

Switching techniques

Traditionally the telephonic system is based on circuit switching; is the main infrastructure for communications (computer) networks => the switching term remains.

Switching techniques used in information transfer are:

- circuit switching**
- message switching**
- packet switching**

Circuit switching

Physical path between communicating parts, achieved using circuit switching –switches (relays)-in the networks nodes.

Three phase communication:

- circuit establishment (setup), establish a (optimum) path between parts; both parts agree communication
- effective data transmission (signal transfer), on this route
- circuit release (disconnection); initiative of one part.

Drawbacks:

- not efficient due to existence of the first phase (it will exist even if there's no data transfer)
- need for covering bandwidth allocation
- important amount of cabling
- no buffers in switches for transmission equalization

Today use of digital **PBX** (Private Branch Exchange)

First circuit-switching: space-division switching (separated signal paths – divided in space): crossbar matrix of I/O full duplex lines

An improvement: multiple-stage switches

Today all telephony: digital time-division techniques (synchronous TDM)

Signaling in digital telephony:

- inchannel

 - in-band: signals using the same band as the voice channel (as payload)

 - out-band: (voice signals do not use whole 4kHz bandwidth)

- common channel – a common signal channel for a number of voice channels

Signaling may use the same (or not) path as the payload (associated/nonassociated modes)

What's signaling?

Signal = control Examples:

- connection setup request = off-hook signal from telephone to switch

- connection setup acknowledge = dial tone

- destination address = pulse or tone dialing

- destination busy = busy tone

- destination available = ringing tone

Other signaling functions: transmission of: dialed number between switches, information about a call not completed, about billing, diagnose and failure isolation

Message switching

Data transfer using **messages** (independent data units, with diff. lengths but similar structures). Types: control and data (embedding control)

Need for addressing (source & destination of message)

Communications nodes are not physical switches, but computing systems (with memory and processing units).

Philosophy is: message *store & forward* .

Not more dedicated communications path; established in an optimum way (cost, network status) by nodes (using routing tables).

Advantages:

- improvement in efficiency (path multiplexing)
- introduces message priority
- equilibrated transmissions.

Drawbacks:

- messages are too long, memory waste and difficult error recovery

Packet switching

Combines the advantages of previous methods. The **packet** has similar message structure but a lower length, up to 1000octets.

Two methods:

- use of **datagrams** (close to message switching)-more speedy and flexible method

- use or not of transmission acknowledgments (ACK)

- use of **virtual circuits** (close to circuit switching)-use of the three phases (connection request, data transfer, disconnect) for a logical connection activation; use of special control packets for that. Also embedding of control information (piggybacking).

A logical connection may be implemented with more different physical connections.

