

Optimal Transport And WGAN

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Outline

- 1 Introduction To Optimal Transport.
- 2 Minkowski Type Problems
 - Picewise Linear Function And Power Diagram
- 3 WGAN

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Optimal Transport

Monge

Objective: Calculate a transport map $T_{\#}\mu = \nu$ which minimize the transport cost

$$c(T) = \int c(x, T(x)) d\mu(x)$$

Kantorovich

Objective: Calculate a transport plane minimize the transport cost

$$c(\Pi) = \int c(x, y) \Pi(x, y)$$

Optimal Transport: Linear Programming View

The optimal transport is a convex problem, which can be formulated as

$$\begin{aligned} \min \quad & \langle C, F \rangle \\ \text{s.t.} \quad & \sum_i F_{i,j} = q_j \\ & \sum_j F_{i,j} = q_i \end{aligned}$$

is a special case of the linear programming:

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b, x \geq 0 \end{aligned}$$

Linear Programming

Consider the dual problem of the linear programming problem.

Dual:

Primal:

$$\min c^T x$$

$$s.t. Ax = b, x \geq 0$$

and we have the relation:

$$\min b^T y$$

$$s.t. A^T y \leq c$$

$$\begin{aligned} \inf_{Ax=b, x \geq 0} c^T x &= \inf_{x \geq 0} \sup_y c^T x + y^T (b - Ax) \\ &\stackrel{?}{=} \sup_y \inf_{x \geq 0} c^T x - y^T Ax + y^T b \\ &= \sup_{A^T y \leq c} y^T b \end{aligned}$$

Kantorovich Dual

First, let us express the constraint $\gamma \in \Pi(\mu, \nu)$ in the following way.

$$\sup_{\phi, \psi} \int_X \phi d\mu + \int_Y \psi d\nu - \int_{X \times Y} (\phi(x) + \psi(y)) d\gamma$$

so that the primal problem can be expressed by

$$\min_{\gamma \geq 0} \int_{X \times Y} + \sup_{\phi, \psi} \int_X \phi d\mu + \int_Y \psi d\nu - \int_{X \times Y} (\phi(x) + \psi(y)) d\gamma$$

then consider interchanging sup and inf:

$$\sup_{\phi, \psi} \int_X \phi d\mu + \int_Y \psi d\nu + \inf_{\gamma \geq 0} \int_{X \times Y} (c(x, y) - (\phi(x) + \psi(y))) d\gamma$$

Kantorovich Dual

If the central notion in the original Monge-Kantorovich problem is cost, in the dual problem it is price.

Imagine that a company offers to take care of all your transportation problem, buying bread at the bakeries and selling them to the cafes. Let $\psi(x)$ be the price at which a baker of bread at the bakery x and selling them to the cafe y at the price $\phi(y)$

Let us maximize the profit:

$$\sup \left\{ \int_Y \phi(y) d\nu(y) - \int_X \psi(x) d\mu(x) \mid \phi(y) - \psi(x) \leq c(x, y) \right\}$$

Kantorovich Dual

It is easy to proof that

$$\begin{aligned} \sup_{\phi - \psi \leq c} \left\{ \int_Y \phi(y) d\nu(y) - \int_X \psi(x) d\mu(x) \right\} \\ \leq \inf_{\pi \in \Pi(\mu, \nu)} \left\{ \int_{X \times Y} c(x, y) d\pi(x, y) \right\} \end{aligned}$$

If we describe a pair of prices (ϕ, ψ) as tight if

$$\begin{aligned} \phi(y) &= \inf_x (\psi(x) + c(x, y)) \\ \psi(x) &= \sup_y (\phi(y) - c(x, y)) \end{aligned}$$

The following formula can be seen as the definition of c -transform.

c-Cyclical Monotonicity

Definition

Once a function $c : X \times Y \rightarrow \mathbb{R} \cup \{+\infty\}$ is given, we say that a set $\Gamma \subset X \times Y$ is c -cyclically monotone if for every $k \in \mathbb{N}$, every permutation σ and every finite family of points $(x_1, y_1), \dots, (x_k, y_k) \in \Gamma$ we have

$$\sum_{i=1}^k c(x_i, y_i) \leq \sum_{i=1}^k c(x_i, y_{\sigma(i)})$$

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Theorem

If γ is an optimal transport plane for the cost c and c is continuous, then $\text{spt}(\gamma)$ is a CM-set.

