

Relational Algebra

Relational Algebra

It's a procedural query language, which takes relation as input and generate Relation as output.

Basic fundamental operators

- ① Selection (σ) → Unary (SQL where)
 - ② projection (Π) → Unary (SQL select)
 - ③ Union (\cup) (OR যুক্তি)
 - ④ set difference (-)
 - ⑤ Cartesian product (\times)
 - ⑥ Rename (ρ) → Unary → $\rho_x(E)$ [E table renamed as x]
-] → Binary

Additional / Derived operations

- ① Join (\bowtie) → (Cartesian Product) ✗
- ② Set Intersection (\cap) → Binary ✗ → (AND যুক্তি)
- ③ Division (\div)

Unary :- মাত্রা একটি relation এ operators করে আগে ঢাকে Unary operators বলে।

Binary :- মাত্রা দুইটি relation/table এ operate করে আগে ঢাকে Binary operators বলে।

Selection operators

Database এর কোনো table থেকে row wise কোনো table select করতে ইন্টি যেকোনো selection operators ব্যবহার করা হয়।

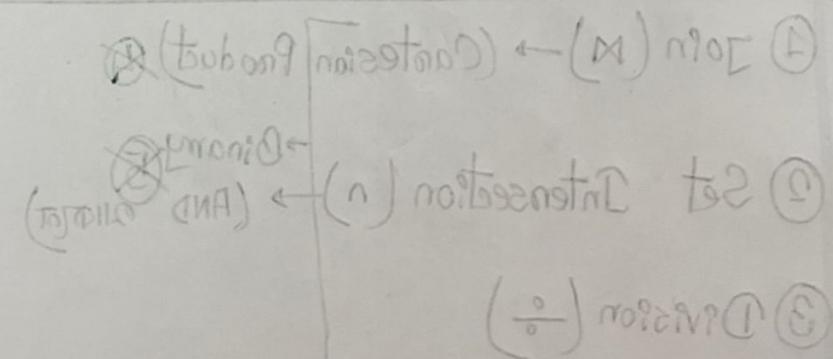
Notation

$\sigma \rightarrow \text{selection}$

$\rho \rightarrow \text{prepositional logic formula}$

$\tau \rightarrow \text{Relation}$

বিভিন্ন operators



Example:-

ID	name	dept-name	salary
22222	Einstein	physics	95000
33456	Gold	physics	87000

To select those tuples of the `instructors` relation where the `instructors` is in the "Physics" department.

$\sigma_{\text{dept-name} = \text{"Physics"}}(\text{instructors})$

We can find all `instructors` with salary greater than \$90,000.

$\sigma_{\text{salary} > 90,000}(\text{instructors})$

To find the `instructors` in Physics with a salary greater than \$90,000

$\sigma_{\text{dept-name} = \text{"physics"} \wedge \text{salary} > 90,000}(\text{instructors})$

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Projection operators:-

Database এর কোন Table টাকে column wise কোন Table select করতে ইট projection operators নামে।

* distinct data ফির কোন duplicate data ফির না।

Notations :- $\Pi_{a_1, a_2, a_3, \dots}^{(r)}$

Π → projection

a_1, \dots, a_n → Attributes

r → Relation

Example:-

① fetch ID, Name, & Board data from student table for those students whose Board = 'Dhaka'.

$\Pi_{ID, Name, Board} (\sigma_{Board='Dhaka'} (Student))$

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- ⑪ fetch Age, Name, Age & Board data from student table for student ID = 2.

$\Pi_{\text{Name}, \text{Age}, \text{Board}} (\sigma_{\text{student ID}=2} (\text{student}))$

Rename Operators

Notation $\rho (R_1, R)$

Example

- ① Query to fetch the data corresponding Name, Age & Board attributes and rename the relation as studentInfo.

$\rho (\text{StudentInfo}, \Pi_{\text{Name}, \text{Age}, \text{Board}} (\text{student}))$

- ② Query to fetch the data for ID, Name, attribute from student table for those students whose Board is Dhaka. And rename the relation as BoardStudentInfo also rename the attributes.

ID = D.ID

Name = D.Name

$\rho (\text{BoardStudentInfo}_{(\text{D.ID}, \text{D.Name})}, \Pi_{\text{ID}, \text{Name}} (\sigma_{\text{Board}=\text{"Dhaka}} (\text{student})))$

Union :-Notation :- $R_1 \cup R_2$ Example :-Rules :-

- Hence, R_1 and R_2 must have the same numbers of attributes.
- Data types of the attributes between two relations should be same.
- Duplicate tuples are automatically eliminated.

Theme:

Intersect Operation:-

Notation:-

$$R_1 \cap R_2$$

(R₁ + R₂) : combines both relations

(R₁ × R₂) : output to one relation

Rules:-

- Same
- Same
- The intersection operation always return the distinct rows.
The duplicate rows will not be returned by the intersect operation.

Set difference operation:-

Notation:-

$$R_1 - R_2$$

Grades
G
J
DGWP
GDWP

GD	Grades	ID
G	I	1
J	O	2
DGWP	I	1
G	S	3
J	I	1
G	S	3

Theme:

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Cartesian Product :-

Notation :-

$$R_1 \times R_2$$

Total no of columns: $R_1(c) + R_2(c)$

Total no of Tuples: $R_1(T) \times R_2(T)$

Division Operation :-

Notation :-

$$R_1 / R_2$$

Fetch those students ID who has taken all the courses

Course Taken

ID	Course
1	C
2	Java
1	DBMS
3	C
1	Java
3	Java

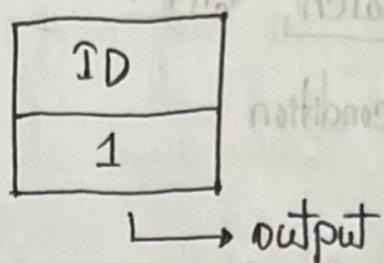
Course

Course
C
Java
DBMS

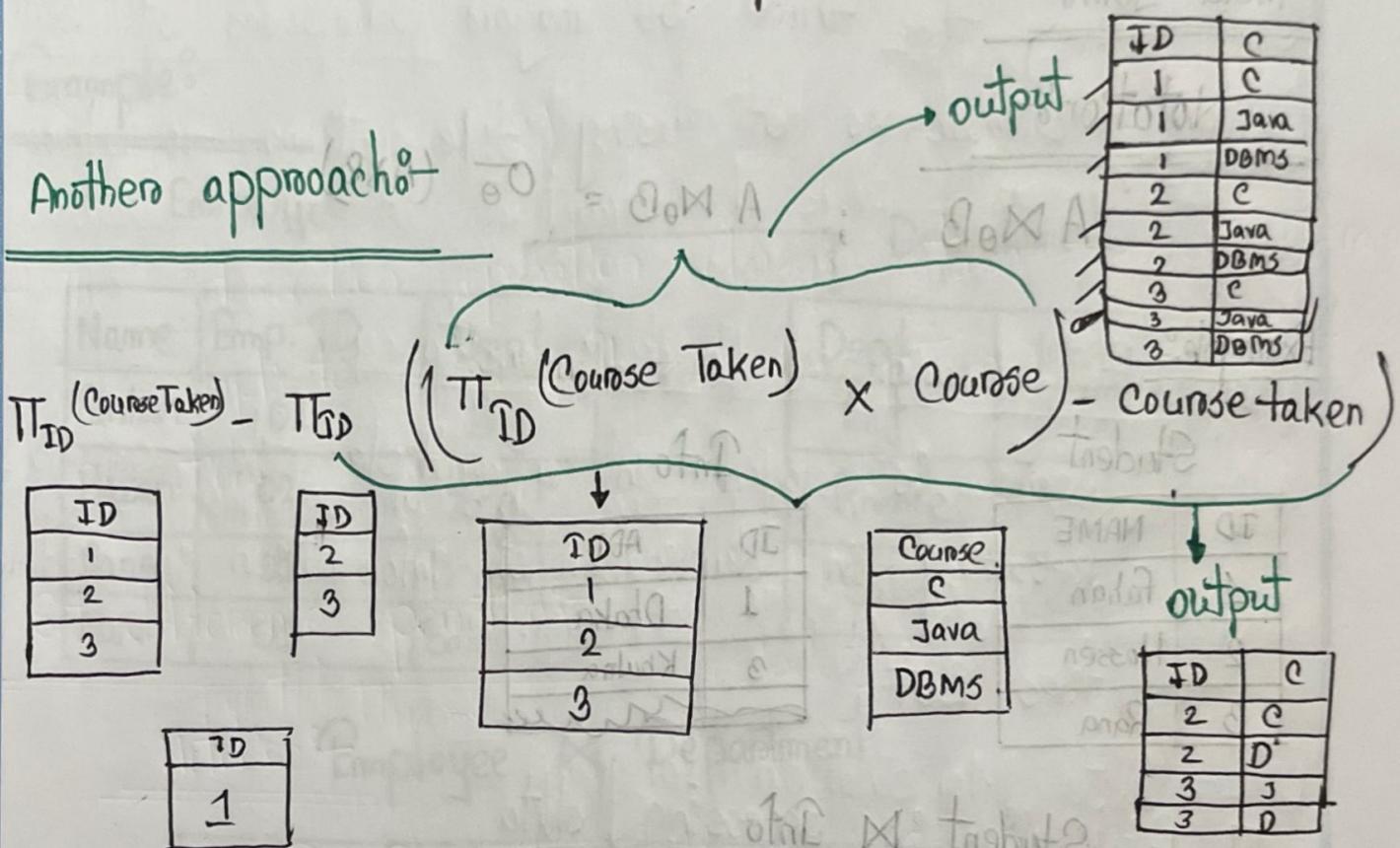
Theme:

