

## LSML #5 Recap

Lin Reg  $x = [x_1, \dots, x_d] \quad d > 10^6$

| Count-min Sketch

Хотим оценивать частоты

$$X \quad \|X\| = n$$

$$\sqrt{\alpha_{ct}} = (\alpha_{1,ct}, \dots, \alpha_{n,ct}) \quad \checkmark$$

$$\forall \alpha_{i,ct} = 0$$

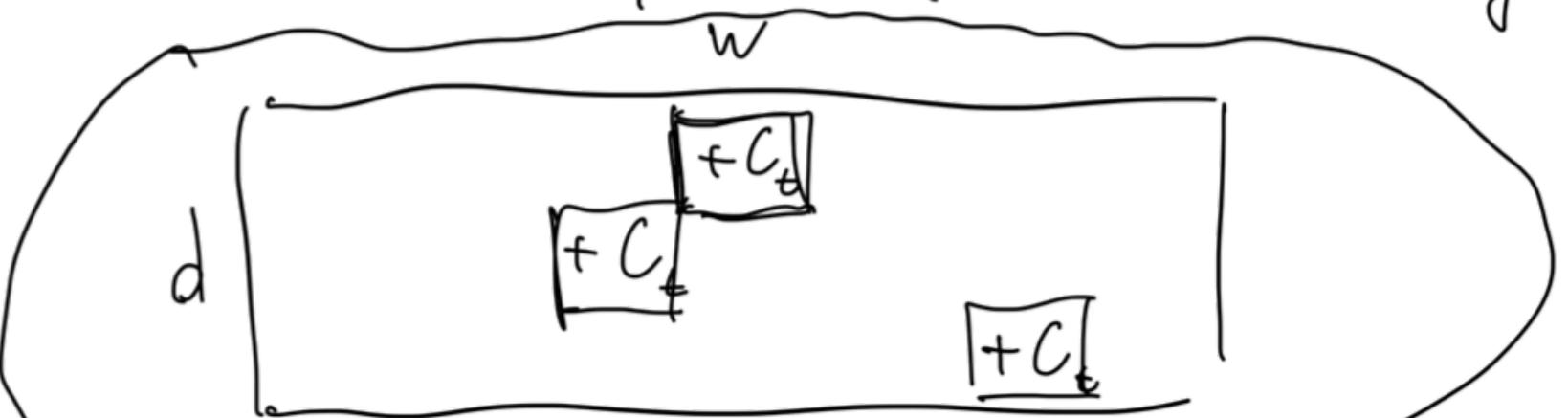
$$\begin{cases} \alpha_{i,t}(t) = \alpha_{i,t-1} + c_t \\ \alpha_{i,t} = \alpha_{i,t-1} \quad \forall i \neq i_t \end{cases}$$

$\boxed{d}$  хешей no w занченоу

$$\sqrt{h_i(x)} \quad i \in [1, \dots, d] \quad - хеши независимы$$

$$i_t \rightarrow h_{i_t} \sim h_{j_t} \quad \checkmark$$

| d-строк u w столбцов



ошибки одна ( $\pm \varepsilon$ ) с вероятностью  $\delta$

$$w^* = \lceil \frac{e}{\varepsilon} \rceil$$

$$d^* = \lceil \frac{\ln 1/\delta}{\varepsilon} \rceil$$

| Теорема  $\underline{\alpha_i \leq \hat{\alpha}_i}$   $\overline{\min_j \text{count}[j, h_j, c_i] = \hat{\alpha}_i}$

$$P(\hat{\alpha}_i > \alpha_i + \varepsilon \cdot \|\hat{\alpha}\|_1) \leq \delta$$

$$P\left(\frac{\hat{\alpha}_i}{\|\hat{\alpha}\|_1} > \frac{\alpha_i}{\|\hat{\alpha}\|_1} + \varepsilon\right) \leq \delta$$

$$\boxed{\sum_{i,j,k} I_{i,j,k}} = \begin{cases} 1 & \text{если } c_i \neq k \wedge (h_j, c_i) = h_j(c_k) \\ 0 & \text{иначе} \end{cases}$$

$$\mathbb{E}(\sum_{i,j,k} I_{i,j,k}) = P(h_j, c_i) = \frac{1}{w} \approx \frac{\varepsilon}{e}$$

$$\boxed{\sum_{i,j} X_{i,j} = \sum_{k=1}^n \sum_{i,j} I_{i,j,k} \Rightarrow \text{count}[j, h_j, c_i] = \alpha_i + X_{i,j}}$$

min count  $> \alpha_i$

$$\mathbb{E}(X_{i,j}) = \mathbb{E}\left(\sum_{k=1}^n I_{i,j,k} \cdot \alpha_k\right) = \sum_{k=1}^n \alpha_k \cdot \mathbb{E}(I_{i,j,k}) \leq \frac{\varepsilon}{e} \cdot \|\hat{\alpha}\|_1$$

$$P(\hat{\alpha}_i > \alpha_i + \varepsilon \cdot \|\hat{\alpha}\|_1) = P(\forall j \text{ min count}[j, h_j, c_i] > \alpha_i + \varepsilon \cdot \|\hat{\alpha}\|_1) \quad \otimes$$

$$\otimes P(\forall j \text{ min count}[j, h_j, c_i] > \varepsilon \cdot \|\hat{\alpha}\|_1) \leq P(\forall j \text{ } X_{i,j} > \varepsilon \cdot \mathbb{E}(X_{i,j})) < e^{-d} \leq \delta$$