Routing in packet-switching networks

Circuit switching vs. Packet switching

Most of WANs based on circuit or packet switching

Circuit switching designed for voice

- Resources dedicated to a particular call

- Much of the time a data connection is idle

- Data rate is fixed

- Both ends must operate at the same rate

Packet switching - Basic Operation

Data transmitted in small packets

- Typically 1000 octets

- Longer messages split into series of packets

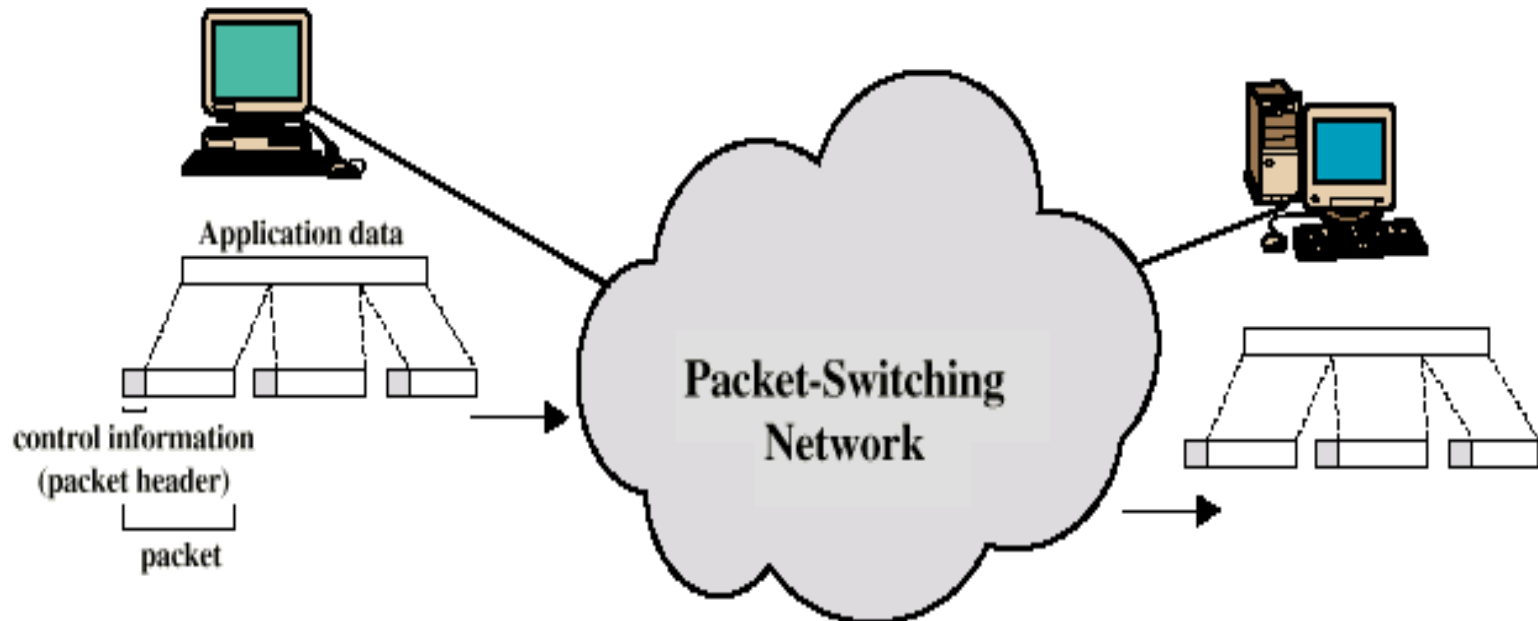
- Each packet contains a portion of user data plus some control info

Use of **Control info**

Routing (addressing) info

Packets are received, stored briefly (buffered) and past on to the next node

Store and forward



Advantages of packet switching

Line efficiency

Single node to node link can be shared by many packets over time

Packets queued and transmitted as fast as possible

Data rate conversion

Each station connects to the local node at its own speed

Nodes buffer data if required to equalize rates

Packets are accepted even when network is busy

Delivery may slow down

Priorities can be used

Packet Switching Technique

Station breaks long message into packets

Packets sent one at a time to the network

Packets handled in two ways: **Datagram** or **Virtual circuit**

Virtual Circuits v Datagram

Virtual circuits

Network can provide sequencing and error control

Packets are forwarded more quickly

No routing decisions to make

Less reliable

Loss of a node loses all circuits through that node

Datagram

No call setup phase

Better if few packets

More flexible

Routing can be used to avoid congested parts of the network

Use of variant with acknowledgements

| Issue | Datagram subnet | Virtual-circuit subnet |
|---------------------------|--------------------------------------------------------------|------------------------------------------------------------------|
| Circuit setup | Not needed | Required |
| Addressing | Each packet contains the full source and destination address | Each packet contains a short VC number |
| State information | Routers do not hold state information about connections | Each VC requires router table space per connection |
| Routing | Each packet is routed independently | Route chosen when VC is set up; all packets follow it |
| Effect of router failures | None, except for packets lost during the crash | All VCs that passed through the failed router are terminated |
| Quality of service | Difficult | Easy if enough resources can be allocated in advance for each VC |
| Congestion control | Difficult | Easy if enough resources can be allocated in advance for each VC |

