

The University of Melbourne

Practice Exam Paper

School of Computing and Information Systems

COMP90038 Algorithms and Complexity

Leo's solution

Question 1

A. Give the names of two *stable* sorting algorithms, together with their worst-case time complexities. Write the names and complexities in the box:

1. Insertion sort, $\Theta(n^2)$
2. Merge sort, $\Theta(n \log n)$

B. Give the names of two *unstable* sorting algorithms, together with their worst-case time complexities. Write the names and complexities in the box:

1. Selection sort, $\Theta(n^2)$
2. Quick sort, $\Theta(n^2)$

Question 2

(4 marks)

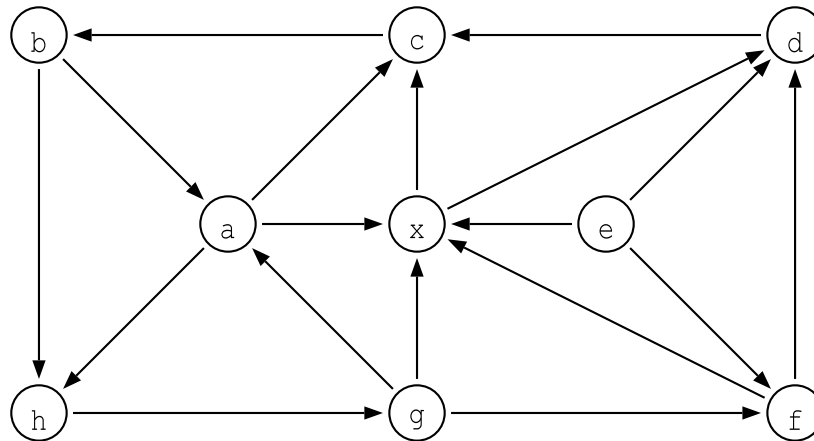
We are given an array A holding n integers, for some large n . **The array is sorted**, and the values in A range from -2147483648 to 2147483647, evenly distributed. Give Θ expressions for the following tasks:

- A. Running the insertion sort algorithm on the array A : $\Theta(n)$
- B. Running the selection sort algorithm on the array A : $\Theta(n^2)$
- C. Performing binary search for integer k which is not in A : $\Theta(\log n)$
- D. Performing interpolation search for integer k not in A : $\Theta(1)$

Question 3

For the directed graph below, list the order in which the nine nodes are visited during a depth-first (DFS) traversal, as well as the order in which they are visited during a breadthfirst (BFS) traversal. As always, assume that any **ties are resolved** by taking nodes in alphabetical order. Write the answers in the boxes given.

(4 marks)



DFS sequence:

a c b h g f d x e

BFS sequence:

a c h x b g d f e

Question 4

(4 marks)

Given the pattern A T G A and the text

T C A T C A T C C A T G C A C A A T G A C T T T

how many character comparisons will Horspool's algorithm make before locating the pattern in the text? Write the number in the box:

$$1+2+1+1+2+4(\text{successful searching}) = 11$$

Question 5

Assume the array A holds the keys 77, 64, 15, 43, 28, 91, 80, 32, 56 in index positions 1 to 9. Show the heap that results after application of the linear-time bottom-up heap construction algorithm. You may show the heap as a tree or as an array.

[91 64 80 56 28 15 77 32 43]

(4 marks)

Question 6

The functions A – D are defined recursively as follows (all divisions round down, to the closest integer):

$$A(n) = 2 A(n/3) + 2, \quad \text{with } A(1) = 1$$

$$B(n) = B(n/2) + n/2, \quad \text{with } B(1) = 1$$

$$C(n) = 512 C(n/8) + 4n^2, \quad \text{with } C(1) = 4$$

$$D(n) = 4 D(n/2) + n^2, \quad \text{with } D(1) = 2$$

In the following table, for each of the four functions, tick the most precise correct statement about the function's rate of growth:

	$O(n)$	$\Theta(n)$	$O(n \log n)$	$\Theta(n^2)$	$O(n^2 \log n)$	$\Theta(n^2 \sqrt{n})$	$O(n^3)$
A	✓						
B		✓					
C							✓
D					✓		

Question 7

For each of **A–D** below, answer yes or no, and, in each case, briefly explain your reasoning (just a justification of your answer, rather than detailed calculations). A yes/no answer that is not justified will not attract marks, even if correct.