c -Cyclical Monotonicity

Theorem

Rockafellar's Theorem

If $\Gamma \neq \emptyset$ is a c -CM set in $X \times Y$ and $c : X \times Y \rightarrow \mathbb{R}$, then there exists a c -concave function $\phi : X \rightarrow \mathbb{R} \cup \{-\infty\}$ such that

$$\Gamma \subset \{(x, y) \in X \times Y : \phi(x) + \phi^c(y) = c(x, y)\}$$

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Proof.

The function ϕ can be defined as

$$\begin{aligned} \phi(x) = \inf \{ & c(x, y_n) - c(x_n, y_n) + c(x_n, y_{n-1}) - c(x_{n-1}, y_{n-1}) + \cdots \\ & + c(x_1, y_0) - c(x_0, y_0) : n \in \mathbb{N}, (x_i, y_i) \in \Gamma \} \end{aligned}$$



Kantorovich Dual

As a result, we can get a theorem as below.

Theorem

If c is C^1 , ϕ is a Kantorovich potential for the cost c in the transport from μ to ν , and (x_0, y_0) belongs to the support of an optimal transport plane γ , then $\nabla_{\phi}(x_0) = \nabla_x c(x_0, y_0)$, provided ϕ is differentiable at x_0 .

Kantorovich Dual

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As an example, if the cost function has the **following form**
 $c(x, y) = h(x - y)$, h is strictly convex. Then there exists an optimal transport plan γ for the cost $c(x, y)$ and is unique of the form $(id, T)_\# \mu$.

Moreover, there exists a Kantorovich potential ϕ and T and the potentials ϕ are linked by

$$T(x) = x - (\nabla h)^{-1}(\nabla \phi(x))$$

Quadratic Case

For the quadratic case $c(x, y) = \frac{1}{2}|x - y|^2$

$$T(x) = x - \nabla \phi(x) = \nabla \left(\frac{x^2}{2} - \phi(x) \right) = \nabla u(x)$$

Theorem

For function $X : \mathbb{R}^n \rightarrow \mathbb{R}$, let us define $u_X = \frac{1}{2}|x|^2 - X(x)$, then we have

$$u_{X^c} = (u_X)^*$$

Proof.

$$u_{X^c}(x) = \sup_y \frac{1}{2}|x|^2 - \frac{1}{2}|x - y|^2 + X(y) = \sup_y x \cdot y - \left(\frac{1}{2}|y|^2 - X(y) \right)$$

Quadratic Case

We go further more on the quadratic case, we only need to minimize the $\int x \cdot y d\gamma$ gives the same result.

We can give the same result easier, actually we have $\phi(x_0) + \phi^*(y_0) = x_0 \cdot y_0$ for $y_0 \in \partial\phi(x_0)$, which means

Theorem

For the quadratic case, there exists unique an optimal transport map T from μ to ν and it is of the form $T = \nabla u$ for a convex function u

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Remark

- The ϕ above is called **Kantorovich Potential**.
- The u here is called **Brenier Potential**.

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Minkowski Problem

Theorem

Suppose Ω is a compact convex polytope with non-empty interior in \mathbb{R}^n are distinct k points and $A_1, A_2, \dots, A_k > 0$ s.t. $\sum_{i=1}^k A_i = \text{vol}(\Omega)$. Then there exists a vector $h = (h_1, \dots, h_k) \in \mathbb{R}^k$, unique up to adding the constant (c, c, \dots, c) , so that the piecewise linear convex function

$$u(x) = \max_{x \in \Omega} \{x \cdot p_i + h_i\}$$

satisfies $\text{vol}(\{x \in \Omega \mid \nabla u(x) = p_i\}) = A_i$

Minkowski Problem

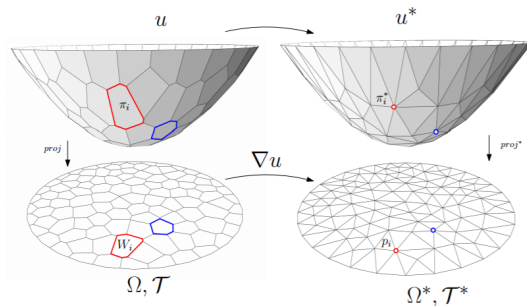


Figure 2: Discrete Optimal Transport Mapping (left to right): map W_i to p_i . Discrete Monge-Ampere equation (right to left): $vol(W_i)$ is the discrete Hessian determinant of p_i .

