCP & DNS
without getting enter into COA

Agent discovery!

MN goes from H

- Agent solicitation
- Agent Advertisement

All HA & FA advertise themselves and make advertise

ICMP - internet control message protocol

Type - 9 default

code - may be 0 or 16 if a mobile to other
16 - mobile to mobile.

checksum - 16th bit is complement
address - all the original address of the mobile nodes

Lifetime - advertising life time
it can much larger if has to be displayed

Router address

Preference level.

mobile extension.

Type = 16 default

length - can be anything

Sequence number : How many advertisements have done,

How much time the mobile node is
present in that current location

R-set \Rightarrow though having temporary IP, still have to
register with CN

B - If HA(D) FN, have mobile no of MN. It is saying
I am busy, so don't come into my area.

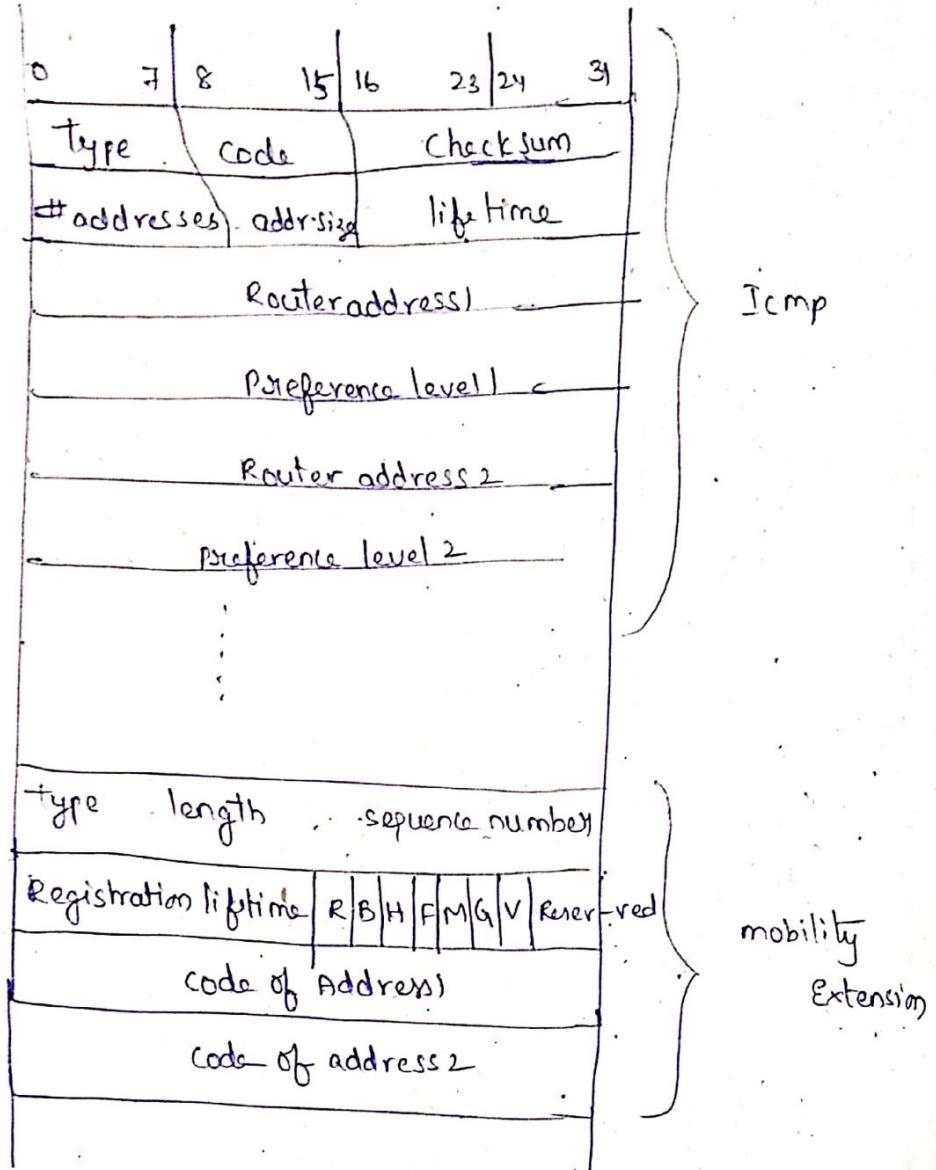
H - mobile node is in its location

F - mobile node is in the foreign location

m - It is performing routing

s - generic routing algorithm

v = 0

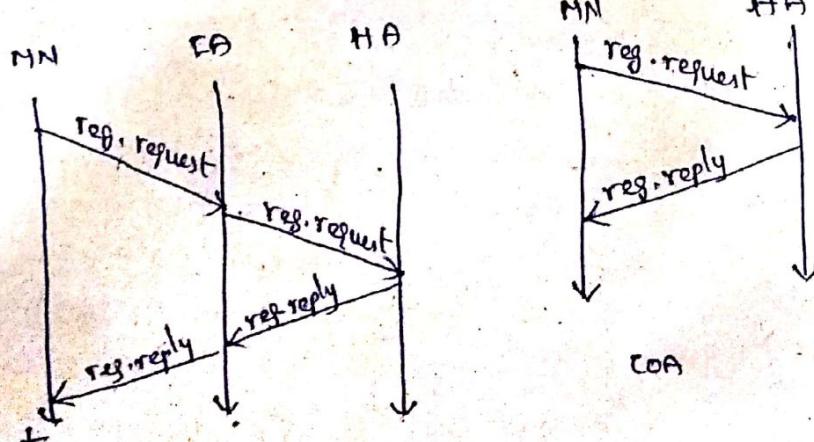


Type - set to 9

code - 0 if agent routes traffic from non-mobile nodes as well & else 16.

~~initial~~

Registration



mobile Ip registration request

0	7 8	15 16	23 24	31
Type	S B D M G Visv			lifetime
				home address
				home agent
				CoA
				Identification
				Extensions ..

UDP Packet on port 343

type=1 for regis. request

S : retain prior mobility bindings

B : forward broadcast packets