Ans^o

Course Taken \rightarrow Course



Another approach



NAME	ID	NAME	ID
Dipika	1	Kiran	1
Rajendra	2	Raju	2

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Join

এই কাছ থেকে only ফেরত আবে tuple/rows পুরো return করবে
 যাত্রে condition match, বলবে।
matching condition

ID
L

tuple

Theta Join

Notation

 $A \bowtie_\theta B$

$$A \bowtie_\theta B = \sigma_\theta (A \times B)$$

Example

Student

ID	NAME
1	Fahim
2	Hossen
3	Rana

Info

ID	ADD
1	Dhaka
3	Khulna

Student \bowtie Info

$$\text{Student.ID} = \text{Info.ID}$$

ID	NAME	ID	NAME
1	Fahim	1	Dhaka
3	Rana	3	Khulna

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Natural join

Attribute name same or ~~পুঁজির নাম~~ (P)

Notation

$$R_1 \bowtie R_2$$

Example

Employee

Name	Emp. ID	Dept.
Farhan	01	Eng
Hossan	02	Eng
Rana	01	Finance
Abir	01	Sales

Department

Dept.	Mngro.
Eng	Rasel
Finance	Monira
Sales	Karim

Employee \bowtie Department

Name	Emp. ID	Dept.	Mngro.
Farhan	01	Eng	Rasel
Hossan	02	Eng	Rasel
Rana	01	Finance	Monira
Abir	01	Sales	Karim

Database systems

A modern database system is a complex software system whose task is to manage a large, complex collection of data.

Summarise the three major purposes of database system. / Drawbacks of using file systems to store

data

1 Data redundancy and inconsistency

data is stored in multiple file formats resulting in duplication of information in different files.

2 Difficulty in accessing data

Need to write a new program to carry out each new task.

3 Data isolation

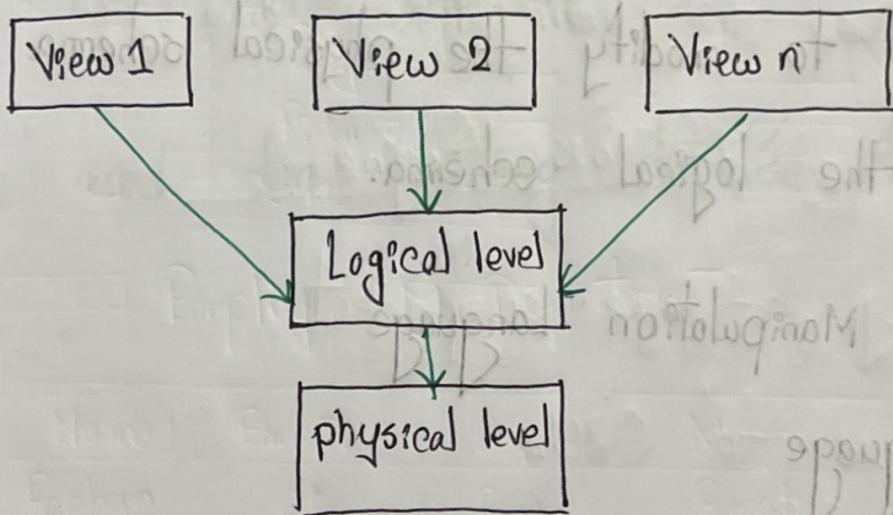
Multiple files and formats.

Security problems

Hard to provide users access to some, but not all data.

Database systems offer solutions to all the above problems.

Level of Abstraction — is a process of hiding unwanted or irrelevant details from the end user.



Theme:

Instances and Schemas

omisding phisical

two forms of schema now showing of both

Logical Schema

The overall logical structure of the database.

Physical Schema

The overall physical structure of the database.

Physical Data Independence

the ability to modify the physical schema without changing the logical schema.

DML → Data Manipulation Language

Query Language

↳ Procedural Language

↳ Declarative Language (Non-Procedural Language)

↳ SQL

DDL

Key Commands

Create - creates new table or objects

Alter - modifies the structure of an existing object

Drop - deletes objects like tables or indexes.

Truncate - removes all data from a table but keeps the table structure.

DDL used for defining database structures (like creating tables)

DML

Key Commands:-

Select → retrieves data from the database

Insert → adds new data

Update → modifies existing data

Delete → removes records

DML is used for manipulating data (like inserting data)

Database Engine

→ underlying software component

→ DBM uses to CRUD

Functions of D.E

- ① Data Storage
- ② Query Processing
- ③ Transaction management component

Storage Management

→ responsible for managing how data stored, retrieved & updated on physical storage media.

It is responsible for

↳ interaction with OS file managers

↳ efficient storing, retrieving & updating data.

The SM components -

- ① Integrity managers
- ② Transaction managers
- ③ File managers
- ④ Buffer managers

Query Processor:-

The query processor is a critical component of a DBMS that ensures SQL queries are properly parsed, optimized and executed.

DDL Interpreter:-

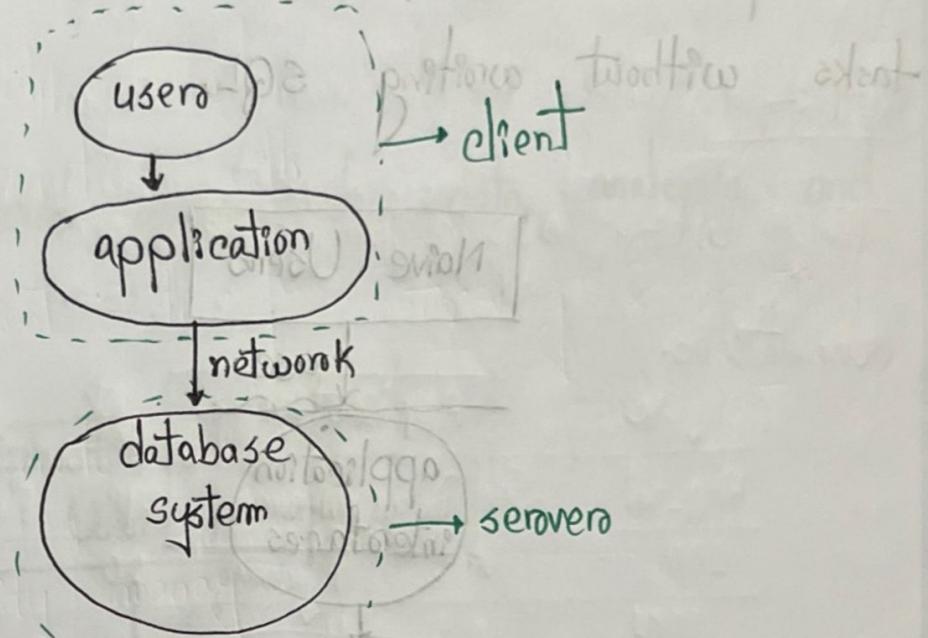
interprets DDL statements and records the definitions in the data dictionary.