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Piecewise Linear(PL) Function

Definition

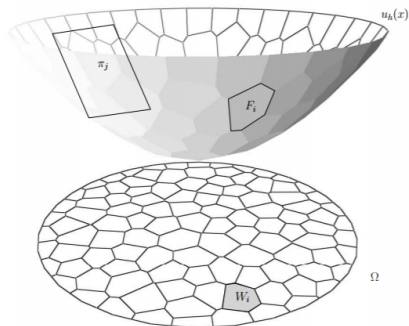
PL Function: For $P = \{p_1, \dots, p_k\}$ and $h = (h_1, \dots, h_k) \in \mathbb{R}^k$, we define the PL convex function $u_h(x)$ to be

$$u(x) = \max\{p_i \cdot x + h_i | i = 1, 2, \dots, k\}$$

The domain $D(u^*)$ of the dual function u^* is the convex hull of P and

$$u^*(y) = \min\left\{-\sum_{i=1}^k t_i h_i \mid t_i \geq 0, \sum_{i=1}^k t_i = 1, \sum_{i=1}^k t_i p_i = y\right\}$$

Piecewise Linear(PL) Function



PL-convex function f defined on a closed convex polyhedron produces a convex subdivision.

Power Diagram

Definition

(Power Distance) Given a point $y_i \in \mathbb{R}^n$ with a power weight ϕ_i the power distance is given by

$$\text{pow}(x, y_i) = |x - y_i|^2 - \phi_i$$

Definition

(Power Diagram) Given weighted points $\{(y_i, \phi_i)\}$, the power diagram is the cell decomposition of \mathbb{R}^n , denote as $V(\phi)$

$$\mathbb{R}^n = \cup_{i=1}^k W_i(\phi), W_i(\phi) = \{x \in \mathbb{R}^n | \text{pow}(x, y_i) \leq \text{pow}(x, y_j), \forall j\}$$

Each cell is a convex polytope

Power Diagram

Now consider a equal construction as let $h_i = \frac{1}{2}(\phi_i - |y_i|^2)$, we construct the convex function

$$u_h(x) = \max_i \{ \langle x, y_i \rangle + h_i \}, \quad W_i(h) = \max_i \{ x \cdot p_i + h_i \}$$

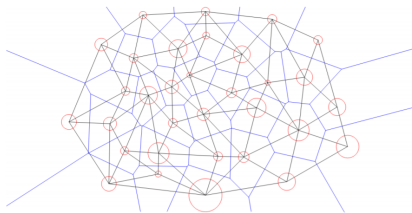


Figure 5: Power diagram (blue) and its dual weighted Delaunay triangulation (black), the power weight ψ_i equal to the square of radius r_i (red circle).

Variation

Proposition

Suppose $\sigma \rightarrow \mathbb{R}$ is continuous defined on a compact convex domain $\Omega \subset \mathbb{R}^n$. If $p_1, \dots, p_k \in \mathbb{R}^n$ are distinct and $h \in \mathbb{R}^k$ so that $\text{vol}(W_i(h) \cap \Omega) > 0$ for all i , then $\omega_i(h) = \int_{W_i(h) \cap \Omega} \sigma(x)$ is a differentiable function in h so that for $j \neq i$ and $W_i(h) \cap \Omega$ and $W_j(h) \cap \Omega$ share a codimension-1 face F ,

$$\frac{\partial \omega_i(h)}{\partial h_j} = -\frac{1}{|p_i - p_j|} \int_F \sigma_F(x) dA$$

where dA is the area form on F and partial derivative is zero otherwise.

