

# Random number generator

HERNÁNDEZ MARTINEZ JORGE IVÁN

Departamento de Física, CUCEI, Universidad de Guadalajara.

jorgeivanhdzmtz@gmail.com

## I. INTRODUCTION

The most common method is to generate the following number from the last generated numbers:

$$x_n = f(x_{n-1}, x_{n-2}, \dots) \quad (1)$$

For example we can use the recurrence relation

$$x_{n+1} = ax_n + c \pmod{m} \quad (2)$$

where  $m$  is called modulus and  $a$  and  $c$  are positive integers called the multiplier and the increment, respectively. But in these recurrence relation will eventually repeat itself with a period no greater than  $m$ . For fix these, we need to choose  $m$  that is the maximal length. The number  $m$  is usually close to the machine's largest representable integer  $\approx 2^{32}$  [1]

## II. PARK AND MILLER'S METHOD

they propose a "Minimal Standard" generator based on the choices

$$a = 7^5 = 16807 \quad m = 2^{31} - 1 = 2147483647 \quad (3)$$

but it represents that a multiply of  $a$  and  $m$ , exceeds the maximum value for a 32 bits integer. This is the reason why the next Schrage's method, based in an approximate factorization of  $m$ , is used.

$$m = aq + r \quad \text{i.e., } q = [m/a], r = m \bmod a \quad (4)$$

with square brackets denoting integer part. If  $r$  is small, specifically  $r < q$ , and  $0 < z < m - 1$ , it can be shown that both  $a(z \bmod q)$  and  $r[z/q]$  lie in the range  $0, \dots, m-1$ , and that [2]

$$az \bmod m = \begin{cases} a(z \bmod q) - r[z/q] & \text{if it is } \geq 0, \\ a(z \bmod q) - r[z/q] + m & \text{otherwise} \end{cases}$$

**Figura 1**

this algorithm use the values  $q=127773$  and  $r=2836$ .

## I. Combined method

L'Ecuyer recommends the use of the two generators  $m_1 = 2147483563$  (with  $a_1 = 40014$ ,  $q_1 = 53668$ ,  $r_1 = 12211$ ) and  $m_2 = 2147483399$  (with  $a_2 = 40692$ ,  $q_2 = 52774$ ,  $r_2 = 3791$ ). Both moduli are slightly less than  $2^{31}$ . The periods  $m_1-1=2 \times 3 \times 7 \times 631 \times 81031$  and  $m_2-1=2 \times 19 \times 31 \times 1019 \times 1789$  share only the factor 2, so the period of the combined generator is  $\approx 2.3 \times 10^{18}$ . For present computers, period exhaustion is a practical impossibility.[1]

## III. DESCRIPTION OF CODE RAN0

I show at first the most simple code programing using Fortran, for a random number generator evenly distributed between 0 and 1.

Start creating a function called `ran0(val)` and defining our variables as integer (Ec.3,  $q$  and  $r$ ) except to `ran0` (our output) and for  $AM$  that correspond to divide  $1/IM$ , and we put it's corresponding value.

```
FUNCTION ran0(val)
INTEGER val,IA,IM,IQ,IR,MASK,k
REAL ran0,AM
PARAMETER (IA=16807,IM=2147483647,AM=1./IM)
PARAMETER ( MASK=123459876,IQ=127773,IR=2836)
```

Then for find a random value we will use equations of Ec.4. But how we see, the first value we have to allowed is zero. for this reason i used the FORTRAN command `ieor`, than performs a subtraction of bits" of each number generating a new one. This command guarantees the exclusion of zero and made or program a little bit more random.

```

val=ieor(val,MASK)
k=val/IQ
val=IA*(val-k*IQ)-IR*k
if (val.lt.0) val=val+IM
ran0=AM*val

END

```

in third line of the code part that we just saw, the operation there is the same as  $Val = \text{mod}(IA * val, IM)$ . so if we change that part, every would remains the same. the 4 and 5 line correspond that we see in figura 1.

We only have to define the principal program that calls to our routine ran0. This program has to open a unit where the numbers will be save. Is in a cycle because we need to save every number on each repetition. But it doesn't matter because each seed bring a different random number

```

program aleatorio
integer::a,b
open(unit=100, file='aleatorios.dat')
do b=1,10
  b=-1-b

  write(100,*)ran0(b)
enddo

end program

```

#### IV. DESCRIPTION OF CODE RAN2

The following routine, ran2 uses the Minimal Standard for its random value, but it shuffles the output to remove low-order serial correlations. A random deviate derived from the  $j$ th value in the sequence,  $I_j$ , is output not on the  $j$ th call, but rather on a randomized later call,  $j + 32$  on average.

First we define our type of variables and their values. I use the values of L'Ecuyer recommends the use for the two generators (subsection 1.1), and it is assigned to idum2, iv and iy a value, but this will change on each cycle.

```

FUNCTION ran2(idum)
      INTEGER idum,IM1,IM2,IMM1,IA1,IA2,IQ1,IQ2,IR1, &
        IR2,NTAB,NDIV,idum2,j,k,iv(NTAB),iy
      REAL ran2,AM,EPS,RNMX
      PARAMETER (IM1=2147483563,IM2=2147483399,AM=1./IM1, &
        IMM1=IM1-1,IA1=40014,IA2=40692,IQ1=53668,IQ2=52774, &
        IR1=12211,IR2=3791,NTAB=32,NDIV=1+IMM1/NTAB, &
        EPS=1.2e-7,RNMX=1.-EPS)
      idum2=123456789 ; iv=NTAB*0 ; iy=0

```

then we have to be sure to prevent idum=0, and give it a new value to idum2. It's important that the seed have to be a negative number to start with this part of code.

```

      if (idum.le.0) then
        idum=max(-idum,1)
        idum2=idum

