# Bounded Model Checking Of C Programs: CBMC Tool Overview

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## Model Checking

Given a model M of a system and a property P, check :

• if  $M \models P(M \text{ models } P)$ , P holds in M

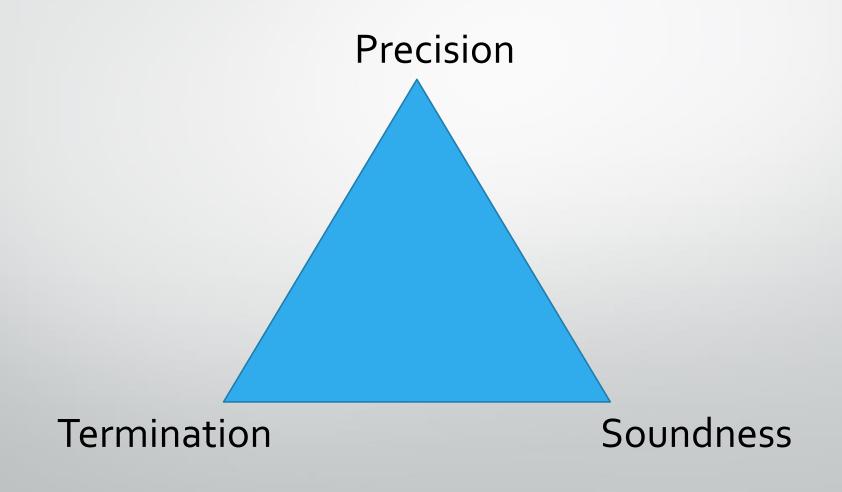
i.e. the system functions according to P.

• if  $M \not\models P(M \text{doesn't model } P)$ , P doesn't hold in M,

and a counterexample is produced, i.e. an execution of the system that does not satisfy P

# Why Bounded Model Checking

In general, there is a compromise to be made between the precision of the analysis and its decidability



## **Bounded Model Checking**

Given a model M of a system, a property P and a bound  $k \, (>\! 0)$ 

- ${}^{\bullet}$  Encode all executions of M of length k into a formula  $M_{\boldsymbol{k}}$
- $\bullet$  Encode all executions of M of length k that violate P into  $\neg P_k$
- If (  $M_k \wedge \neg P_k$ ) is un-satisfiable then P holds in M of length k
- if  $(M_k \land \neg P_k)$  is satisfiable then P doesn't hold in M of length k , and a <code>counterexample</code> is <code>produced</code>

# Simplified Safety Properties

- Array bounds (Buffer Overflows)
- Division by zero
- Pointer checks (i.e., NULL pointer dereference)
- Arithmetic overflow
- Custom assertions (i.e., assert (i > j) )

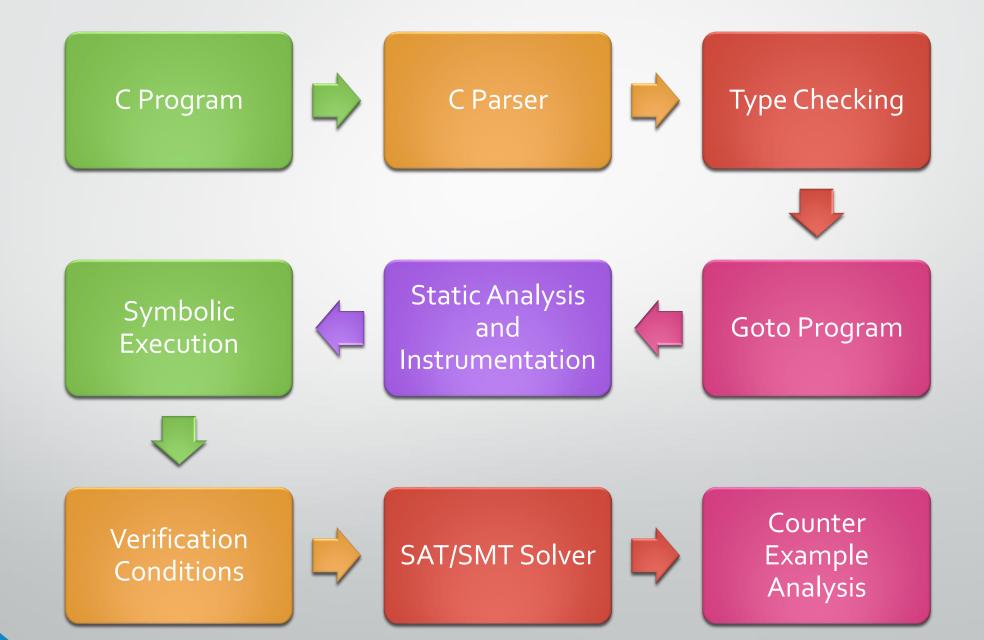
# C Bounded Model Checker **Assertion Proved** Annotated C CBMC Out of Resources Program Counter Example **Unrolling Bound CBMC – Developed and Maintained by Dr Daniel Kröning** www.cprover.org/cbmc/