Question	Answer/explanation
A. Is $\sqrt{n} \in \Omega(n)$?	No, because $n/(\text{root } n) = \text{root } n$

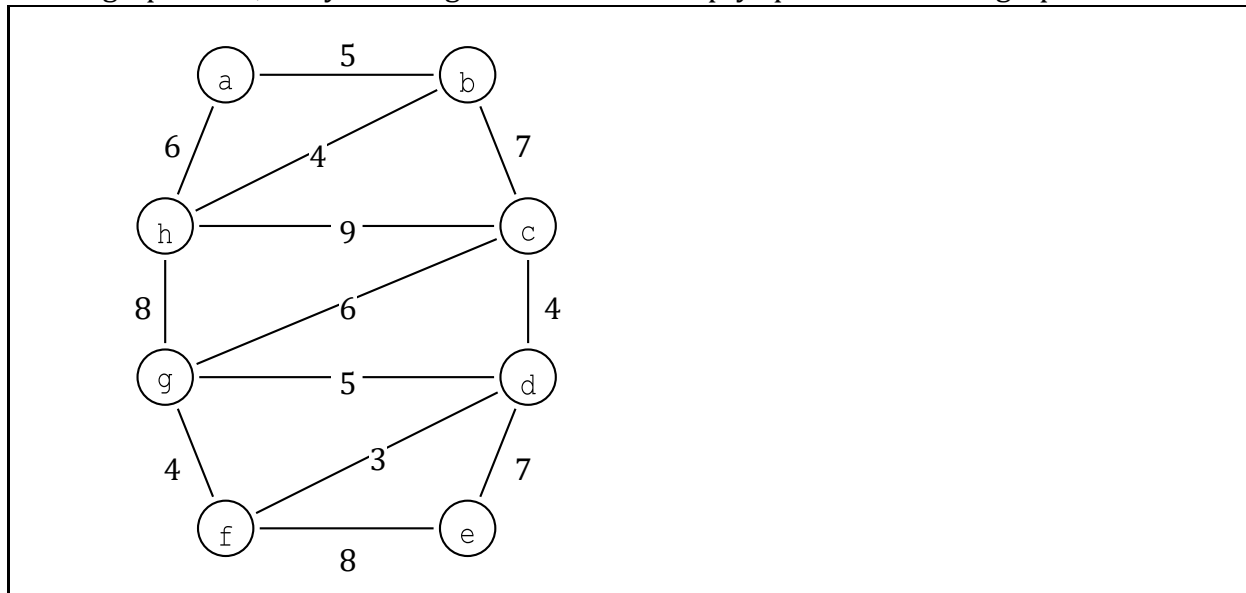
(4 marks)

<p>B.</p> <p>Is $n^{(2\log n)} \in O(n^{(\log n)})$?</p>	<p>No, because</p> $\frac{n^{(2\log n)}}{n^{(\log n)}} = n^{(\log n)}$
<p>C.</p> <p>Is $\Theta(\log(n^{2\log n})) = \Theta(\log(n^{\log n}))$?</p>	<p>Yes, because</p> $\frac{\log(n^{2\log n})}{\log(n^{\log n})} = \frac{2\log n}{\log n} = 2$
<p>D.</p> <p>Is $\Theta(\log(2^n)) = \Theta(\log(3^n))$?</p>	<p>Yes, because</p> $\frac{\log(2^n)}{\log(3^n)} = \frac{(\log(3^n) * \log(2)(3))}{\log(3^n)} = \log(2)(3)$

Question 8

(6 marks)

A. The box below contains a weighted undirected graph with eight nodes. Give a minimum spanning tree for the graph. You may do that either by outlining a minimum spanning tree on the graph itself, or by drawing the tree in the empty space next to the graph.



These edges were chosen: a-b, b-h, b-c, c-d, d-f, f-g, d-e

B. Given a weighted graph $G = \langle V, E \rangle$, a subgraph $\langle V, E' \rangle$ (that is, $E' \subseteq E$) which is a tree with *minimum* weight is a *maximum spanning tree* for G .

We want a transformation of the graph G so that we can run Prim's algorithm on the transformed graph G' , and the algorithm will find a maximum spanning tree for G .

