

Chapter-17

i) Inserting n elements using

a) aggregate method:

The table doubles its size when it runs out of space. So if the original size is 1, after insertion it doubles the size to 2 after 2 more insertions it doubles to size 4 etc.

In general after k doublings the size is 2^k .

Pseudo code:

Initialize table with capacity = 1

for $i = 1$ to n :

if table is full:

new table = create new table with
size $2 \times$ current size

copy elements then from old table
to new table

table = new table

insert element 1 into table

Let, $k = \log(n+1) - 1$

Total cost = $O(n)k$

= $O(n \log n)$

Cost per insertion = $O(\log n)$

Run time per insertion is $O(\log n)$

Total time is $O(n) \times \log(n+1)$

(b) Accounting Method:

Charge 2 units for each ~~insert~~ insertion
when the table doubles in size from m to $2m$,
credit m units,

Total credit is $m + 2m + 4m + \dots$

$$n/2 \times m = O(n)$$

Pseudo code:

Initialize table with capacity = 1

for $i = 1$ to n :

if table is full:

newtable = create newtable
with size \times current size

copy element from old table to new
table

table = newtable

insert element i into table

initialize charges = 0

initialize credits = 0

for $i = 1$ to n :

charges \leftarrow 2

if table doubled in size

from m to $2m$

Credits $t = m$

Total Charges $= 2 \times n = O(n)$

Total Credits $= m + 2m + \dots$
 $= n/2 \times m = O(n)$

Cost per Insertion $= \text{total} / n$

$= O(n) / n$

$= O(1)$

Runtime per Insertion $= O(1)$

Total time $= O(n)$