D : co-located soft address \Rightarrow MN decapsulates packets.

mobile Ip registration reply

0	7 8	15 16	31
Type = 3	code	life time	
			home address
			home agent
			Identification
			Extensions

Registration successful

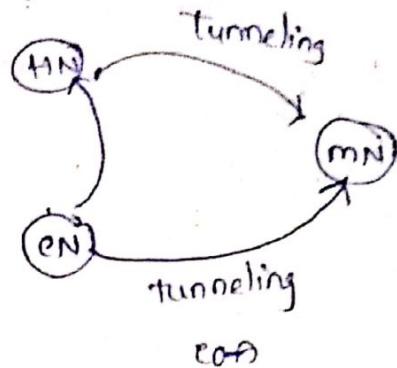
0 - registration accepted

1 -

Tunneling & Encapsulation

- where encapsulation starts called tunnel

through tunnel a packet moves from



Types of encapsulation

- Ip-in-ip encapsulation
- minimal encapsulation
- Generic Routing Encapsulation (GRE)

- cache invalidation

- Time based invalidation.

- command based invalidation

Assignment → - Dynacache API and CACTIEV L table invalidation

- Group based invalidation.

→

<timeout> value <timeout> value may be anything
in seconds

<inactivity> value <inactivity>

→

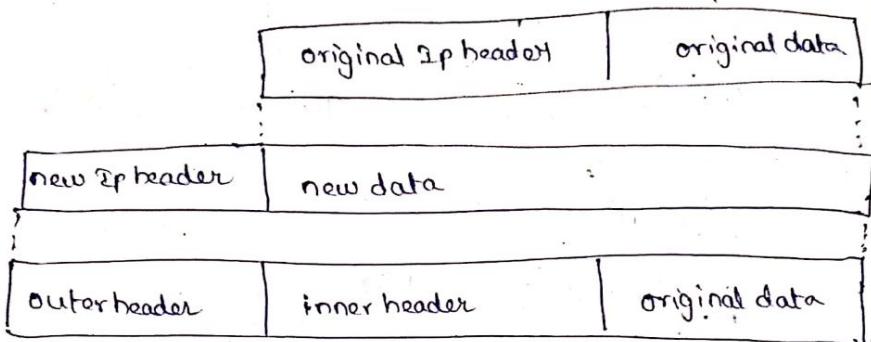
through commands

→

Power aware computing

- based on the switches & routers power consumption is measured.
- goal is to conserve the power
- two types of
 - unicast 1 to 1
 - multicast 1 to many
- 5 Routing metrics
 - Energy consumed by one packet
 - time taken/w partition
 - variance in power level across mobiles
 - cost for one packet
 - maximum mobile cost
- Power aware computing methods
 - data caching : data is cached at mobile device
 - cache invalidation mechanism : remove old data
 - Normalization of records :
Before sending data, it should be normalized as 8x duplicate records should be suppressed and not trans