#### Use Cases

- Verification of operating systems
- Verification of Linux device drivers
- Equivalence Checking (Code Generators / Translation Validation / FBD -> C)
- Runtime Verification, Reachability Verification
- Model Based testing
- WCET analysis

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#### **CBMC** Architecture



# **CBMC Simplifications and Transformations**

- Removal of side effects
- Transforming all explicit loops into while loops
- All non-linear control flow is replaced by guarded goto.
- Generates one CBMC goto Program per function
- Performs light weight static analysis to resolve function pointers.
- Implicit assertions are introduced.

# Goto Program

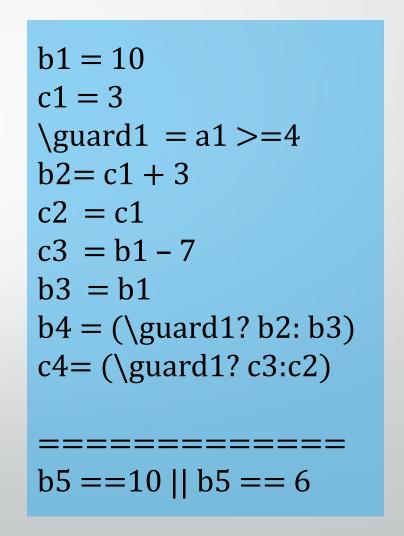
```
int main() {
  int a,SIZE;
  a =2;
  SIZE = 8;
  int i, sn=0;
  for(i=SIZE; i>=0; i--)
  {
    if (i<6)
      sn = sn + a;
  }
  assert(sn<=SIZE*a);
}</pre>
```



```
main /* main */
    signed int a;
    signed int SIZE;
    a = 2;
    SIZE = 8;
    signed int i;
    signed int sn;
    sn = 0;
    i = SIZE;
  1: IF !(i >= 0) THEN GOTO 3
    IF !(i < 6) THEN GOTO 2
    sn = sn + a;
  2: i = i - 1;
    GOTO 1
  3: ASSERT sn <= SIZE * a
    IF !(sn <= SIZE * a) THEN GOTO 4
  4: dead sn;
    dead i;
    dead SIZE;
    dead a;
    main#return_value = NONDET(signed int);
    END_FUNCTION
```

#### Generation of Verification Conditions

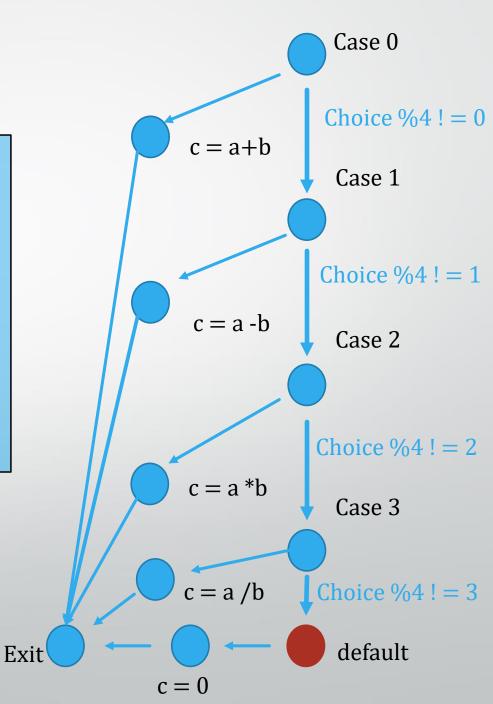
```
int a;
int b = 10, c = 3;
if(a>3)
       b = c + 3;
else
       c = b - 7;
assert(b == 10 || b
==6)
```



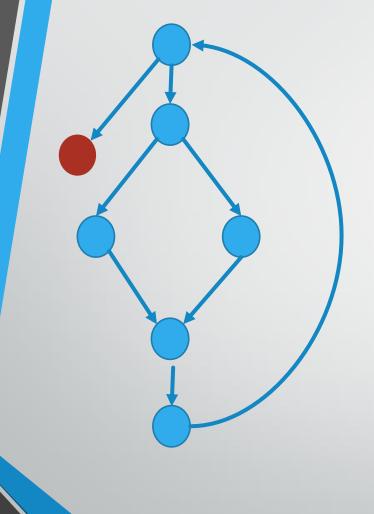
```
int choice;
if(choice > 0)
switch (choice%4)
case 0:
         c = a+b;
         break;
case 1:
         c = a-b;
         break;
case 2:
         c = a*b;
         break;
case 3:
         c = a/b;
         break;
default:
         assert(0);
         c = 0;
         break; }
```