Comparison of datagrams & virtual-circuit subnets

Network's External and Internal Operation

Interface between station and network node

Connection oriented

Station requests logical connection (virtual circuit)

All packets identified as belonging to that connection & sequentially numbered

Network delivers packets in sequence

External virtual circuit service

e.g. X.25

Different from internal virtual circuit operation

Connectionless

Packets handled independently

External datagram service

Different from internal datagram operation

Combinations

External virtual circuit, internal virtual circuit

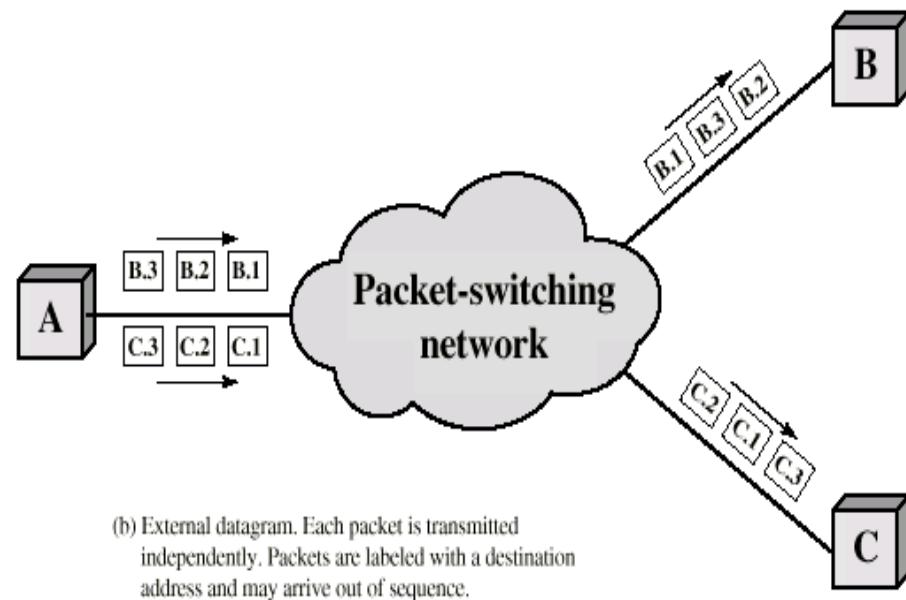
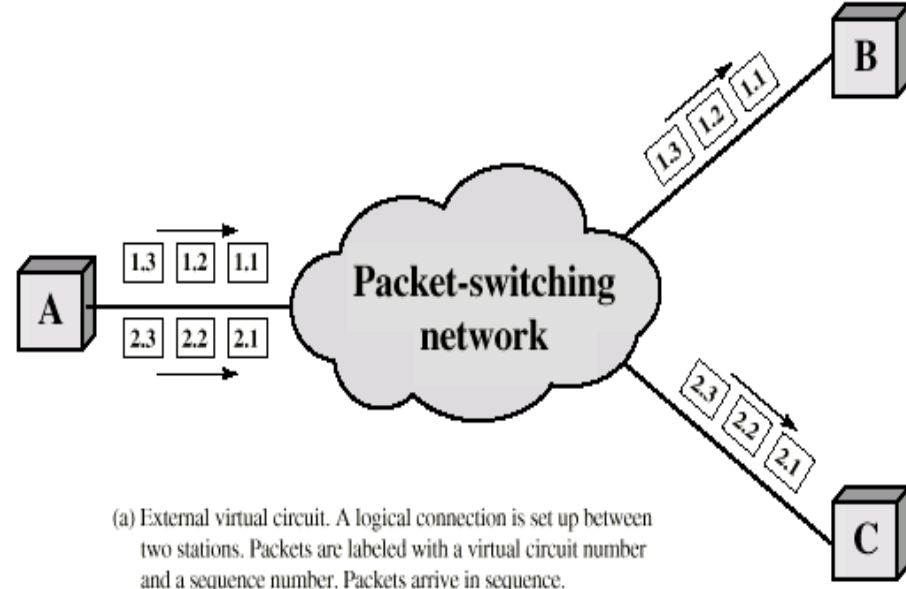
Dedicated route through network