Encapsulation



Encapsulation methods

IP-in-IP encapsulation packet format

ver.	IHL	DS (TOS)	length					
IP identification		flags fragment offset						
TTL	IP-in-IP		IP checksum					
IP address of HA								
core-of address of COA								
ver	IHL	DS (TOS)	length					
IP identification		flags	fragment offset					
TTL	IP-in-IP		IP checksum					
IP address of CN								
IP address of MN								
TCP/UDP --- Pay load								

IP-in-IP encapsulation

- ver - IP f
- IHL - inter
- TOS - typ
- Length -
- IP id, +
- TTL -
- IP-in

Minimal IP

ver			
IP id			
TTL			
layer 4			

IP-in-IP encapsulation

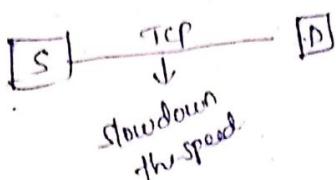
- ver - IP protocol version no.
- IHL - internal header length
- TOS - type of services (copied from inner header)
- Length - complete encapsulated packet length
- IP id, flags, flag offset - used for fragments
- TTL - time to live
- IP-in-IP upper layer protocol

Minimal Routing Encapsulation

ver	IHL	DS (TOS)	length
		IP identification	flags fragment offset
TTL		min-encapsulation	IP checksum
		IP address of HA	
		care-of address of COA	
layer 4 protocol	S	reserved	IP checksum
			IP address of MN
			original sender IP address (if S=1)
			ICP/UDP/--- payload

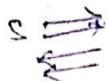
~~Traditional~~ Traditional TCP / GECN

- congestion - due to congestion, the packet may drop
- speed of up link > speed of dp link



- If congestion is seen - it will start slow start mechanism
- Upp - is not reliable, may drop packet in b/w the transmission.
- congestion window increases by doubling the previous value
 - the growth is exponentially.