#### Paths

```
choice > 0
choice \% 4 \neq 0
choice \% 4 \neq 1
choice \% 4 \neq 2
choice \% 4 \neq 3
```



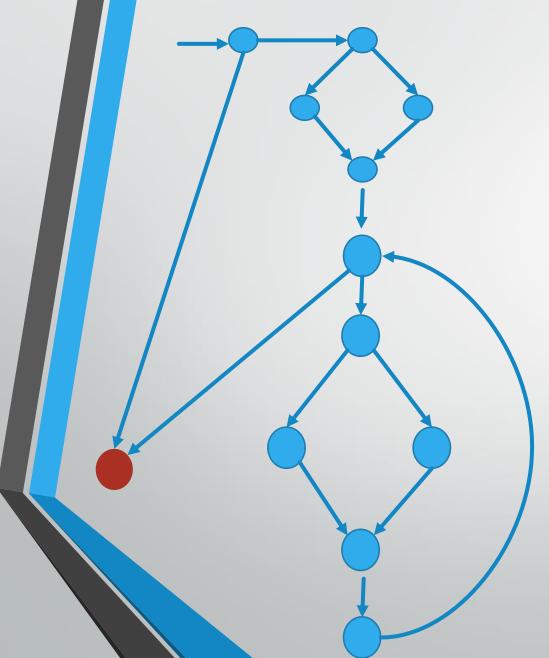
# Loops



```
while( some condition holds)
{
  do something here
}
```

```
int a, b;
while(a < b)
{
    a = a+1;
}</pre>
```

# Loop Unrolling



```
if( some condition holds)
{
  do something here
  while( same condition holds)
      do same thing here
}
...
```

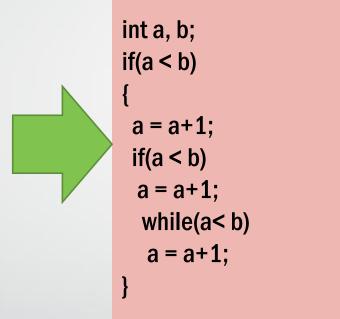
#### **Unwinding Assertions**

```
if( some condition holds)
{
  do something here
  if( same condition holds)
  {
   do same thing here
   while( same condition holds)
      do same thing here
}
```

```
if( some condition holds)
{
  do something here
  if( same condition holds)
  {
    do same thing here
    while( same condition holds)
       do same thing here
    assume(!cond)
}
```

```
if( some condition holds)
{
  do something here
  if( same condition holds)
  {
    do same thing here
    while( same condition holds)
       do same thing here
    assert (!cond)
}
```

```
int a, b;
if(a < b)
{
    a = a+1;
    while(a < b)
        a = a+1;
}</pre>
```





```
int a, b;
if(a < b)
a = a+1;
 if(a < b)
 a = a+1;
   if(a<b)
    a = a+1;
_assume(a \ge b)
```

#### Assumptions and Assertions

- The \_\_CPROVER\_assume statement restricts the program traces that are considered and allows assume-guarantee reasoning.
- The \_\_CPROVER\_assert statement aborts the program successfully if the condition evaluates to false.

```
int onetoten ()
{
int value=nondetint();
__CPROVER_assume ( value>=1 &&
value <=10);
return value ;
}</pre>
```

#### **CBMC Checks**

- Bounds check
- Div by Zero checks
- Pointer Checks
- Memory Leak checks

- Unsigned Overflows
- Float overflows
- Nans
- Signed Overflow Checks

# Supported Language Features: Arithmetic Operators

- Supports all ANSI-C Boolean operators on scalar variables
- Support all integer and floating arithmetic operators on scalar variables
- Full support for arithmetic type casts
- Different rounding modes are currently not supported
- CBMC allows to model user-input by means of non-deterministic choice functions.