DML Compiler:-

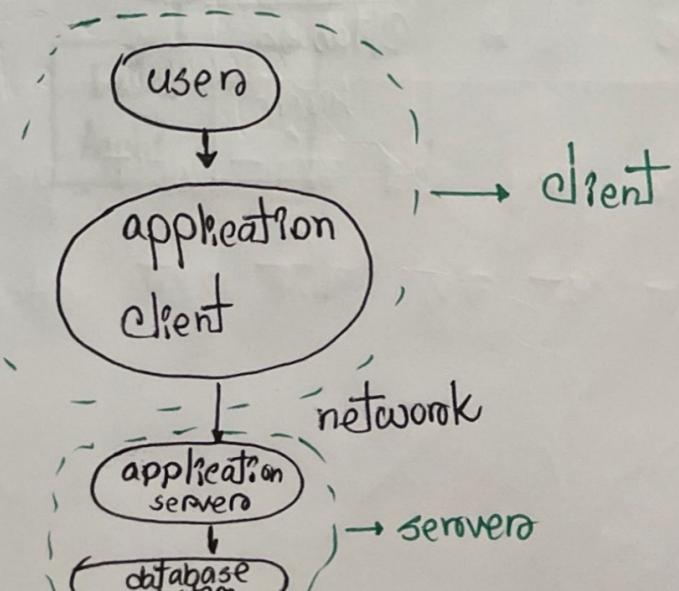
it performs query optimization.

Theme:

Two tier ~~and~~ Architecture is a common client servers architecture used in database system, where the application is divided into two tiers. client and the servers.



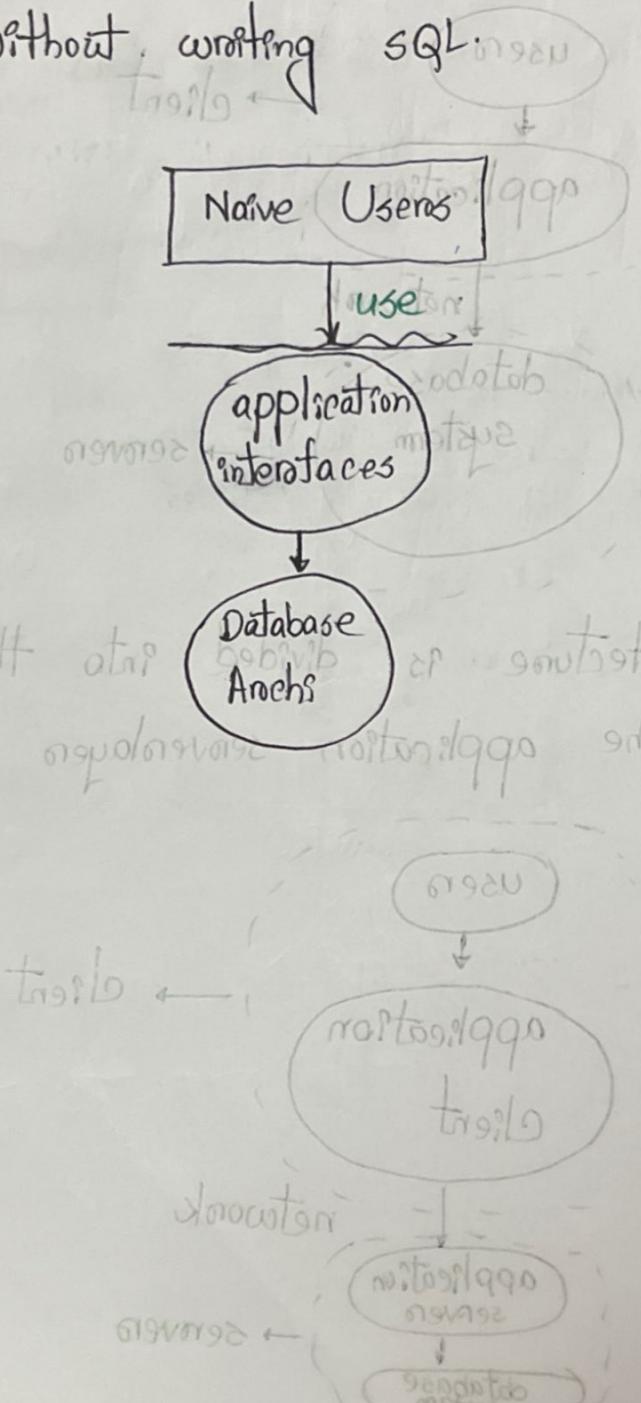
Three tier architecture is divided into three tiers, the client layer, the application servers layer and the database layer.



Naive users: (telkens, agents, web users)

Role / Responsibility

Use pre-defined queries or forms to perform repetitive tasks without writing SQL.

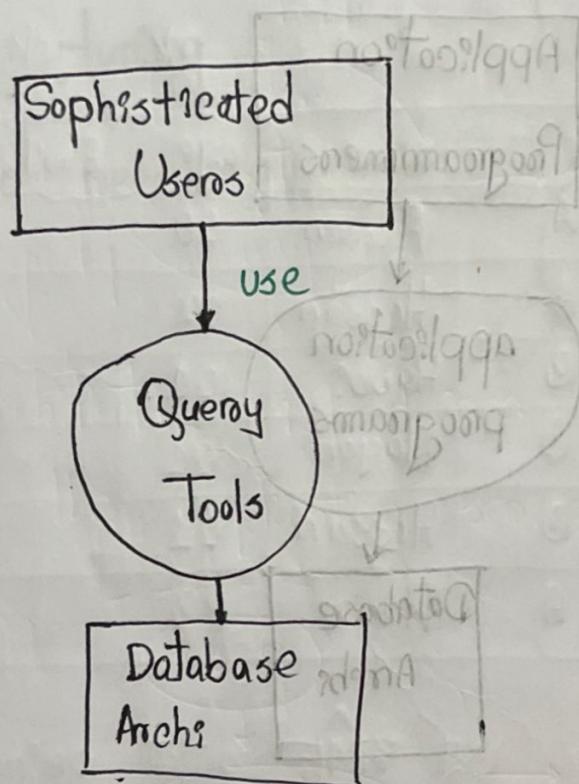


Sophisticated Users :- (Analyst)

These users interact directly with the database using complex SQL queries.

Role:-

Write advanced queries for data analysis and reporting.

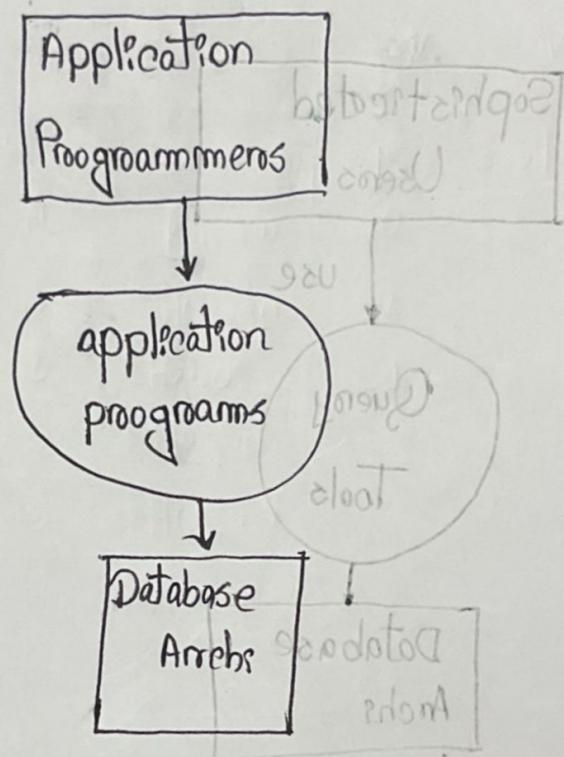


Application Programmers :-

There are developers who build applications that interact with the database.

Role:-

They write application code and handles database connections.

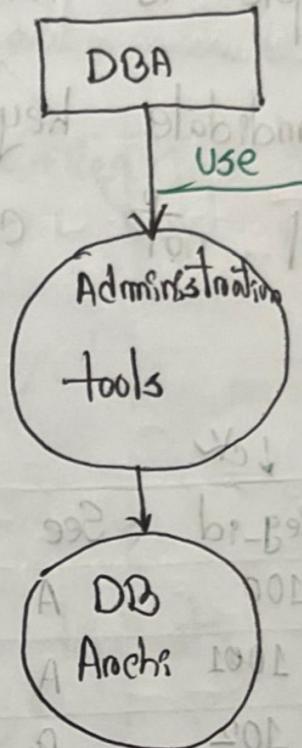


Database Administrators:-

The DBA is responsible for the overall management of the database.

Responsibilities:-

- ① Schema definition
- ② Access control
- ③ Performance tuning
- ④ Backup and recovery
- ⑤ Security



Candidate key :-

The minimal set of attributes that can uniquely identify a tuple is known as a candidate key.

■ Candidate key ^{attributes} can contain null values, but

will not consider same value.

Reg-id	Sec	name	age	Email-id	Phone-no
100	A	Farhan	20		
1001	A	Hossan	22		
102	B	Rabbi	21		
103	B	Nayon	21		
104	B	Hossan	21		
105	A	Rana	23		

■ candidate key এর মধ্যে আকে primary key ফিল্ড
ক্ষেত্র হচ্ছে তা বাইরে থাকি গুলো alternative key.

Theme:

Super key:-

A set of attributes that can uniquely identify each tuple in a relation.

- ① Reg-id
- ② Email-id
- ③ Phone-no

} SK

Sec + name \rightarrow x sk

Sec + Name + Reg Id \rightarrow sk

লেখা করুন

একটি ck must

লাগবে

Primary key:- (not null & not duplicate)

A candidate key chosen to uniquely identify tuples in a relation.

Example:-

'ID' is typically chosen as the primary key.

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foreign keyo-

An ~~other~~ attribute in one relation that refers to the primary key of another relation.

(Parent table-এর মুক্তির একটি কলার হিসেবে child table-এর একটি কলা।)

(student টেবিল এর মুক্তির একটি কলা) -> foreign keyo-

student প্রতিশ্রুতি প্রদর্শন করে গোড়ে যে স্থানে A
র মুক্তির একটি কলা।

Foreign keyo-

ID, এর মুক্তির একটি কলা।

Chapter-3

SQL is a standardized language for interacting with database. It includes-

- DDL
- DML
- Transaction Control

Parts of SQL

DML → Deals with querying, inserting, updating & deleting data.

Integrity Constraints → Defines rules to maintain accurate data.

View definition → allows creating views to simplify querying.

Transaction Control → manage database transactions.

Authorization → sets permissions for accessing data.

Theme:

DDL defines the schema for databases.

Data types for each attribute

Integrity constraints.

Indexes and physical storage.

Domain Types in SQL:-

char(n) - Fixed-length character strings.

varchar(n) = variable length strings.

int, smallint - integer data types

numeric(p,d) - fixed precision numbers

float(n), real, double precision - floating-point numbers.

Creating Table

Create table instructors

ID char (5),

name varchar (20),

dept-name varchar (20),

Salary numeric (8,2)

;

Integrity Constraints

- primary key
- foreign key ----- references table - r
- not null

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Example

Create table instructor
(
ID char (5)
name varchar (20) not null
dept-name varchar (20),
salary numeric (8,2),
primary key (ID),
foreign key (dept-name) references dept
)

Updates to tables

insert

insert into instructor values ('10211', 'Smith', 'Bio', 66000)

① Delete (Removes all tuples)

delete from student

Drop table

drop table n

Alter :-

alter table n add A D → new column data type
↓ ↓
table new column name

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The select Clause (Same as projection operator)

Select clause, যেমন desired attribute এর show করা.

Example:-

Select TITLE from course → যেই table এর কোনো কোনো মান
Key word Key word Key word
 কোনো column মান
 বিকল কোনো নাই

SQL names are case insensitive

Example:- Name ≡ NAME ≡ name

To remove all duplicate tuples:-

Select distinct dept_name from course;

To show all attribute:-

Select * from course;

To rename a selected column:-

Select course_id as C_ID from course;

Theme:

To select more than one column

To rename a column

To do arithmetic operation on a column

Select budget / 13 as Budget_of_1_room, dept_name
from department;

The where Clause (selection operator)

Where clause কোন খুঁতির প্রয়োজন হলে একটি শর্ত দিব।
কর্ণ করে একটি শর্ত দিব।

To find a column with specific condition

using connectives

using comparison operation

Select name, dept_name from student where dept_name =
where dept_name = 'CSE' and tot_marks > 100

Theme:

The from Clause (Cartesian Product)

To do cartesian product:-

Select * from student, department;

To find "building" of "individual person's" using dept-name:

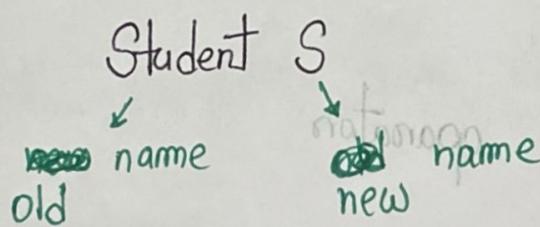
Select name, building from student, department

where student.dept = name = department.dept_name;

The Rename Operation

We rename table on attributes, old-name as new-name

OR



String Operations

Names starting from 'j': (exp: jisan, Juddha etc)

Select name from student where name like '%j%';

Names ending with 'mi': (exp: Urmil, Nawmi etc)

Select name from student where name like '%omi%';

Names which have 'ai' somewhere: (exp: naina, Traini etc)

Select name from student where name like '%a%';

Names which has exactly 4 letters

Select name from student where name like '____';

Names that has 'a' after 1 letter

Select name from student where name like '_a%';

Names that starts with 'N' and has exactly 4 letters

Select name from student where name like 'N____';

escaping special characters

Select name from instructors where name like '100%\$%'
ESCAPE'

Special characters এর টমর backslash দিব কো খাব ESCAPE টমর

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Ordering

For ascending

(to sort first : qx) : 'i' most tribute amount

order by name asc

Using Comparison operators

Select name from student where tot_encl between 50 and 100;

(to print union : qx) : name is 'in' George Paul most tribute amount

: "in" all name starts tribute most Simon total

, starts with most tribute amount

: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

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: "in" all name starts tribute most Simon total

: "in" all name starts tribute most Simon total

Normalization

→ It is the process of organizing the data in the database.

→ It is used to remove or reduce the redundancy from a relation.

একটি দ্বিমুখীয় data এবং
একটি দ্বিমুখীয় value যদি

একটি table এ multiple
time আসে,

মনে করা যাক, একটি table
এর duplicate value থাকে।
তাহলে একটি table
থেকে একটি ক্ষেত্র উৎপন্ন হলে Normalization.
একটি table এ যদি duplicate value থাকে তাহলে
reduce করা হবে।

↓ pk

Emp_id	Emp_name	Contact_no	Dept.	Dept_name
101	A	01789.....	Eng.	FH
102	B	0198.....	HR	MR
103	C	0188.....	Marketing	MJ
104	D	0177.....	Eng	FH
105	E	0154.....	Marketing	MJ
→ 101	A	01789....	Eng.	FH

① Insert Anomaly →

② Update Anomaly →

③ Delete Anomaly

First Normal Form (1NF)

→ An attribute of a relation can't hold multiple values.

■ In a column multi-values not allowed.

■ Attribute Domain should same.

■ Every column will have a unique name

■ Doesn't matter in which order data is stored.

Student_id	Name	Course
101	A	OS, DBMS
102	B	OS
103	C	DBMS
104	D	Algorithm, DBMS



Student-id	Course
101	OS
101	DBMS
102	OS
103	DBMS
104	Algorithm
104	DBMS

Student-id	Name
101	A
102	B
103	C
104	D

Second Normal Form (2NF) :-

- Table should be in 1NF first
- Partial dependency not allowed.
Need to remove Partial Dependency

Dependency বলতে কোনো Non-key এবং column থেকে আসে।
যদি fully depended হয়ে pk এর দ্বা দ্বারা নির্ভর করে তবে এটি অসম্ভব।

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Theme:

Student-ID	Name	Address	Course-ID	Course-Name	Credit	Grade
101	A	DH	1	C	3	A+
101	A	DH	2	Java	3	A
102	B	GN	1	C	3	B
102	B	GN	2	Java	3	A



C-id	C-Name	Credit
1	C	3
2	Java	3

S-id	Name	Add
101	A	DH
102	B	GN

S-id	C-id	Grade
101	1	A+
101	2	A
102	1	B
102	2	A

Third Normal Form (3NF)

- ① Table must be in 2NF form
- ② Transitive Dependency not allowed.

