Inspection of Julia GC for multithreaded programs

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Intro to Julia

What: High-level, high-performance, dynamic programming language - <u>aimed</u> to employ

the best of both worlds

Who: GP, but mainly for scientific community - <u>numerical</u> analysis

~10,000 companies, 1500 universities[1]

When: 2012, MIT

Features: Multiple dispatch, dynamic typing, parallel and distributed computing, no need for

a vectorized code, package management

Compiler: Written in Julia, C/C++, JIT



VS



VS



Syntax in short

Q: What makes Julia to perform close to C?

A: (1) Julia features + (2) Efficient user code

Features:

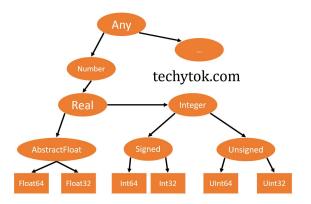
Dynamic type inferring

```
d = Int[1,2,3,4,5] vs d = [1,2,3,4,5]
```

- Multiple dispatch (don't overuse)
- Using function profiling

User:

- Type stability to infer types beforehands correctly $NO d = 1 \dots d = 2.3 \dots d = [1,2,3,4,5]$
- Common sense: no global vars, static arrays, @inbounds, etc[2]
- Identifying bottlenecks



Types hierarchy in Julia [2]

```
function test1(x::Int)
    println("$x is an Int")
end

function test1(x::Float64)
    println("$x is a Float64")
end

function test1(x)
    println("$x is neither an Int
nor a Float64")
end
```

Multiple dispatch example [2]

Motivation

Source of issues:

- 1) > Non parallel, mark and sweep GC
 - > multi threading
 - > more data to collect during "stop the world" (stw)
 - > higher "stw" latency
 - > lower average performance
- Some non-allocational threads might not never yield to a safe(precise) point
 - > Hence, allocating threads might stall
 - > lower average performance

Motivation

Solutions:

- Inspecting the GC to identify which specific applications are more vulnerable to the "stw" latencies
- 2) Identifying which parts of the GC collections and sweeping relatable for this

GC in Julia

- Julia's garbage collector algorithm is called mark and sweep.
- It is <u>precise</u>, <u>generational</u>, <u>non-compacting</u>, <u>non-parallel</u>.
- The mark phase: Where all objects that are not garbage are found and marked so
- The sweep phase: where all unmarked objects are cleaned
- Starts <u>earlier</u> than memory overflow
- The top Julia GC source file is GC.c

Triggering manual GC

- When the garbage collector detects that the program can no longer access the object, then it will run the finalizer, and then collect (free) the object.
- Note that the garbage collector can still access the object, even though the program cannot.

Example of Manual GC

```
julia> r = Ref(0)
Base.RefValue{Int64}(0)

julia> finalizer(r) do r
    println(r)
    end
Base.RefValue{Int64}(0)

julia> r = nothing

julia> GC.gc()
Base.RefValue{Int64}(0)
```

DTrace and BPFTrace

- DTrace(Solaris originally) and bpftrace(Linux) are tools that enable lightweight instrumentation of processes.
- You can turn the instrumentation on and off while the process is running, and with instrumentation off the overhead is minimal.

Tracing interfaces

	kernel	userland
static	tracepoints	USDT* probes
dynamic	kprobes	uprobes

[5]

- We will focus on userland for now, since that is the most useful feature for understanding applications, like Julia
- USDT User-level Statically Defined Tracing

DTrace and BPFTrace

Using nm we can list all of the exported runtime functions of Julia.

```
nm -D /usr/lib/julia/libjulia-internal.so | grep jl_gc_
00000000000015215 T jl_gc_add_finalizer
0000000000015221 T jl_gc_add_finalizer_th
000000000001522d T jl_gc_add_ptr_finalizer
0000000000015239 T jl_gc_alloc
000000000015245 T jl_gc_alloc_0w
0000000000015251 T jl_gc_alloc_1w
000000000001525d T jl_gc_alloc_2w
0000000000015269 T jl_gc_alloc_3w
0000000000015275 T jl_gc_allocobj
0000000000015281 T jl_gc_alloc_typed
000000000001528d T jl_gc_big_alloc
```