# Supported Features: Arrays

- Arrays are encoded with the WITH and [] operators.
- Can be mapped to store and select Operators in theory of arrays

$$i2 = 0$$
  
 $x2 = x1$  WITH  $[0_1 := j1]$   
 $p2 = x2[i1_1] + k1$   
 $q2 = j1 + k1$ 

#### Supported Features: Pointers

Exact points-to analysis is performed.

if(d ==0)  

$$x = &a$$
  
else  
 $x = &b$   
\* $x = c;$ 

\guard1 = (d1 == 0)  

$$x2 = &a1$$
  
 $x3 = x1$   
 $x4 = &b1$   
 $x5 = \guard1? &a1 : &b1$   
 $b2 = (x5 == &b1)? c1: b1$   
 $a2 = (x5 == &b1)? a1: c1$ 

#### Supported Features: Structures

Structures are handled as aggregate data structures like arrays

```
p = &y
if(choice == 1)
       p->a[1]=1;
       p->b = 'c';
else
       p->a[1]=1;
       p->b = 'c';
```

```
p2 = &y1
\gaurd1 = (choice1 == 1)
y2 = y1 WITH [.a:=y1.a WITH [1] = 1]]
y3 = y2 WITH [.b = 'c']
y4 = y1
y5 = y4 WITH [.a = y4.a WITH [1] = 2]]
y6 = y5 WITH [.b = 'c']
y7 = \guard1 ? y3 : y6
```

#### Supported Features: Unions

- Unions are allowed in CBMC, but access to elements of the unions is not allowed across fields.
- Access to a field is permitted if only it was the field last updated.

```
union myunion
{
  char a[2];
  int b;
};

union myunion mu;
  mu.a[0]= 0;
  mu.a[1] = 3;
  k = mu.b;
  assert(k ==3);
```

# Trace for main.assertion.1: $mu=\{ .a=\{ 0, 3 \} \} (\{ 00000000, 00000011 \})$ State 21 file /home/prateek/workspace/workshop/examples/union.c line 17 function main thread 0 State 22 file /home/prateek/workspace/workshop/examples/union.c line 18 function main thread 0 k=768 (000000000000000000001100000000) Violated property: file /home/prateek/workspace/workshop/examples/union.c line 19 function main assertion k == 3k == 3

```
mu=\{ .a=\{ 0, 3 \} \} (\{ 00000000, 00000011 \})
```

State 21 file /home/prateek/workspace/workshop/examples/union.c line 17 function main thread 0

-----

State 22 file /home/prateek/workspace/workshop/examples/union.c line 18 function main thread 0

-----

k=536871680 (001000000000000000001100000000)

Violated property:

file /home/prateek/workspace/workshop/examples/union.c line 19 function main assertion k ==768

k == 768

#### **Function Calls**

- Support functions by in-lining.
- Preserves the locality of the parameters and the non-static local variables by renaming.
- Supports Recursion by finite unwinding

### Other Supported Features

Test case generation:

CBMC can be used to automatically generate test cases following a certain code coverage criterion.

- 1. MC/DC
- 2. Decision Coverage
- 3. Branch
- 4. Path

#### Test Case Generation contd...

```
int foo(int a,int b, int c)
int d;
if (a > 12 && b < 45)
    if (c > 4)
          d = 1;
    else
          d = 2;
else
    d = 3;
return d;
```

./cbmc simple\_test\_case.c --cover mcdc --function foo

```
** 11 of 11 covered (100.0%)

** Used 5 iterations

Test suite:

a=13, b=45, c=-2147483644

a=-2147483636, b=45, c=-2147483644

a=-2147483636, b=-2147483634, c=-2147483644

a=13, b=-2147483634, c=5

a=13, b=-2147483634, c=5
```

# Supported features

CBMC has support for concurrency. Supports various memory models: SC, PSO, TSO

```
void* foo(void *arg)
{
   __CPROVER_atomic_begin();
   ++i;
   __CPROVER_atomic_end();
   return 0;
}
```

./cbmc --mm sc thread\_example.c --trace

```
int i=0;
int main()
{
  pthread_t th1, th2;
  pthread_create(&th1, 0, foo, 0);
  pthread_create(&th2, 0, foo, 0);
  pthread_join(th1, 0);
  pthread_join(th2, 0);
  assert(i==1);
  assert(i==2);
}
```

```
** Results:
[main.assertion.1] assertion i==1: FAILURE
[main.assertion.2] assertion i==2: SUCCESS

** 1 of 2 failed (2 iterations)
VERIFICATION FAILED
```

#### References

- CBMC C Bounded Model Checker, Tools and Algorithms for the Construction and Analysis of Systems, TACAS 2014 – Daniel Kroening et al
- CBMC Ver 1.0 Slides by Daniel Kroenig
- CBMC CMU Website: www.cprover.org/cbmc
- Introduction to CBMC Arie Gurfinkel, December 5, 2011
- Rui Goncalo : Automated test Generation using CBMC
- Ajith K J: VBMC A VHDL bounded Model Checker, Feb 20, 2017

Thank you.