

CS 348 Computer Networks Lec 8

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Disclaimer: These slides are adapted from Computer Networking: A Top-down Approach by Kurose & Ross, 7th ed. and lecture slides of cs 168-2020 (http://cs168.io/) by Prof. Sylvia Ratnasamy

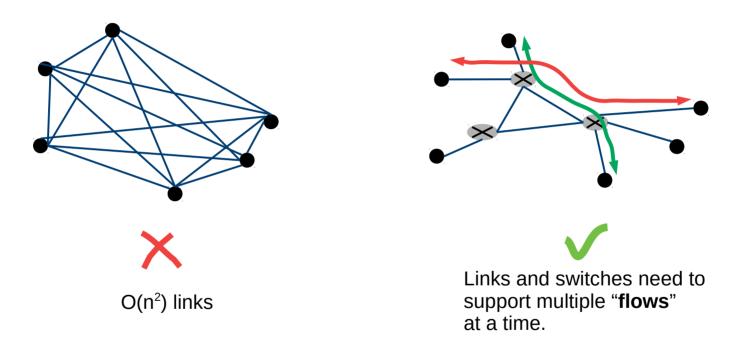
Questions

• Why does the Internet use **Packet Switching**?

What are the alternatives?

• What is the history of Packet switching?

Connecting together multiple hosts



Network resources (links, switches) must be shared between end-hosts.

Sharing Network Resources

How can network resources (links, switches) be shared between flows?

Two approaches:

• **Reservations:** end-hosts explicitly reserve bandwidth along the entire path at the start of a communication session, and release it when done.

• **Best-effort:** just send data packets across. The packets use up resources (link bandwidth, switch buffer space) as and when needed ("on demand").

Sharing Network Resources

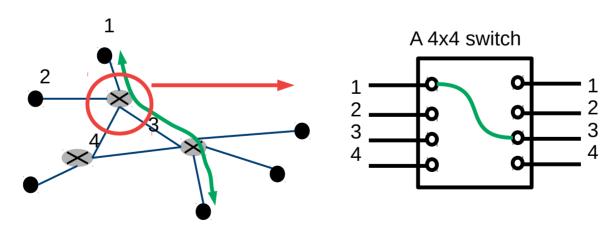
Two approaches:

• **Reservations:** this is the approach used by **Circuit switching**

• **Best-effort:** this is the approach used by **Packet switching**

Circuit Switching

- Network resources **along a route** are **reserved** between the end-hosts for the duration of a communications session.
- The terminology comes from early telephone exchanges. A "circuit" was literally formed between the end-hosts by configuring switches in telephone exchanges.



Circuit Switching

- Network resources **along a route** are **reserved** between the end-hosts for the duration of a communications session.
 - source sends a reservation request to the destination
 - switches "establish a circuit"
 - source starts sending data
 - source sends a "teardown circuit" message
- Currently used in the Public Switched Telephone Network (PSTN)

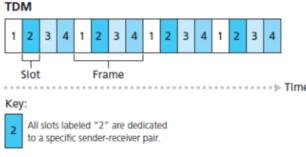
Circuit Switching

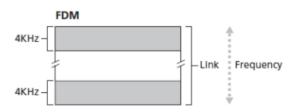
• One link doesn't imply one communication channel. A single link can support multiple "channels" at a time:

Time division multiplexing (TDM)

Frequency division multiplexing (FDM)

• Thus "channels" (not links) along a path are reserved for a communication session.





How do they compare along the following axes?

- Service model for applications (are there any performance guarantees?)
- Efficiency
- Handling failures
- Implementation complexity

Service model for Applications:

- **Circuit switching:** guaranteed bandwidth for the duration of the session
- Packet switching: no bandwidth/performance guarantees for a flow:
 - no guarantee for allocation of space inside switch buffers
 - no guarantee for link bandwidth

Efficiency:

- **Circuit switching**: less efficient for "bursty" traffic.
 - Reserved bandwidth cannot be given to other "flows" even when unutilized.
 - Time spent for setting up/tearing down circuit. Inefficient for short flows.
- **Packet switching:** more efficient for "bursty" traffic.

Handling Failures: What happens if a link goes down?

• Packet switching:

- Network must detect failure
- Network recalculates routes (Job of the routing protocols)
- Endhosts and individual flows do nothing special except cope with the temporary loss of service.

Circuit switching:

- All of the above
- In addition, Endhosts must detect failure, teardown old reservations, send a new reservation request. All impacted endhosts must do this, for each impacted flow

Implementation Complexity

Packet switching:

 Network does not need to maintain per-flow "state" (the network does not need to remember anything about which "flows" are active)

• Circuit switching:

- Setup/teardown
- Network needs to maintain per-flow state.

History of Packet Switching

- 1961: Leonard Kleinrock proposes the idea of packet switching in his MIT doctoral thesis.
- 1964: Packet switching studied and developed further by Paul Baran at the RAND corp. Motivation was to develop resilient and survivable communication systems.
- 1965: Davies designed a store-and-forward packet switching system and coined the term "packet".
- 1969: ARPANET is designed to work on packet switching and becomes functional with a few nodes

(Ref: https://en.wikipedia.org/wiki/Packet_switching)

Questions (Going forward)

- How are forwarding tables populated? How/when do routers update the forwarding tables when some link goes down?
- How does TCP provide reliable transmission over an unreliable channel?
- How does TCP perform congestion control?
- What are some Application-layer protocols? How do they work?
- What is the interface between the Application layer and the transport layer?

Reference

• Kurose and Ross, Section 1.3.2