DTrace and BPFTrace

An alternative way to find all valid uprobe is to use bpftrace:

```
sudo bpftrace -l 'uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_*'

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_add_finalizer

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_add_finalizer_th

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_add_ptr_finalizer

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_0w

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_1w

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_2w

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_3w

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_page

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_alloc_typed

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_allocobj

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_allocobj

uprobe:/usr/lib/julia/libjulia-internal.so:jl_gc_big_alloc
```

Using uprobe[5]

 The function allocator is going to run forever and allocate an array of size N bytes.

```
function allocator(range)
while true
N = rand(range)
buf = Array{UInt8}(undef, N)
end
end
```

Running it: julia -L allocator.jl -e "allocator(64:128)"

Results: uprobe

- Memory allocation distribution by process ID
- This triggers all the processes using the library, but it is possible to trigger only by PID with -p PID

Using USDT probes[6]

- Enabling support:
 Using WITH DTRACE=1 in Make.user file
- Available probes, uprobes.d:
 - •julia:gc_begin: GC begins running on one thread and triggers stop-the-world.
 - •julia:gc__stop_the_world: All threads have reached a safepoint and GC runs.
 - •julia:gc_mark_begin: Beginning the mark phase
 - •julia:gc__mark_end(scanned_bytes, perm_scanned): Mark phase ended
 - •julia:gc__sweep_begin(full): Starting sweep
 - •julia:gc__sweep_end(): Sweep phase finished
 - •julia:gc_end: GC is finished, other threads continue work
 - •julia:gc_finalizer: Initial GC thread has finished running finalizers

Using USDT probes

```
using Base.Threads
fib(x) = x <= 1 ? 1 : fib(x-1) + fib(x-2)
beaver = @spawn begin
    while true
        fib(30)
        GC.safepoint()
    end
end
allocator = @spawn begin
    while true
        zeros(1024)
    end
end
wait(allocator)</pre>
```

- Run this using julia -t 2 (This means we are using 2 threads)
- @spawn allocates a Task to an available thread and executes

Results: USDT probe for "stop the world"

```
ubuntu@euca-10-1-4-194:~/julia$ sudo bpftrace --usdt-file-activation contrib/bpftrace/gc stop_the_world_latency.bt Attaching 4 probes...

Tracing Julia GC Stop-The-World Latency... Hit Ctrl-C to end.
^C
```

• Latency distribution of the stop-the-world phase in the executed Julia process.

Contents of gc stop the world latency.bt

```
#!/usr/bin/env bpftrace
BEGIN
    printf("Tracing Julia GC Stop-The-World Latency... Hit Ctrl-C to end.\n");
usdt:usr/lib/libjulia-internal.so:julia:gc__begin
   @start[pid] = nsecs;
usdt:usr/lib/libjulia-internal.so:julia:gc__stop_the_world
/@start[pid]/
   @usecs[pid] = hist((nsecs - @start[pid]) / 1000);
    delete(@start[pid]);
END
    clear(@start);
```

Using USDT probes

```
using Base.Threads

fib(x) = x <= 1 ? 1 : fib(x-1) + fib(x-2)

beaver = @spawn begin
    while true
        fib(30)
        GC.safepoint()
    end
end

allocator = @spawn begin
    while true
        zeros(1024)
    end
end</pre>
```

wait(allocator)

- GC.safepoint() should be called from a non-allocating thread to make sure that all threads are checked for a precise state.
- It triggers segmentation fault by loading from a protected memory space under the hood and forces GC to execute.
- Future Julia versions expected to include built-in function entry GC.safepoints()

Project challenges

Never use M1 chip Mac for GC tracing!:)

Using linux servers with sudo access is a prerequisite

References

- [1] https://en.wikipedia.org/wiki/Julia (programming language)#History
- [2] https://techytok.com/code-optimisation-in-julia/
- [3] https://juliacomputing.com/
- [4] https://github.com/JuliaLang/julia
- [5] https://vchuravy.dev/notes/2021/08/bpftrace/
- [6] https://docs.julialang.org/en/v1.8-dev/devdocs/probes/#Available-probes

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