

25, 10, 2017

(nkr-7)
(n-1)

Content review

Divides, gcd, lcm

Union, intersection, subset

Cardinality, power set, Cartesian product

Venn diagrams, set operation laws
(boolean algebras)

$\lambda, \Sigma^*, \Sigma^{\leq k}, \Sigma^k$

lexicographic, lex ordering

Functions & relations

Relation $R : R \subseteq A \times B$
 $|R| = 2^{mn}$

Converse $R^{\leftarrow} \subseteq B \times A, R^{\leftarrow}$

Image $R(x) \subseteq B$

Function $f: A \rightarrow B$, ~~1-1~~ $1 \rightarrow 1$ relation

Image

Injection, surjection, bijection

Inverse $f^{-1}: B \rightarrow A$ iff. f is bijection

Inverse image $f^{\leftarrow}(x)$ image of x under
converse of f
exists for any function

Binary relations

Reflexivity, Symmetry, AS,
transitivity

partial order: R, AS, T.
Equivalence: R, AS, T

glb, lub.

Hasse diagram

Topological sort

Graphs

Basic definition

Eulerian circuit/path, Hamiltonian
circuit/path
edge once
vertex once
0 or 2 vertices of odd.

Graph coloring, chromatic number

Planarity

K_5 , $K_{3,3}$

Logic

Conjunction, Disjunction, Negation,
wff.

logical equivalence \equiv
logical entailment \models

CNF/DNF Karnaugh maps

Boolean algebras

Satisfiable

φ satisfiable
if some assignment makes φ true

φ tautology if all assignment makes φ true

CNF some
Derive Φ a formula from literals
 $p, \neg p, q, r$
which is not satisfiable.

$$\langle \neg p \wedge \neg p \rangle \wedge \langle q \wedge r \rangle$$

$$p \wedge \neg p$$

It is in CNF b/c it is a conj. of disj.
of literals p and $(\neg p)$

It is not satisfiable according to
De Morgan's laws. $p \wedge \neg p = \perp$

$$\frac{b' - a'}{1}$$

Show that logical equivalence is an
equivalence relation.

Let $\varphi \equiv \psi$ be arbitrary formulae

(R): $\varphi \equiv \varphi, \psi \equiv \psi$ ~~for all~~
for all formulae φ, ψ .

Take any evaluation $v, v(\varphi) = v(\psi)$
 $\varphi \equiv \psi$

\therefore It is reflexive

(S) Suppose $\varphi \equiv \psi$ be arbitrary formulae

for any evaluation $v(\varphi) = v(\psi)$
 $v(\psi) = v(\varphi)$
 $\psi \equiv \varphi$

(T) Suppose $\varphi \equiv \psi, \psi \equiv \theta$ for arbitrary formulae

For any evaluation $v(\varphi) = v(\psi), v(\psi) = v(\theta)$
 $v(\varphi) = v(\theta)$
 $\varphi \equiv \theta$

Induction and recursion
one or two base cases,
one or two inductive steps

Structural induction

Running time of algorithms

Big-O

~~AD~~
Master Theorem

counting

k objects from n w/ replacement k^n

k objects from n without replacements: $\binom{n}{k}$

k balls into n boxes, many balls in boxes.

$$\binom{n+k-1}{n-1}$$

$$\binom{n+k-1}{k-1}$$

Basic probability and expectation

Sample space, uniform distribution

Conditional probability & independence

not in the exam