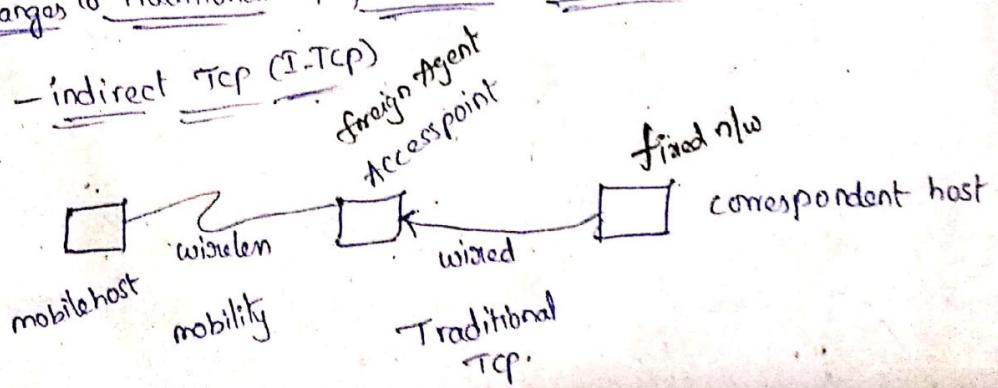
$$\text{congestion window} = x 2^A$$



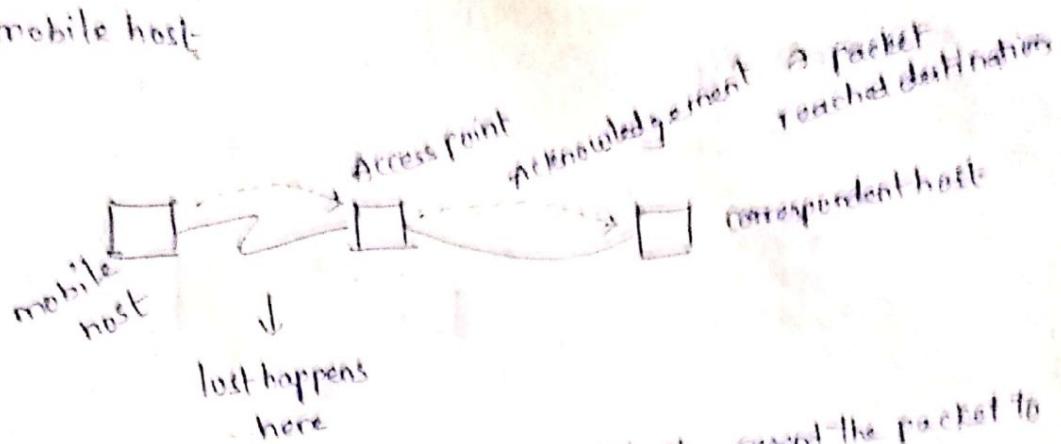
when the ~~the~~ value reaches the ~~the~~ congestion threshold, then it will show linear relation / linear growth.

Changes to Traditional TCP / Improvements to TCP

Indirect TCP (I-TCP)



The mobile host



- then Access point takes responsibility to resend the packet to the mobile host from buffer.
- from mobile host to correspondent host, if packet lost then mobile host takes responsibility to resend

Advantages

- Indirect TCP doesn't require any changes in TCP
- Segmentation & transmission errors will not disturb each other.
- Testing can be performed in original standard TCP
- short round trip time and easy receiving.

Disadvantage

- Loss of end-to-end semantics of TCP
- Handovers may create more problems if it takes a while to register to the new location
- Foreign agent must be a trusted entity because the TCP connection ends at this point, and it has to be integrated into all security mechanisms.

- Snoping TCP



- Using one

MH requires

- Using IP

numbers

schemes

— Mobile T

Advantages:

- End-to-End semantics is maintained here
- No improvements & enhancements are needed to the correspondent host
- Packet is buffered to new location of another foreign agent, if handover is performed at mobile host
- It doesn't matter if the next foreign agent uses the same enhancement or not

Disadvantages:

- The problems of the wireless link are now visible following correspondent CH as it is not fully isolated from the original standard TCP.

Adva

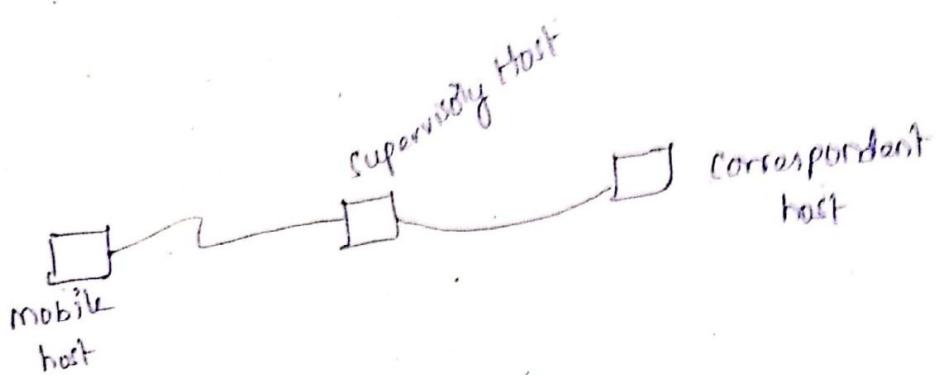
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- Using 3rd duplicate acknowledgement between FA8L
mH requires additional mechanisms on the FA8L
- Using IP encapsulation security snooping on the sequence
numbers will no longer work because many security
schemes prevent replay attack

Mobile TCP (M-TCP)



SH-request forces the correspondent node (cnode) to set the
window size = 0 with help of TCP. It will go to
persistant mode. (when mH doesn't get connected to
internet)

Advantages:

- It maintains the TCP end-to-end semantics
- The SH doesn't send any acknowledge by itself, but forwards the acknowledgement from the mH
- If mH gets disconnected it avoids useless retransmissions & slow start, breaking the connection.

Disadvantages

- SH doesn't act as proxy as in Indirect-Tcp,Packet
- loss on the wireless link is propagated to the sender
- A modified Tcp on the wireless link not only requires modifications but also a new flow element like Bandwidth manager