$$\boxed{P(X \geq t) \leq \frac{\mathbb{E}(X)}{t}} \quad \forall t > 0$$

$$-d \leq \ln \delta \quad d \geq -\ln \delta$$

~~доказать~~

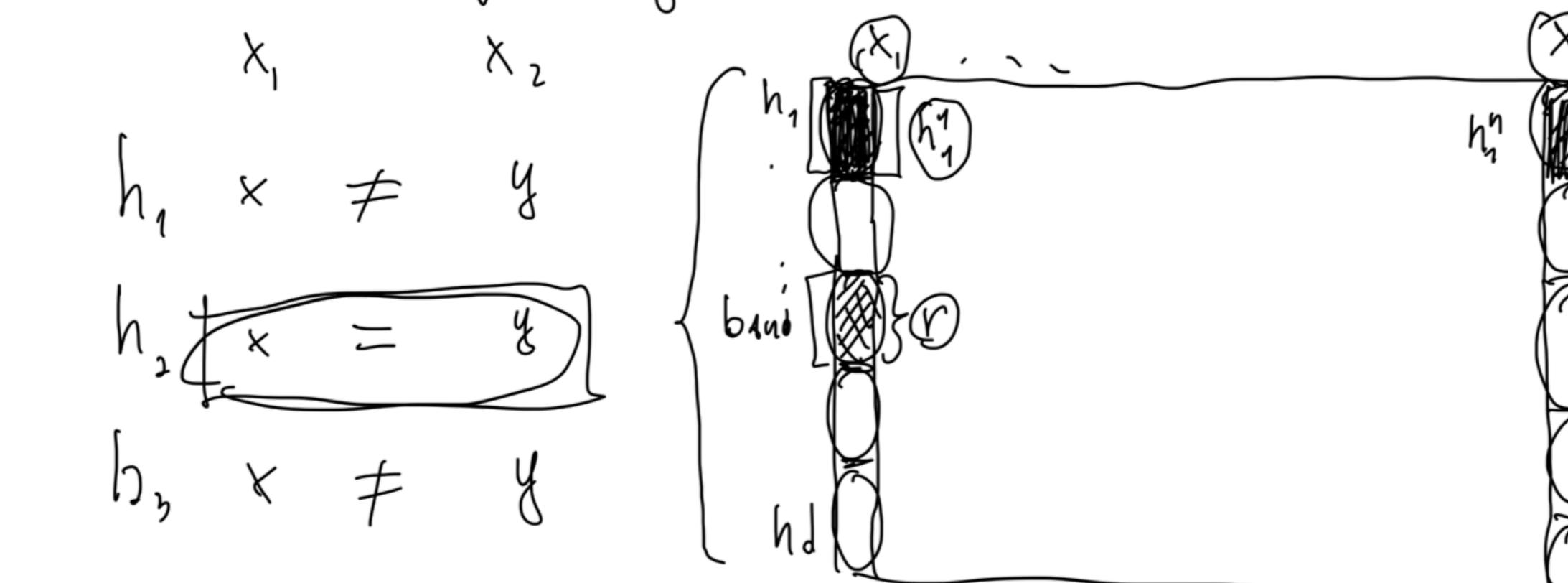
$$d \geq \ln \frac{1}{\delta}$$

$$d \geq \ln \left( \frac{1}{\delta} \right)$$

$$d^* = \lceil \ln \left( \frac{1}{\delta} \right) \rceil$$

## LSH locally sensitive hashing

Наше: схема  $h(x)$ : в один элемент  
хеша попадут одинаково  $\exists i \neq j$



$$b+r=s$$

$$h_i = [1 \dots k]$$

$$\left(\frac{1}{k}\right)^r$$

$$b \approx 20/10$$

$$r \approx 5-10$$

$$(1 - \left(\frac{1}{k}\right)^r)^B$$

$$0.8888$$

| Online learning

X  $\|X\| > HOD$  Datasets known

loss  $\rightarrow \min \equiv \text{offline learning}$

$$\boxed{\text{Regret} = \left[ \sum_{i=1}^T \underline{l_i(c_w)} - \underline{l_i(c^{*})} \right]}$$

$$x_1 \rightarrow f_1 \rightarrow y_1 \rightarrow w_2 \quad x_2 \rightarrow f_2 \rightarrow y_2 \rightarrow w_3$$

| FTL = follow the leader

$$\boxed{w_{t+1} = \arg \min_w \left[ \sum_{i=1}^t \underline{l_i(c_w)} \right]} \quad \begin{array}{l} \text{-- good и хоро} \\ \text{-- best of letator} \end{array}$$

$$w_{t+1} = \arg \min_w \left( \sum_{i=1}^t \underline{l_i(c_w)} + r_i(c_w) \right)$$

$$r_i(c_w) = \lambda \cdot \|w - w_i\|^2$$

| Strong convexity!

$$f(y) \geq f(x) + \nabla f(x)(y-x) + m \cdot \|y-x\|^2$$

$$\boxed{\underline{l_i(c_w)} = \underline{l_i(c_{w_i})} + \nabla l_i^T(c_{w_i})(w - w_i)}$$

$$\boxed{\text{① Regret} = \sum_{i=1}^T \underline{l_i(c_w)} - \underline{l_i(c^{*})}}$$



$$\boxed{\underline{l_i(c_w^*)} = \underline{l_i(c_{w_i})} + \nabla l_i^T(c_{w_i})(w^* - w_i)}$$

$$\boxed{\underline{l_i(c_w^*)} \geq \underline{l_i(c_{w_i})} + \nabla l_i^T(c_{w_i})(w^* - w_i) + m \cdot \|w^* - w_i\|^2}$$

$$m > 0$$