```

Then we make iterations and compute schrange's method using values for  $m_1$

```

      do j=NTAB+8,1,-1
        k=idum/IQ1
        idum=IA1*(idum-k*IQ1)-k*IR1
        if (idum.lt.0) idum=idum+IM1
        if (j.le.NTAB) iv(j)=idum
      enddo
      iy=iv(1)
endif

```

then if the number is g.t 0, start here, and compute the schrange's method using values for  $m_1$  and  $m_2$

```

k=idum/IQ1
idum=IA1*(idum-k*IQ1)-k*IR1
if (idum.lt.0) idum=idum+IM1

k=idum2/IQ2
idum2=IA2*(idum2-k*IQ2)-k*IR2
if (idum2.lt.0) idum2=idum2+IM2

```

Then we have to shuffle idum and idum2 and are combined to generate output (line 2 of the next part of code)

```

j=1+iy/NDIV
iy=iv(j)-idum2
if (iy.lt.1) iy=iy+IMM1
ran2=min(AM*iy,RNMX)
END

```

to finish we just only have to define the principal program that calls to our routine ran2. This program has to open a unit where the numbers will be save. Is in a cycle because we need to save every number on each repetition. It is important remember that seed is a negative number

```

PROGRAM aleatorioran2

INTEGER::q,a

OPEN(unit=100, file='aleatoriosran2.dat')
do q=1,1000
    a=-1-q
    write(100,*)ran2(a)
enddo

END PROGRAM

```

## V. SIMULATION

Now we will check if the random numbers obtained with the last method repeat among themselves. If we graph the numbers in a plane, we have to get a random distribution. Using Next code of Matlab

```
datos= textread('C:\Users\omarc\Desktop\universidad ...  
Ivan\estadistica\programas\random\aleatoriosran2.dat');  
y = datos(:,1);  
n=length(y);  
x=1:1:n;  
c=unique(x);  
c1=length(c)  
plot(x,y, ' . ')
```

we get a 'cloud' of point what it looks like never cross.

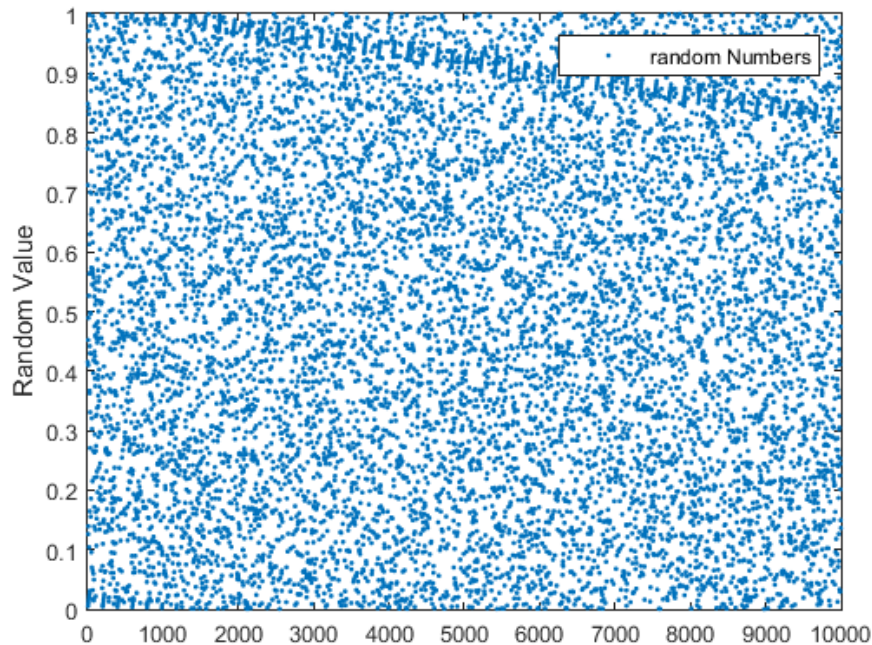


Figura 2

in the fist instance we se the value of random numbers never cross, but to make

sure of that, using the command `unique(x)` on matlab like i did in the code, it is possible to read the vector's length. If this length is less than 10000 (this number is because i generate 10000 random values) is because one or more numbers repeat. but in figura 3 we see that length of unique vector has the same length that the vector of random numbers.

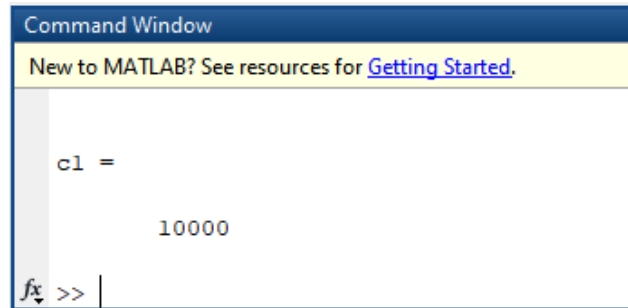


Figura 3

## REFERENCIAS

- [1] S. Teukolsky *Numerical Recipes in FORTRAN* chapter7. Get from [http : //nuclear.fis.ucm.es/wordpress/wp - content/uploads/2011/09/RandomNumbers.pdf](http://nuclear.fis.ucm.es/wordpress/wp-content/uploads/2011/09/RandomNumbers.pdf)
- [2] David G. Carta *Two Fast Implementations of the "Minimal Standard" Random Number Generator*. Volume 33 P. 87-88
- [3] Herbert Hoeger *GENERACION DE NUMEROS ALEATORIOS*. Get from [hhttp : //webdelprofesor.ula.ve/ingenieria/hhoeger/simulacion/PARTE4.pdf](http://webdelprofesor.ula.ve/ingenieria/hhoeger/simulacion/PARTE4.pdf)