Course-Name	Teacher-id	Teacher Name	Credit
C	1	Farhan	3
OS	2	Hossan	3
DBMS	1	Farhan	3
CLAB	1	Farhan	2



C-Name	T_id	Credit
C	1	3
OS	2	3
DBMS	1	3
CLAB	1	2

T_id	T_Name
1	farhan
2	Hossan

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Chapter-6

Design Phases

Client description → ERD → Schema Diagram → Table with constraints

logical design

↳ Business decision

↳ Computer science decision



Entity Sets (represented by attributes)

An entity is an object that exists and is distinguishable from other objects.

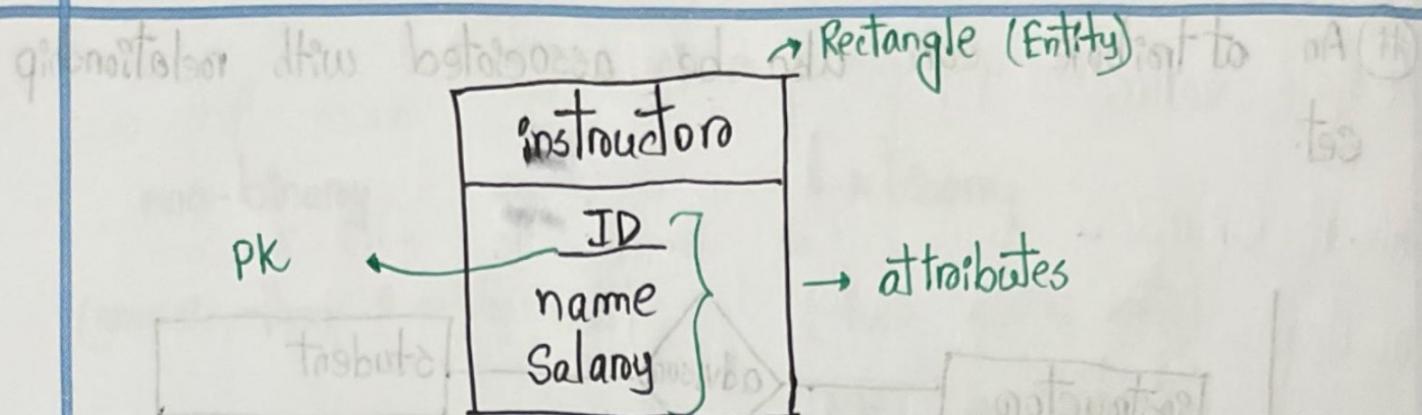
Example:- specific person, company, event, plant

An entity set

↳ set of entity

↳ same properties

Ex: set of all persons, companies, holidays.

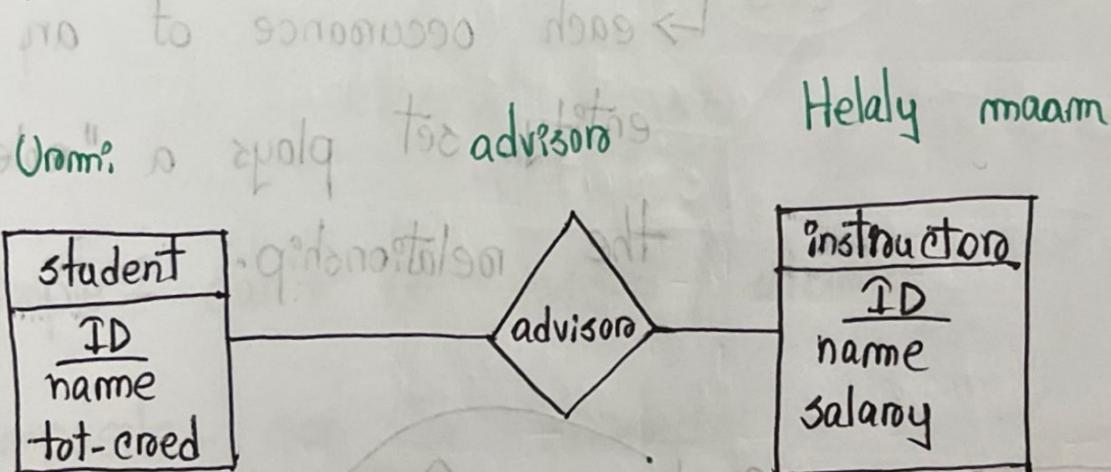


Relationship sets

represented
by diamond

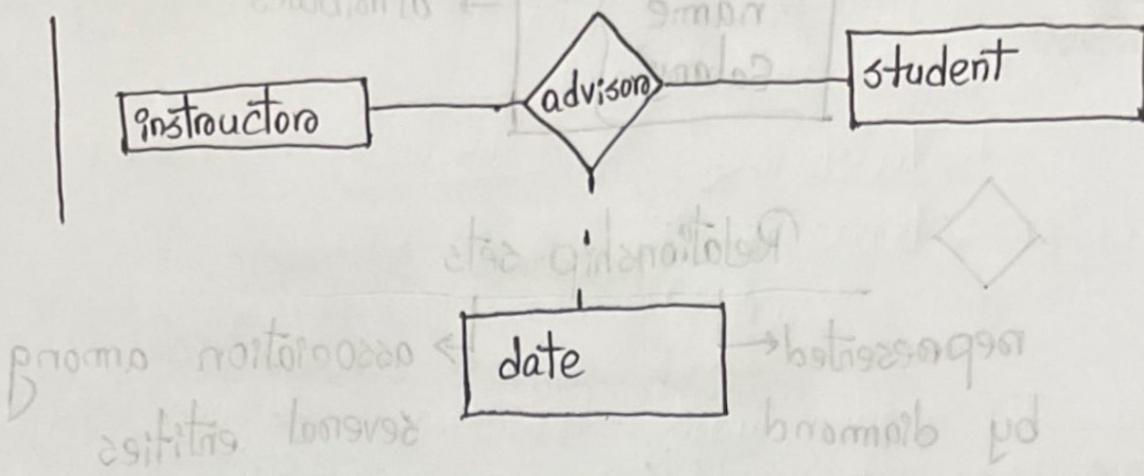
association among
several entities

Exp:



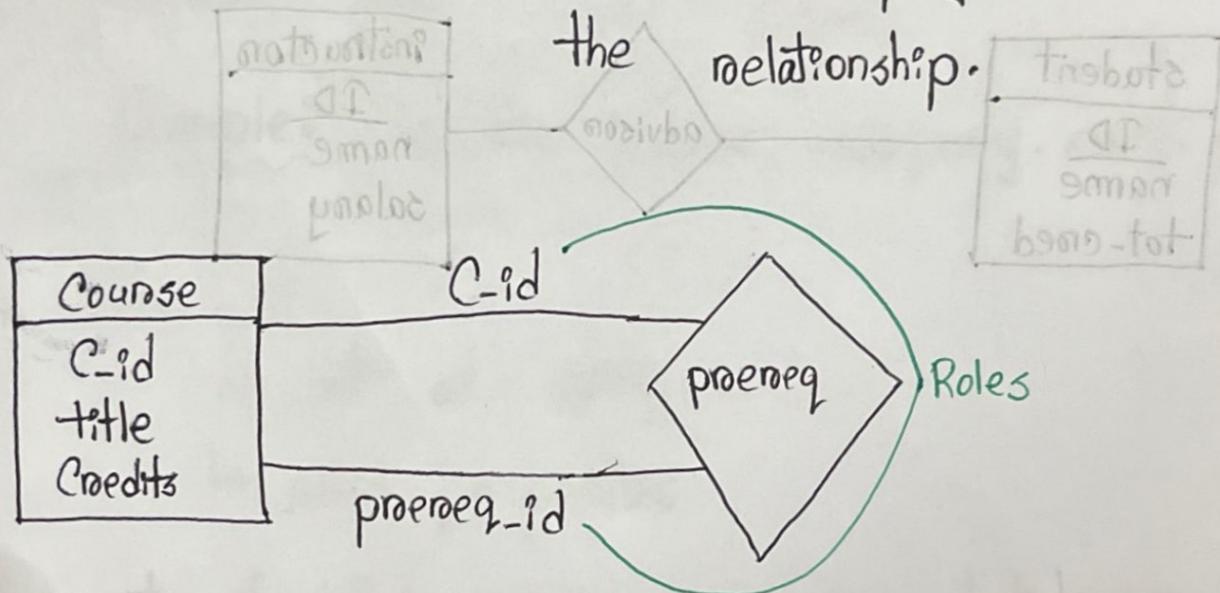
Theme:

- # An attribute can also be associated with relationship set.