External virtual circuit, internal datagram

Network handles each packet separately

Different packets for the same external virtual circuit may take different internal routes

Network buffers at destination node for re-ordering



External datagram, internal datagram

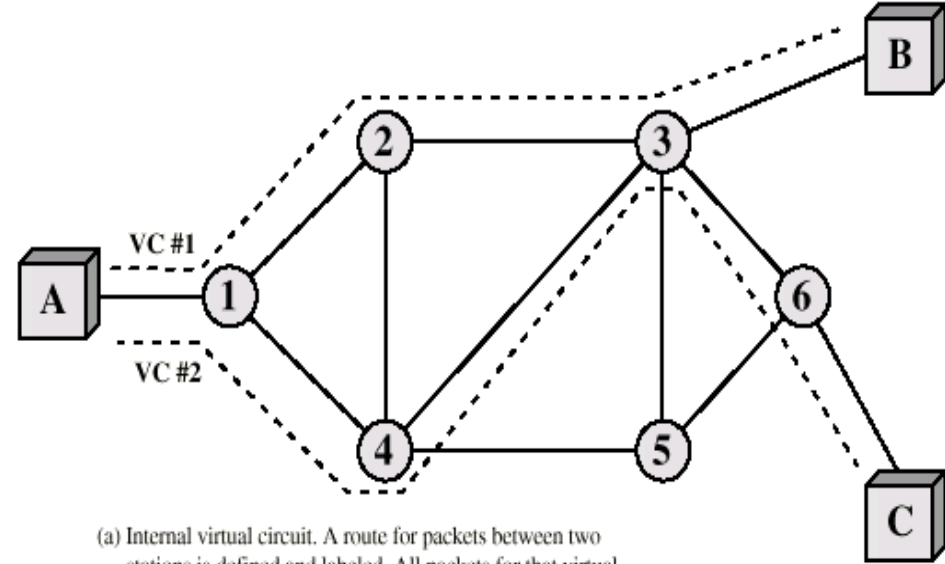
Packets treated independently by both network and user

External datagram, internal virtual circuit

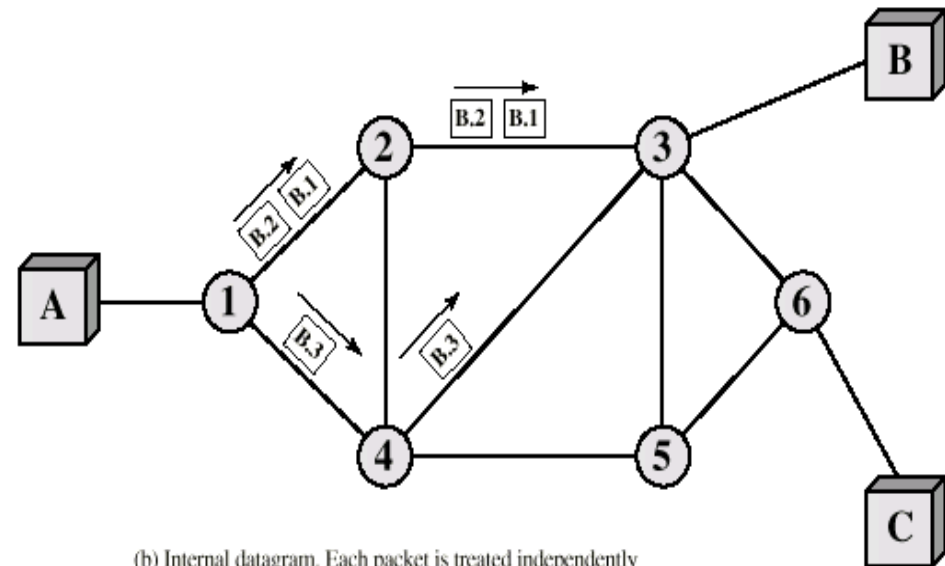
External user does not see any connections

External user sends one packet at a time

Network sets up logical connections



(a) Internal virtual circuit. A route for packets between two stations is defined and labeled. All packets for that virtual circuit follow the same route and arrive in the same sequence.



