```
Answers:
1.(a)
 gcd(288, 120)
= gcd(288 \mod 120, 120)
= gcd(48, 120)
= gcd(120 \mod 48, 48)
= gcd(48, 24)
= 24
(b)
 lcm(-91, 52)
= |-91 * 52| / gcd(-91,52)
 gcd(-91, 52)
= gcd(39, 52)
= gcd(13, 53)
= 13
hence:
 lcm(-91, 52)
= |-91 * 52| / 13
= 364
(c)
If n = 0, n + 1 = 1, 1|0, 1|1, so gcd(0, 1) = 1
If n \neq 0:
 gcd(n, n+1)
= gcd(n+1-n, n)
= gcd(1, n)
= 1
hence
gcd(n, n + 1) = 1 \text{ for } n \in N
```

$$Pow(\emptyset) = \{\emptyset\}$$

$$Pow(Pow(\emptyset)) = \{\emptyset, \{\emptyset\}\}\$$

$$card(Pow(Pow(\emptyset))) = 2$$

(b)

A n (B
$$\oplus$$
 C)

$$= A \cap ((B \cap C^{c})) \cup (C \cap B^{c}))$$

=
$$(A \cap (B \cap C^{c})) \cup (A \cap (C \cap B^{c}))$$

$$= ((\emptyset \cap B) \cup (A \cap (B \cap C^{c}))) \cup ((\emptyset \cap C) \cup (A \cap (C \cap B^{c})))$$

$$= (((A \cap A^{c}) \cap B) \cup (A \cap (B \cap C^{c}))) \cup (((A \cap A^{c}) \cap C) \cup (A \cap (C \cap B^{c})))$$

$$= (((A \cap B) \cap A^{c}) \cup ((A \cap B) \cap C^{c})) \cup (((A \cap C) \cap A^{c}) \cup ((A \cap C) \cap B^{c}))$$

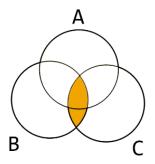
=
$$((A \cap B) \cap (A^{c} \cup C^{c})) \cup ((A \cap C) \cap (A^{c} \cup B^{c}))$$

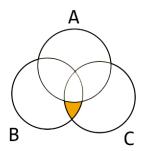
$$= ((A \cap B) \setminus (A \cap C)) \cup ((A \cap C) \setminus (A \cap B))$$

$$= (A \cap B) \oplus (A \cap C)$$

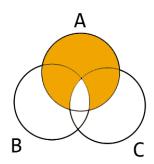
(c)

B ∩ C:

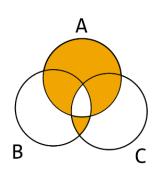




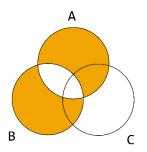
A \ (B n C):



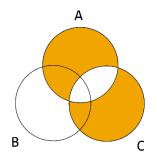
 $A \oplus (B \cap C) = ((B \cap C) \setminus A) \cup (A \setminus (B \cap C))$:



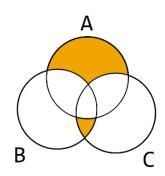
A \oplus B:



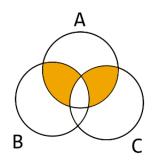
$A \,\oplus\, C\colon$



$(A \oplus B) \cap (A \oplus C)$:



Hence, $(A \oplus (B \cap C)) \setminus ((A \oplus B) \cap (A \oplus C))$:



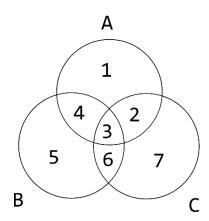
The same Venn diagram above can also be written as:

$$(A \cap B) \oplus (A \cap C)$$

Hence

 $(A \oplus (B \cap C)) \setminus ((A \oplus B) \cap (A \oplus C)) = (A \cap B) \oplus (A \cap C).$ if $(A \cap B) \oplus (A \cap C) \neq \emptyset, (A \oplus (B \cap C)) \neq ((A \oplus B) \cap (A \oplus C)).$

For instance, if A = {1, 2, 3, 4}, B = {3, 4, 5, 6}, C = {2, 3, 6, 7}, then $(A \oplus (B \cap C)) = \{1, 2, 4, 6\}$, and $((A \oplus B) \cap (A \oplus C)) = \{1, 6\}$, Hence, in this case, $(A \oplus (B \cap C)) \neq ((A \oplus B) \cap (A \oplus C))$



3.