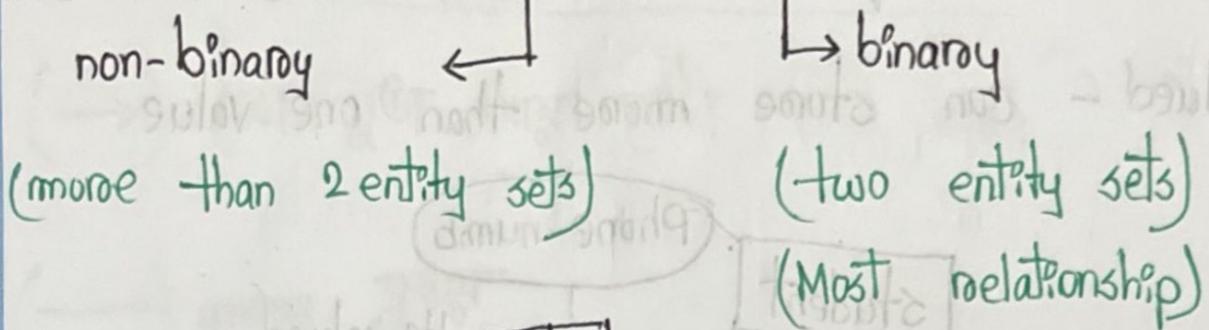


Role

→ each occurrence of an entity set plays a "role" in the relationship.



Degree of a relationship set



project

instructors

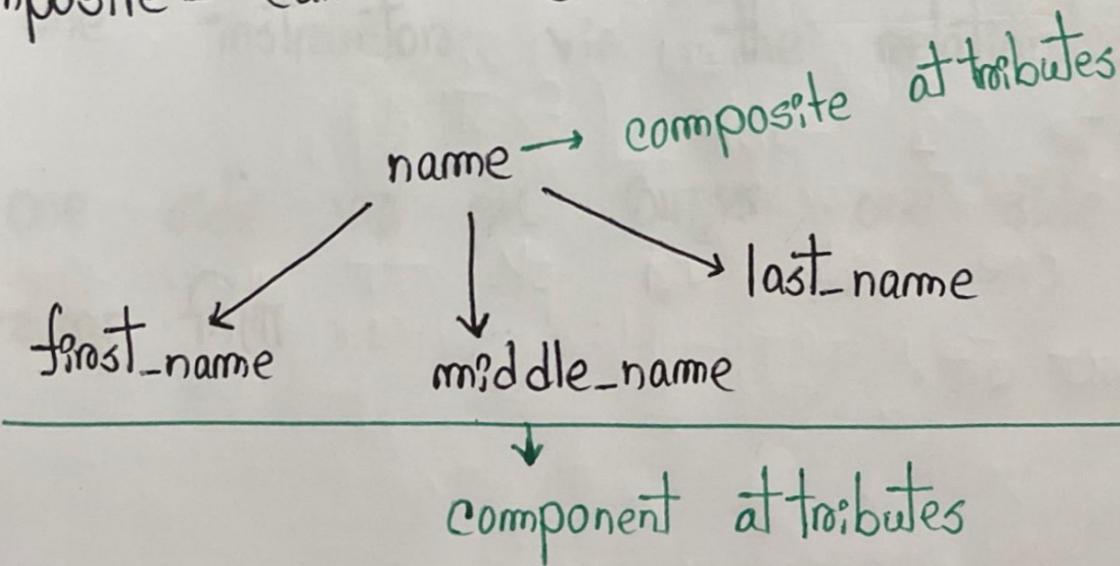
Student

proj-guide

Attribute type

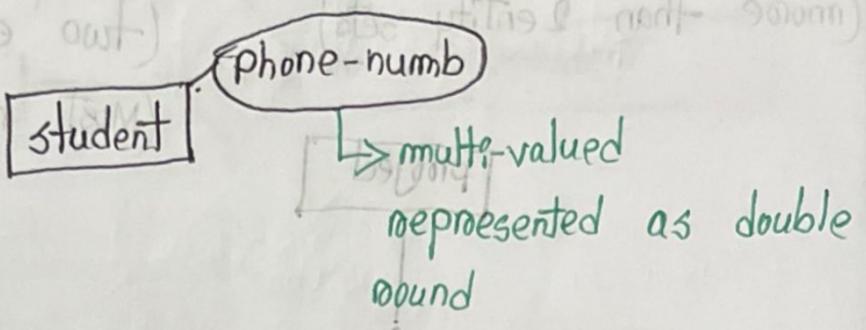
Simple - can't be divided / atomic

Composite - can be divided

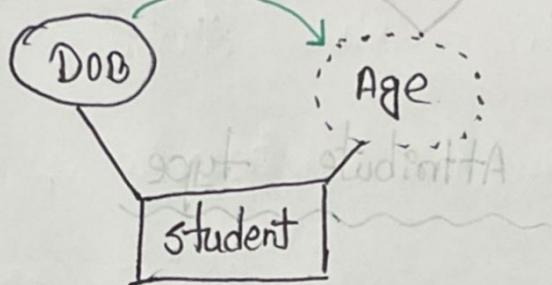


single-valued - unique / can't store more than one value

multi-valued - can store more than one value



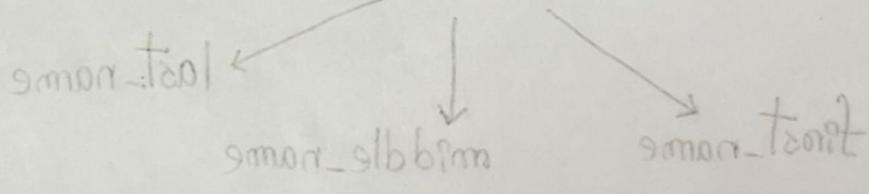
derived - can be computed from other attributes



Composite - can be divided into smaller components

Composite - can be divided into smaller components

student to composite



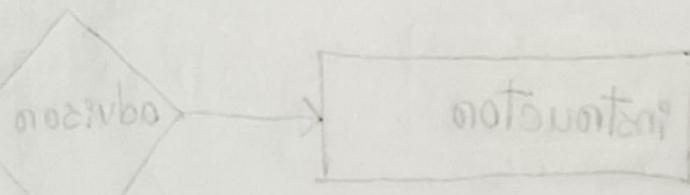
student to transposed

Mapping Cardinalities

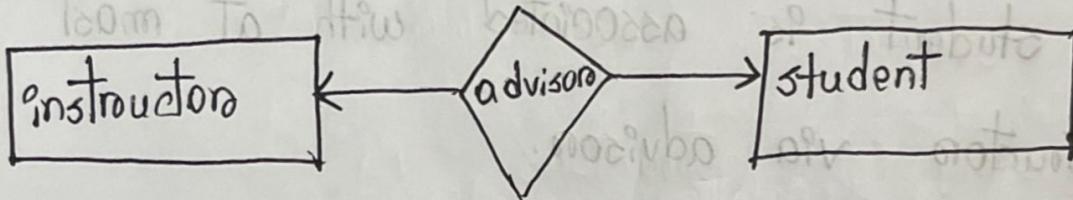
(→) একটি শুটর আছে তার্ক একটি one side

(—) many বুম্বাম্ব

One to one:-



each instance of one entity is associated with exactly one instance of another entity.