(b) Internal datagram. Each packet is treated independently by the network. Packets are labeled with a destination address and may arrive at the destination node out of sequence.

Routing

Complex, crucial aspect of packet switched networks

Characteristics required

Correctness

Simplicity

Robustness

Stability

Fairness

Optimality

Efficiency

Performance Criteria

Used for selection of route

Minimum hop

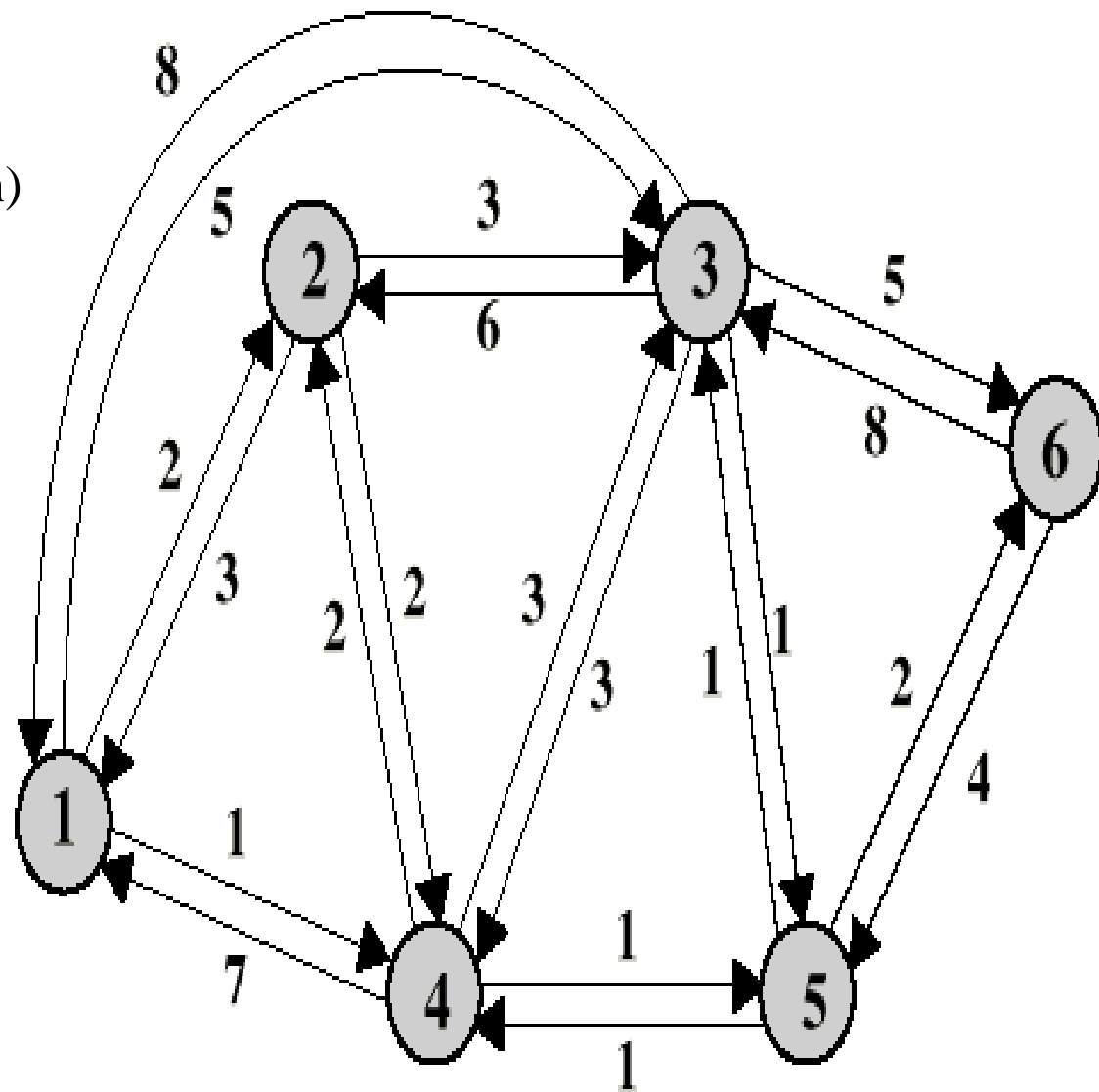
Least cost algorithms (shortest path)