- (b) $\Sigma^{\leq 3} = \{\lambda, a, b, aa, ab, ba, bb, aaa, aab, aba, aba, baa, bab, bba, bbb\}$

4.

(a)

f:
$$f(a) = 0$$
, $f(b) = 0$, $f(c) = 0$

f:
$$f(a) = 0$$
, $f(b) = 0$, $f(c) = 1$

f:
$$f(a) = 0$$
, $f(b) = 1$, $f(c) = 0$

$$f: f(a) = 0, f(b) = 1, f(c) = 1$$

f:
$$f(a) = 1$$
, $f(b) = 0$, $f(c) = 0$

f:
$$f(a) = 1$$
, $f(b) = 0$, $f(c) = 1$

f:
$$f(a) = 1$$
, $f(b) = 1$, $f(c) = 0$

f:
$$f(a) = 1$$
, $f(b) = 1$, $f(c) = 1$

(b)

- (i) For each $b \in B$, we can find m members from A to establish functions. Hence number of functions = n^m
- (ii) Relations between A and B: $R = A \times B$. Hence $|R| = |A| \times |B| = m \times n$

```
(c)
Pow(a, b, c) = \{\emptyset, a, b, c, ab, ac, bc, abc\}
|Pow(a, b, c)| = 2^{|\{a, b, c\}|} = 2^3 = 8
Number of f: \{a, b, c\} \rightarrow \{0, 1\} = n^m = 2^3 = 8
|Pow(a, b, c)| = Number of f : \{a, b, c\} \rightarrow \{0, 1\}
To explain this, for f : \{a, b, c\} \rightarrow \{0, 1\}, each possible function
corresponds to some subset A \subseteq Pow(a, b, c), in the way:
              for each element x \in A, f(x) \mapsto 1, and
              for each element y \in (Pow(a, b, c) \setminus A), f(y) \mapsto 0
Hence
       the number of all possible functions = card(Pow(a, b, c)) = 8
5.
(a)
       (i) (abab, baba)
       (v) (\lambda, bbb)
(b)
(R) In the case of (a, a) \in R, for any v \in \Sigma^*, av = av. Hence R is
reflexive.
(S) For (a, b) \in R, for all v \in \Sigma^*,
       (av \in L \land bv \in L) \oplus (av \notin L \land bv \notin L) = true.
Since the Boolean " \Lambda" is commutative:
       (bv \in L \land av \in L) \oplus (bv \notin L \land av \notin L) = true
Hence in the case of (b, a), for all v \in \Sigma^*, R is symmetric.
(T) Let (a, b) \in R and (b, c) \in R, hence for all v \in \Sigma^*, v' \in \Sigma^*:
(av \in L \land bv \in L) \oplus (av \notin L \land bv \notin L) = true, (bv \in L \land cv \in L) \oplus
(bv) \notin L \land cv \notin L = true.
With av \in L, bv \in L, since v and v' are arbitrary, bv \in L, cv \in L:
       Since v and v' are arbitrary, av' \in L, bv' \in L, bv \in L, cv \in L also
       establish, so av \in L, cv \in L, a R c.
With av \notin L, bv \notin L, since v and v' are arbitrary, bv` \notin L and cv` \notin L:
```

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Since v and v' are arbitrary, av' \notin L, bv' \notin L, bv \notin L and cv \notin L also establish, so (av \notin L, cv \notin L).
```

Hence R is transitive.

Concluding from (R), (S), (T), R is an equivalence relation.

```
(c)  X = \{ w \in \Sigma^* \colon 3 | length(w) \} 
 Y = \{ w \in \Sigma^* \colon 3 | (length(w) + 1) \} 
 Z = \{ w \in \Sigma^* \colon 3 | (length(w) + 2) \}
```