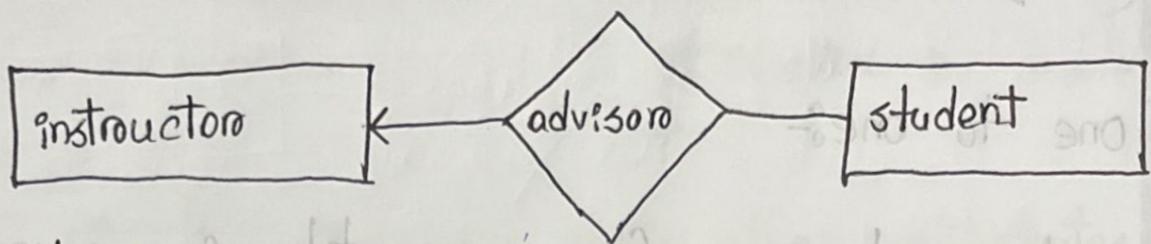


a student is associated with at most one instructor via the relationship advisor.

one side এক pk টাকের one side এক Fk
বান্দা টিকা

one - to - many / many - to - one

one instance of an entity is associated with multiple instances of another entity.



an "instructors" is associated with several (including 0) students via "advisors".

a student is associated with at most one "instructors" via "advisor".

O 2 M / M 2 O relationship

① Create 2 table (Existing)

one side pk many side-pk

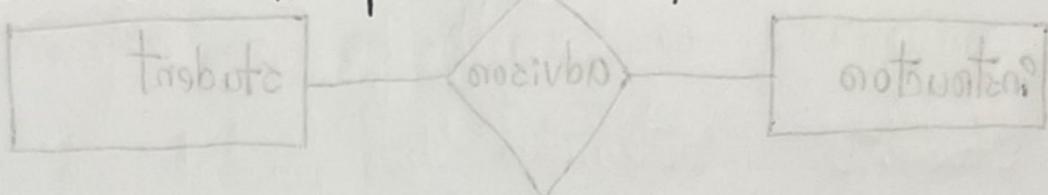
Theme:

ପ୍ରଯୋଗେ one side ଏବଂ data repeat ହେତୁ ଏହି data

repetition ହେଲୁ ଥାଏ ମଧ୍ୟ ମଧ୍ୟ ମଧ୍ୟ ମଧ୍ୟ ମଧ୍ୟ ମଧ୍ୟ table ଏବଂ

⑩ create ୩ table (new one)

both side ଏବଂ pk ନିଷ୍ଠା କରିବାରେ ଏହି table.



(ଓ ପିଲାଙ୍କରୁ) ଲୋକେ କ୍ରୀତିମାଣ ଏବଂ ଜୀବନକାଳୀନ ନାମରେ
 ଲୋକେ କ୍ରୀତିମାଣ ଏବଂ ଜୀବନକାଳୀନ ନାମରେ

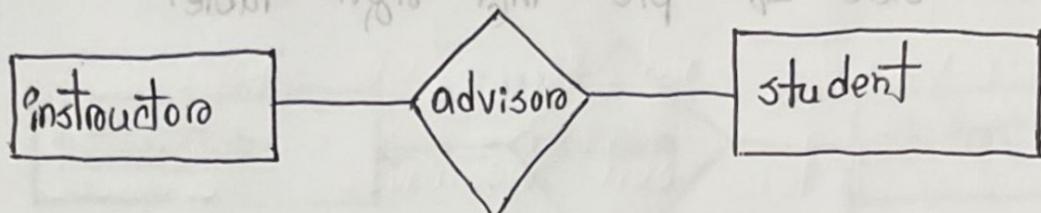
ଲୋକେ କ୍ରୀତିମାଣ ଏବଂ Trabute ଏବଂ

ଲୋକେ କ୍ରୀତିମାଣ ଏବଂ ଜୀବନକାଳୀନ (ଓ ପିଲାଙ୍କରୁ)

Theme:

many to many

each entity can be associated with multiple instances of ^{other} entity



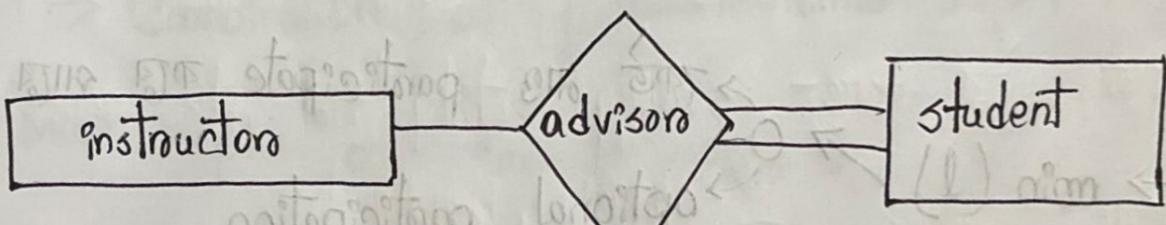
- # an instructors is associated with several (possibly 0) students via advisors.
- # a student is associated with several (possibly 0) instructors via advisor

নতুন একটি table গৈরিক কর্তৃত থবে both side এর pk
ছিন্ন যা ক্ষয় table কে refers করে।

চুল্টি pk ছিন্ন নতুন table এর composite pk বানাতে
থবে।

Total participation

every entity in the entity set must participate in at least one relationship in the relationship set.



all students need an instructor, meaning participation is total for students in this relationship.

Theme:

Partial Participation

Some entities in the entity set may not participate in any relationship in the relationship set.

- # An instructors may not be an advisor for any students. (instructors has partial relationship with advisor)

Minimum & Maximum Cardinality

↳ (l.....h notation)

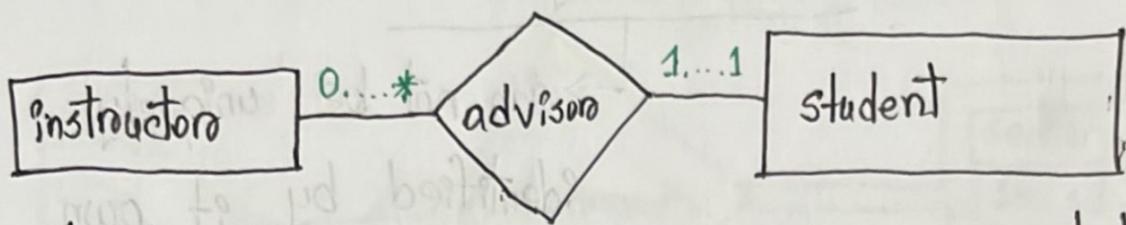
↳ minimum (l) and max (h) কৃত্যাব এন্টিটি সেট এন্টিটি
 রেলিশপ-এ প্রতিক্রিয়া করবে।

↳ min (l) → ০ কর্তৃ নাও প্রতিক্রিয়া করতে পারে
 optional participation

↳ max (h) → 1 → ∞ কর্তৃ একজনের প্রতিক্রিয়া
 single participation

* → 1 → ∞ কর্তৃ অবশ্যই করতে পারে
 unlimited participation

Theme:



→ instructors can advise 0 or more students.

→ a student must have 1 advisor, can't have multiple advisors.

Primary key

Relationship sets

(Combining pk/
attributes of
relationship itself)

← Entity sets

(unique for each each
entity)

M2M → Combination of pk → minimal super key

O2M / M2O → pk of "many" side → " "

O2 → any pk from both entity → " "

Weak Entity Sets

↳ can not be uniquely

identified by its own

attributes (can't form PK)

↳ relies on strong entity sets

↳ participate in special relationship.

weak entity sets existence depends on identifying entity set (strong entityset)