Variation

It is easy to observe that $\frac{\partial \omega_i}{\partial h_j} = \frac{\partial \omega_j}{\partial h_i}$, thus we can give our main theorem.

Theorem

Theorem 4.3 (Gu-Luo-Sun-Yau[12]) *Let Ω be a compact convex domain in \mathbb{R}^n , $\{y_1, \dots, y_k\}$ be a set of distinct points in \mathbb{R}^n and μ a probability measure on Ω . Then for any $\nu_1, \dots, \nu_k > 0$ with $\sum_{i=1}^k \nu_i = \mu(\Omega)$, there exists $h = (h_1, \dots, h_k) \in \mathbb{R}^k$, unique up to adding a constant (c, \dots, c) , so that $w_i(h) = \nu_i$, for all i . The vectors h are exactly maximum points of the concave function*

$$E(h) = \sum_{i=1}^k h_i \nu_i - \int_0^h \sum_{i=1}^k w_i(\eta) d\eta_i \quad (19)$$

on the open convex set

$$H = \{h \in \mathbb{R}^k | w_i(h) > 0, \forall i\}.$$

Furthermore, ∇u_h minimizes the quadratic cost

$$\int_{\Omega} |x - T(x)|^2 d\mu(x)$$

among all transport maps $T_{\#}\mu = \nu$, where the Dirac measure $\nu = \sum_{i=1}^k \nu_i \delta_{y_i}$.

Semi-discrete Optimal Mass Transport

For our empirical distribution is defined as the sum of several Dirac measure $\nu = \sum_{j=1}^k \nu_j \delta(y - y_j)$

Define the discrete Kantorovich potential $\phi : Y \rightarrow \mathbb{R}$, $\phi(y_j) = \phi_j$, then

$$\int_Y \phi d\nu = \sum_{j=1}^k \phi_j \nu_j$$

Define the c -transformation of ϕ is given by

$$\phi^c(x) = \min_{1 \leq j \leq k} \{c(x, y_j) - \phi_j\}$$

and each cell is defined as

$$W_i(\phi) = \{x \in X | c(x, y_i) - \phi_i \leq c(x, y_j) - \phi_j, \forall 1 \leq j \leq k\}$$

Brenier's Approach

We only consider the situation that the cost function is the L^2 distance. Here

$$u_h(x) = \max_{i=1}^k \{ \langle x, y_i \rangle + h \}$$

Then

$$W_i(h) = \{x \in X | \nabla u_h(x) = y_i\} \cap \Omega$$

and at the same time

$$\nabla u_h : W_i(h) \rightarrow y_i, i = 1, 2, \dots, k$$

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3 **WGAN**

GAN/WGAN

Objective Function:

$$\min_{u \in U} \max_v \mathbb{E}_{x \sim D_{real}} [\phi(D_v(x))] + \mathbb{E}_{x \sim D_G} [\phi(1 - D_v(x))]$$

- GAN: $\phi = \log$
- WGAN: $\phi = \text{id}$

GAN/WGAN

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Rewrite the WGAN Objective:

$$\min_{u \in U} \max_v \mathbb{E}_{x \sim D_{real}} [D_v(x)] + \mathbb{E}_{x \sim D_G} [1 - D_v(x)]$$

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equal to

$$\min_{u \in U} \max_v \mathbb{E}_{x \sim D_{real}} [D_v(x)] - \mathbb{E}_{x \sim D_G} [D_v(x)]$$

GAN/WGAN

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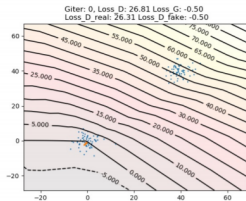
$$\min_{u \in U} \max_v \mathbb{E}_{x \sim D_{real}} [D_v(x)] - \mathbb{E}_{x \sim D_G} [D_v(x)]$$

For if $c(x, y) = |x - y|$, then $\phi^c = -\phi(1\text{-Lip})$, **approximate to W_1 Distance**

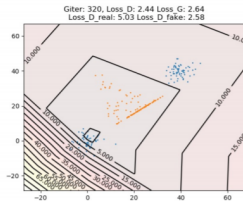
Geometric Generative Model

- Encoding/Decoding process: This step maps the samples between the image space X and the latent space Z by deep neural networks, the encoding map is denoted as $f_\theta : X \rightarrow Z$ and decoding map is $g_\xi : Z \rightarrow X$
- Probability measure transformation process: This step transform a fixed distribution $\xi \in P(Z)$ to any given distribution $\mu \in P(Z)$, the mapping is denoted as $T : Z \rightarrow Z$, $T_{\#}\xi = \mu$. This step can either use conventional deep neural network or use explicit geometric/numerical methods.

