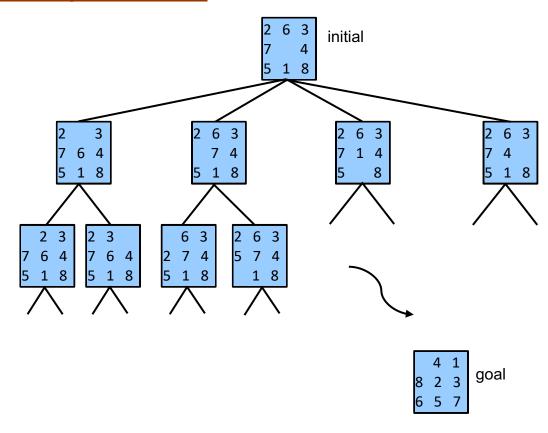
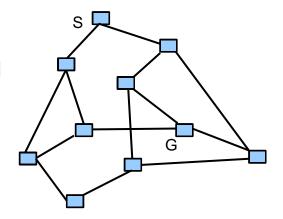
State space search



- How to find a path from initial board to goal board
- State space search
 - start at the initial state
 - search the states in a systematic order
 - stop when the goal state is found
 - display the path from initial state to goal state
- Two types of state space search
 - uninformed / blind search
 - informed / heuristic search

Applications of state space search

- State space search may be used in problems that involve
 - multiple states
 - actions that change states
 - initial state and goal state
 - finding path from initial to goal
- Some applications
 - board puzzles
 - two player board games
 - automous vehicle driving
 - robot task planning



<u>Uninformed / blind search</u>

- Search is done blindly
- Search does not use any problem specific information
- Search may take long time
- Uninformed search algorithms
 - breadth first search
 - depth first search
 - depth limited, iterated, bidirectional search
 - uniform cost search

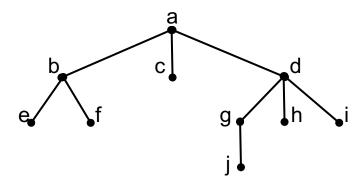
Informed / heuristic search

- Search is done intelligently
- Search uses some problem specific information

- Search is likely to run fast
- Informed search algorithms
 - best first search
 - A* search

Breadth first search on tree

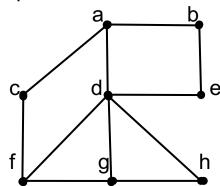
- Tree



- Breadth first search order - a b c d e f g h i j

Breadth first search on graph

- Graph



- Initial state a
- Breadth first search order a c d b f g h e

Breadth first search algorithm

```
- open list has initial state
 closed list is empty
 while open list is not empty
 {
     S = remove the first state from open list
     insert S into closed list
     if S is goal state
        display path from initial state to S
        stop search
     else
         generate successors of S
         for each successor C
             if C is not in open list and C is not in closed list
                 insert C at the end of open list
    }
 say there is no path
```

- Open list is a queue, operates in first in first out order, FIFO
- Closed list is not needed if search is done on tree
- Run time O(b^d), space usage O(b^d), b is branching factor, d is depth of goal state

Open/closed lists in breadth first search

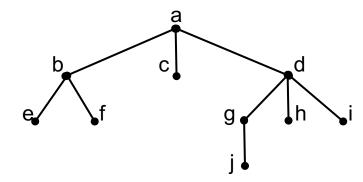
- Open and closed lists for the tree shown above

```
open list look closed lista a
```

```
bcd
         b
                 а
c d e f
                 a b
         С
d e f
         d
                 a b c
efghi
                 abcd
         е
fghi
         f
                 abcde
                 abcdef
ghi
                 abcdefg
hij
         h
                abcdefgh
ij
                 abcdefghi
                 abcdefghij
```

Depth first search on tree

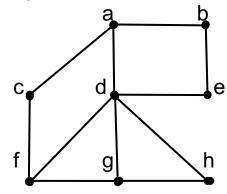
- Tree



- Depth first search order - a b e f c d g j h i

Depth first search on graph

- Graph



- Initial state a
- Depth first search order a c f d e b h g

Depth first search algorithm

```
- open list has initial state
  closed list is empty
  while open list is not empty
{
     S = remove the first state from open list
     insert S into closed list
     if S is goal state
          display path from initial state to S
          stop search
     else
     {
          generate successors of S
          for each successor C
               if C is not in open list and C is not in closed list
                    insert C at the beginning of open list
          }
     }
     say there is no path
```

- Open list is a stack, operates in last in first out order, LIFO
- Closed list is not needed if search is done on tree
- Run time O(b^m), space usage O(b^m) if search is done in graph run time O(b^m), space usage O(bm) if search is done in tree b is branching factor, m is depth of tree

Open/closed lists in depth first search

- Open and closed lists for the tree shown above

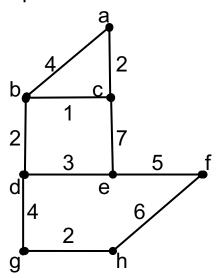
 open list 	look	closed list
а	a	
b c d	b	а
e f c d	е	a b
fcd	f	a b e
c d	С	a b e f
d	d	abefc
ghi	g	abefcd
j h i	j	abefcdg
h i	h	abefcdgj
İ	i	abefcdgjh
		abefcdgjhi

Depth limited, iterated, bidirectional search algorithms

- Depth limited search
 - choose a depth limit
 - when the depth limit is reached, do not create successors
 - depth limit depends on time/space constraints
- Iterated depth limited search
 - do depth limited search with limit = 1, limit = 2, limit = 3, limit = 4, and so on
 - number of iterations depends on time/space constraints
- Bidirectional search
 - search forward from initial to goal
 - search backward from goal to initial
 - stop when forward and backward searches meet

Uniform cost search on graph

- Finding shortest path from initial state to goal state
- Graph



- Initial state a goal state e
- Shortest path a c b d e cost 8

- open list	look	closed list
a/0	а	
b/4 c/2	С	а
b/3 e/9	b	a c
e/9 d/5	d	a c b
e/8 g/9	е	acbd
g/9		acbde

Uniform cost search algorithm

```
- open list has initial state with cost = 0
 closed list is empty
 while open list is not empty
 {
     S = remove the minimum cost state from open list
     insert S into closed list
     if S is goal state
        display path from initial state to S
        stop search
     else
         generate successors of S
         for each successor C
             cost of C = cost of S + edge cost from S to C
             if C is not in closed list
                 if C is not in open list
                    insert C into open list
                 else
                 {
                     compare costs of new copy of C and
                     old copy of C in open list
                     if new cost < old cost
                        replace old copy of C in open list with
                        new copy of C
             }
     }
 say there is no path
```

- If all edges have the same cost then the uniform cost search behaves like the breadth first search