Dijkstra's Algorithm

Implemented in link state packet
routing algorithms

Bellman-Ford algorithm

Used by distance vector based



Routing: Decision Time and Place

Time

On packet or virtual circuit basis

Place

Distributed routing

Made by each node

Centralized routing

Source-based routing

Network Information Source and Update Timing

Routing decisions usually based on knowledge of network (not always)

Distributed routing

Nodes use local knowledge

May collect info from adjacent nodes

May collect info from all nodes on a potential route

Central routing

Collect info from all nodes

Update timing

When is network info held by nodes updated

Fixed - never updated

Adaptive - regular updates

Routing Strategies

Fixed

Flooding

Random

Adaptive

Fixed Routing

Single permanent route for each source to destination pair

Determine routes using a *least cost algorithm*

Route fixed, at least until a change in network topology

| | | CENTRAL ROUTING DIRECTORY | | | | | |
|---------|---|---------------------------|---|---|---|---|---|
| | | From Node | | | | | |
| | | 1 | 2 | 3 | 4 | 5 | 6 |
| To Node | 1 | — | 1 | 5 | 2 | 4 | 5 |
| | 2 | 2 | — | 5 | 2 | 4 | 5 |
| | 3 | 4 | 3 | — | 5 | 3 | 5 |
| | 4 | 4 | 4 | 5 | — | 4 | 5 |
| | 5 | 4 | 4 | 5 | 5 | — | 5 |
| | 6 | 4 | 4 | 5 | 5 | 6 | — |

| Node 1 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 2 | 2 |
| 3 | 4 |
| 4 | 4 |
| 5 | 4 |
| 6 | 4 |

| Node 2 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 1 | 1 |
| 3 | 3 |
| 4 | 4 |
| 5 | 4 |
| 6 | 4 |

| Node 3 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 1 | 5 |
| 2 | 5 |
| 4 | 5 |
| 5 | 5 |
| 6 | 5 |

| Node 4 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 1 | 2 |
| 2 | 2 |
| 3 | 5 |
| 5 | 5 |
| 6 | 5 |

| Node 5 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 1 | 4 |
| 2 | 4 |
| 3 | 3 |
| 4 | 4 |
| 6 | 6 |

| Node 6 Directory | |
|------------------|-----------|
| Destination | Next Node |
| 1 | 5 |
| 2 | 5 |
| 3 | 5 |
| 4 | 5 |
| 5 | 5 |

Flooding

No network info required

Packet sent by node to every neighbor

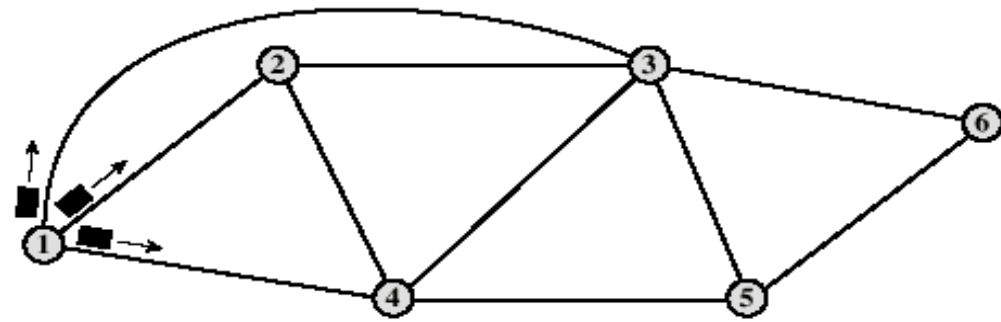
Incoming packets retransmitted on every link except incoming link