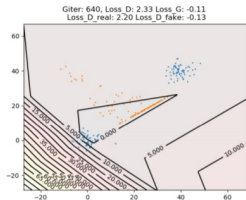
WGAN



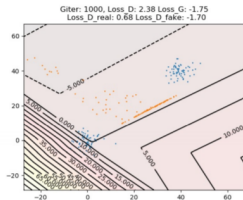
(a) initial stage



(b) after 320 iterations



(c) after 640 iterations



(d) final stage, after 1000 iterations

Geometric OMT

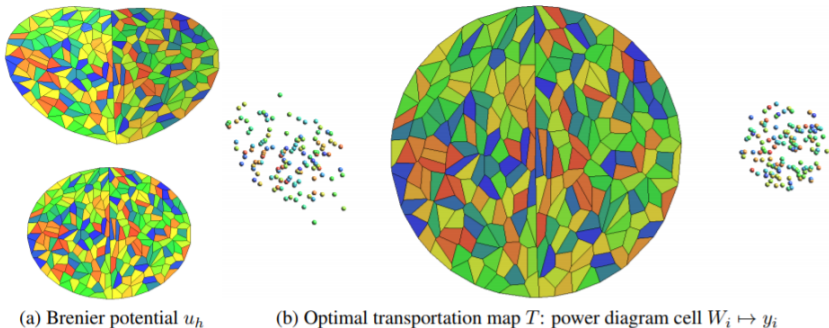


Figure 9: Geometric model learns the Gaussian mixture distribution .

Geometric Method

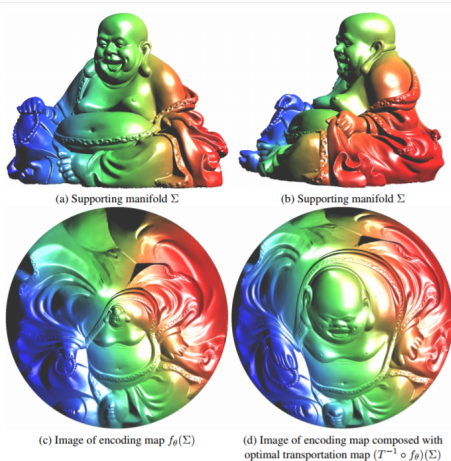


Figure 10: Illustration of geometric generative model.

Geometric Method

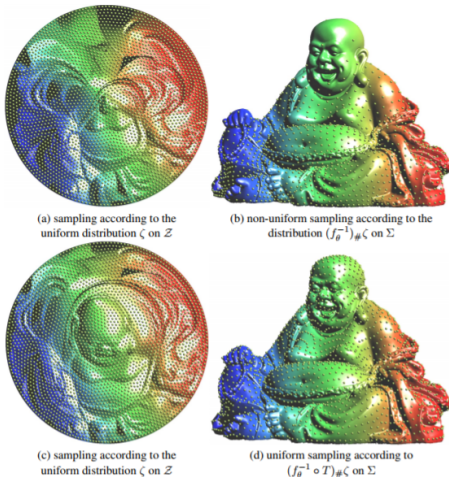


Figure 11: Illustration of geometric generative model.

Conclusion

1. 生成器：最优映射等价于Power胞腔分解，将每个胞腔 W_i 映到 y_i ，
2. 判别器：Wasserstein距离中 $W_c(\mu, \nu)$ 中的 ψ 等于power 权重，
3. 判别器：Wasserstein距离Kantorovich势能 φ 等于power距离，

$$\varphi(x) = \min_i \{\text{pow}(x, y_i)\}$$
4. 生成器：Brenier势能等于Power Diagram的上包络。