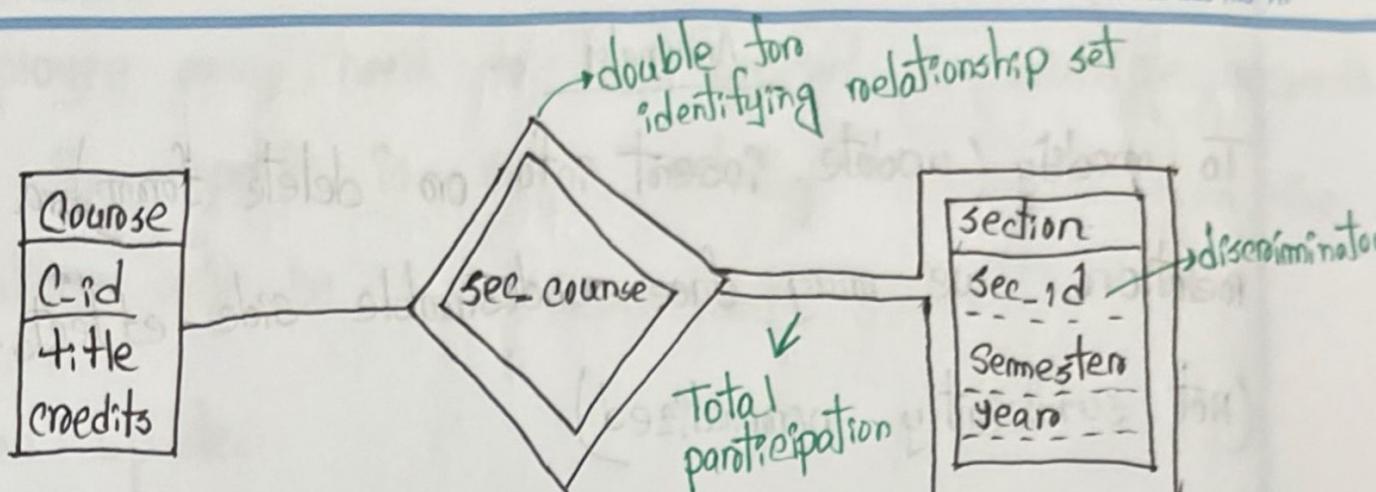
weak entity set এর attribute গুলোকে dissemination

বন্টে, এটার জায়ে strong entity এর pk combine

করলে weak entity uniquely identify করা থার্জ.

one side এর pk কে many side-এ add করত

বিবে।



SL#	E-obj	E-id
प्राप्ति	प्राप्ति	H10
प्रतिक्रिया	प्रतिक्रिया	H10
प्राप्ति	प्राप्ति	281

Theme:

Anomaly

To modify (update, insert info or delete from) a relation, we may face undesirable side-effects.
 (Not sufficiently normalized)

Update Anomaly (addressed by ensuring data is stored in only one place)

It occurs when the same data is stored in multiple rows and update must be applied to all occurrences. If some occurrences are not updated, the database ends up in inconsistent state.

E_id	E_add	Skill
214	Razabazar	Beauty
214	Razabazar	Handwriting
186	Panthalpath	Running

a change in the address for a particular

Theme:

employee may need to be applied to multiple records. If the update is only partially successful then the relation will be in inconsistent state.

After update:

E_id	E_add	Skills
214	Razabazar	Beauty
214	Monipura	Handwriting
186	Panthapath	Running

Theme:

Date: / /
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Insertion Anomaly (reduced by creating separate tables)

Occurs when certain data can't be inserted into the database without the presence of other data.

FID	Ename	Date_hire	Course
212	Helaly	Feb 2015	CSF 201
213	Nadeem	Jan, 2008	CSF 211

200	Nayeema	Jan, 2019
-----	---------	-----------

Until the new faculty member, Nayeema, is assigned to teach at least one course, his/her details can't be recorded. We have to assign course to null. This phenomenon is known as insertion anomaly.

Deletion Anomaly:- (keep important data in independent tables).

Occurs when the deletion of data inadvertently results in the loss of additional, valuable data.

F-id	F-name	Hire-date	C-code
201	Hedaly	Feb, 1995	CSE 201
205	Tahira	Jan, 2000	CE 101

→ delete

All information about Hedaly is lost if she temporarily assigned to another course. This is deletion anomaly.

Transaction

Theme:

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- # Suppose, you were withdrawing from an ATM booth. You put the card inside the machine, gave password correctly & clicked 'OK' button. After few minutes, the machine failed & gave your card back to you. Which ACID property(s) can play role for the management of this transaction? Explain that property.

101 30 000 00 101 200

101 30 000 00 101 200

- # Suppose, you were withdrawing money from an ATM booth in Dhaka. At the same time, your father was depositing cash money in another branch. Which ACID property(s) can play role for the management of this transaction? Explain that property.

- # Discuss the four transaction properties of ACID in transaction management.

Ans:-

Atomicity:-

A transaction is an indivisible unit of work that is executed entirely or not at all. It ensures that the database is never left in a partial or inconsistent state due to transaction failures. The database system keeps track (on disk) of the old values of any data on which a transaction performs a write. The information is written to a file called the log. If the transaction does not work complete its execution, the database system restores the old values from the log to make it appear as though the transaction never executed.

Theme:

Example:-

A fund transfer transaction deducts \$50 from account A and adds \$50 to account B.

If a system crash occurs after deducting \$50 from account A, but before adding it to account B, atomicity ensures that either both operations complete or nothing takes effect.

Consistency:-

A transaction must take the database from one consistent state to another, maintaining all defined rules. It ensures that the database remains valid according to its defined rules & relationship.

Ensuring consistency for an individual transaction is the responsibility of the application programmers who codes the transaction. This task may be facilitated by automatic testing of integrity constraints.

Example:-

A transaction that transfers funds between accounts must ensure that the sum of all balances before & after the transaction remains the same.

Isolation

Transactions are executed as if they are the only operations happening in the system, preventing interferences from other concurrent transactions. It ensures that concurrent transactions do not affect each other's results. A way to avoid the problem of concurrently executing transactions is to execute transactions serially. Ensuring the isolation property is the responsibility of a component of the database system called the concurrency-control system.

Theme:

Example:-

Two transactions, T1 & T2

T1 reads Account A's balance, \$50 and writes back.

T2 at the same time reads the balance & updates it.

Without isolation, T2 might see intermediate results or overwrite T1's updates.

Durability:-

Once a transaction commits, its changes must be permanent, even in the event of a system failure. It ensures that committed data is not lost. We can guarantee durability by ensuring that the updates carried out by the transaction have been

Theme:

written to disk before the transaction completes.

The recovery system of the database is responsible for ensuring durability.

Example:-

After a fund transfer transaction completes, the updated balances of account A & account B must be stored persistently so that they survive crashes or power failures.

A transaction is a unit of program execution that accesses and possibly updates various data items.

Transaction access data:

Read(x):

it allows transactions to access the current value of data item without modifying it.

Write(x)

it allows transactions to make changes to data that are eventually stored in the data that are eventually stored in the database.

Example:-

T_1 : read (A)

A: A - 50;

read (B)

B: B + 50;

write (A)

write (B);

Theme:

Draw the state diagram of transaction states with a brief discussion.

Ans:-

We establish a simple abstract transaction model. A transaction in a database system goes through a series of states from its initiation to its completion. A transaction must be in one of the following states:

Active:-

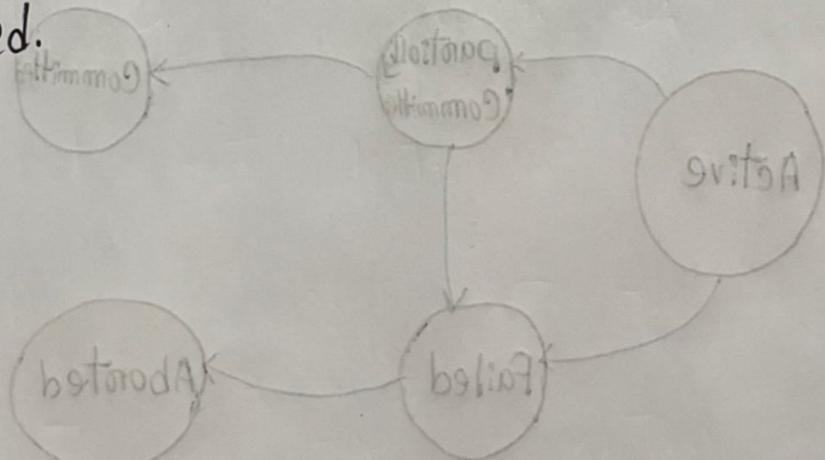
The initial state of a transaction while it executes its operations. If failure occurs during execution, it transitions to the failed state.

Partially Committed:-

The transaction executes but database has not recorded permanently. This state exists because some operations may still be buffered in memory. Even in this state if a failure occurs, it transitions to the failed state.

Failed:-

The transaction has encountered an error or failure during ~~executio~~ execution. The database is in inconsistent state due to partial updates made by the transaction. The transaction transitions to the aborted state after rollback is completed.



Theme:

Aborted :-

After the transaction has been rolled back & the database has been restored to its state prior to the start of the transaction. From hence the transaction may restart or terminate.

Committed :-

The transaction has successfully completed all of its operations; & the changes have been made permanent in the database. This is the final state & can't be rolled back.