Eventually a number of copies will arrive at destination

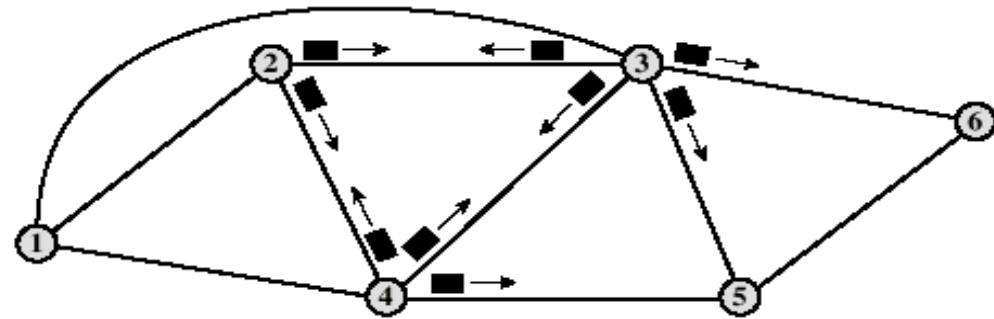
Each packet is uniquely numbered so duplicates can be discarded

Nodes can remember packets already forwarded to keep network load in bounds

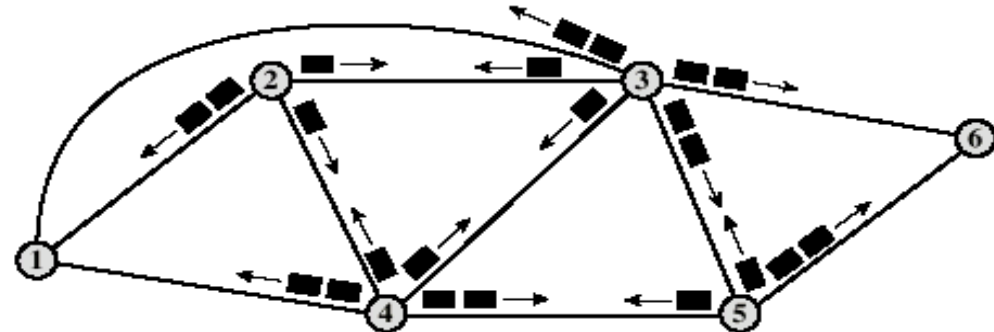
Can include a hop count in packets



(a) First hop



(b) Second hop



(c) Third hop

Properties of Flooding

All possible routes are tried

Very robust

At least one packet will have taken minimum hop count route

Can be used to set up virtual circuit

All nodes are visited

Useful to distribute information (e.g. routing)

Random Routing

Node selects one outgoing path for retransmission of incoming packet

Selection can be random or round robin

Can select outgoing path based on probability calculation

No network info needed

Route is typically not least cost nor minimum hop

Adaptive Routing

Used by almost all packet switching networks

Routing decisions change as conditions on the network change

- Failure

- Congestion

Requires info about network

Decisions more complex

Tradeoff between quality of network info and overhead

Advantages

Improved performance

Aid congestion control

Complex system

- May not realize theoretical benefits

Classification

Based on information sources for network state

Local (isolated)

Route to outgoing link with shortest queue

Can include bias for each destination

Rarely used - do not make use of easily available info

Adjacent nodes – select information based on the neighbour's experience (network delays or outages)

All nodes – used for source based routing