Describe such a (systematic) transformation from G to G' .

(I think the red word should be "**maximum**", otherwise this question is unreasonable.)

For each edge in G , we change the weight from (x) to $(-x)$. For example, the weight of "a to b" is 5 originally, now we let it become -5.

Then we can run Prim's algorithm to find the maximum tree.

Question 9

(6 marks)

Consider the function F below. The function takes as input an integer array A , and the size n of A . The array indices run from 1 to n . The division used is integer division, that is, it rounds down to the closest smaller (or equal) integer value.

In the box, give a Θ expression for the function's time complexity.

function $F(A[\cdot], n)$

$s \leftarrow 0$

$m \leftarrow n$

while $m > 0$ **do**

for $i \leftarrow 1$ **to** m **do**

$s \leftarrow s + A[i]$

$m \leftarrow m/2$

$$n + (n/2) + (n/4) + \dots + 1 = 2n - 1 = \Theta(n)$$

Question 10

Using pseudo-code, give an algorithm for deleting the smallest element of a binary search tree (a BST). Assume a non-empty binary tree T has attributes $T.left$, $T.right$, and $T.root$ which denote T 's left sub-tree, right sub-tree, and the key of T 's root node, respectively. You can use these tests if they seem useful: $\text{isLeaf}(T)$ tests whether the binary tree T is a leaf, and $\text{isEmpty}(T)$ tests whether it is empty.

Function $\text{deleteSmallest}(T)$

if $\text{isEmpty}(T.left)$

$T = T.right$ *//Delete T*

else

$\text{deleteSmallest}(T.left)$

(10 marks)

Question 11

Consider an array A of n distinct integers (that is, all elements are different). It is known that A was originally sorted in ascending order, but A was then right-rotated r places, where $0 < r < n$. In other words, the last r elements were moved from the end of the array to the beginning, with all other elements being pushed r positions to the right. For example, for $n = 7$ and $r = 3$, the result may look like this:

[43,46,58,12,20,29,34]

For $r = 5$, the result, based on the same original array, would be

[29,34,43,46,58,12,20]

You know that the given $A[0..n-1]$ has this particular form, that is, for some r , the sequence $A[r], \dots, A[n-1], A[0], \dots, A[r-1]$ is in ascending order, but you do not know what r is. Design an algorithm to find the largest integer in A . Full marks are given for an algorithm that works in time $O(\log n)$; half marks are given for a solution that is correct, but less efficient.

```
Function findMax(A[])
    return find(A, 0, A.length-1)
```

```
Function find(A[], start, end)
    mid = (start+end)/2
    if A[mid]>A[mid+1]
        return A[mid]
    if A[mid]>A[0]
        return find(A, mid+1, A.length-1)
    return find(A, 0, mid-1)
```

(10 marks)

Question 12

Two programmers face the following problem. Given an array containing n random integers in random order, find the largest integer. The integers are placed in cells $A[1] \dots A[n]$.

Programmer X has come up with the code shown below, on the left. (In the programming language used, arrays are indexed from 0, but X's method does not use $A[0]$.)

```
function X( $A[\cdot], n$ )
   $max \leftarrow A[1]$ 
   $i \leftarrow 2$ 
  while  $i \leq n$  do
    if  $A[i] > max$  then
       $max \leftarrow A[i]$ 
     $i \leftarrow i + 1$ 
  return  $max$ 
```

```
function Y( $A[\cdot], n$ )
   $i \leftarrow n$ 
  while  $i > 0$  do
     $A[0] \leftarrow A[i]$ 
     $i \leftarrow i - 1$ 
    while  $A[0] > A[i]$  do
       $i \leftarrow i - 1$ 
  return  $A[0]$ 
```

Programmer Y has solved the same problem differently, as shown above on the right.

Compare the two solutions using three criteria: Correctness, time complexity class, and the number of comparisons performed. Write your analysis in the box:

	Left one	Right one
Correctness	Correct	Correct
Time complexity	$O(n)$	$O(n)$
Number of comparisons	$(n-1) + n$	$n + x$ (where x may be 2 to $n+1$)

*The green part is for “if $A[i] > max$ ” (left) or “while $A[0] > A[i]$ ” (right).

*The red part is the number of comparisons for “while $i \leq n$ ” (left) or “while $i > 0$ ” (right).