# INFORMATION RETRIEVAL

# PRACTICAL-2

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Aim: Implement Dynamic programming algorithm for computing the edit distance between strings s1 and s2.

## **DESCRIPTION:**

Levenshtein distance algorithm

The Levenshtein algorithm calculates the least number of edit

operations that are necessary to modify one string to obtain another

string

### Steps:

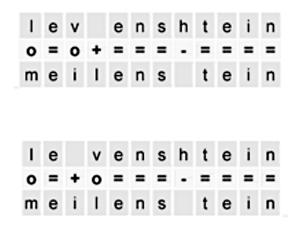
- 1. A matrix is initialized measuring in the (m, n) cell the Levenshtein distance between the m-character prefix of one with the n-prefix of the other word.
- 2. The matrix can be filled from the upper left to the lower right corner.
- 3. Each jump horizontally or vertically corresponds to an insert or a delete, respectively.
- 4. The cost is normally set to 1 for each of the operations.

- 5. The diagonal jump can cost either one, if the two characters in the row and column do not match else 0, if they match. Each cell always minimizes the cost locally.
- 6. This way the number in the lower right corner is the Levenshtein distance between both words.

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Levenshtein Distance Algorithm
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Example:
EDITDISTANCE(s_1, s_2)
  1 int m[i,j] = 0
 2 for i \leftarrow 1 to |s_1|
 3 do m[i, 0] = i
  4 for j ← 1 to |s<sub>2</sub>|
      \mathbf{do}\ m[0,j]=j
      for i \leftarrow 1 to |s_1|
 7
      do for j \leftarrow 1 to |s_2|
 8
          do m[i,j] = \min\{m[i-1,j-1] + \text{if } (s_1[i] = s_2[j]) \text{ then } 0 \text{ else } 1\text{fi},
 9
                                 m[i-1,j]+1,
10
                                 m[i, j-1]+1
11 return m[|s_1|, |s_2|]
```

		m	е	İ		е	n	S	t	е	İ	n
	0	1	2	3	4	5	6	7	8	9	10	11
	1	1	2	3	3	4	5	6	7	8	9	10
е	2	2	1	2	3	3	4	5	6	7	8	9
٧	3	3	2	2	3	4	4	5	6	7	8	9
е	4	4	3	3	3	3	4	5	6	6	7	8
n	5	5	4	4	4	4	3	4	5	6	7	7
S	6	6	5	5	5	5	4	3	4	5	6	7
h	7	7	6	6	6	6	5	4	4	5	6	7
t	8	8	7	7	7	7	6	5	4	5	6	7
е	9	9	8	8	8	7	7	6	5	4	5	6
İ	10	10	9	8	9	8	8	7	6	5	4	5
n	11	11	10	9	9	Ø	8	8	7	6	5	4



```
print('Name and Roll No')
a=input("Enter a String: ")
b=input("Enter a String: ")
A=[[0 \text{ for } i \text{ in } range(len(b)+1)] \text{ for } j \text{ in } range(len(a)+1)]
for i \text{ in } range(len(a)+1):
A[i][0]=i
for j \text{ in } range(len(b)+1):
A[0][j]=j
for i \text{ in } range(1,len(a)+1):
for j \text{ in } range(1,len(b)+1):
if a[i-1]==b[j-1]:
A[i][j]=A[i-1][j-1]
else:
```

```
delete=1+A[i-1][j]
    replace=1+A[i-1][j-1]
    insert=1+A[i][j-1]
    A[i][j]=min(insert,delete,replace)
print("Levenshtein Distance",A[len(a)][len(b)])
```

#### Screenshot:

