The Bucket Sort Algorithm

Presentation by: John Patrick Tobias

The **Bucket Sorting Algorithm**, also known as bin sort, from the name of the algorithm we can derive a general idea on how the algorithm would use or work in order to sort an array of values.

In this sorting algorithm, we are going to use **buckets** in order to **segregate different values into their respective buckets** and sort the values inside the buckets with the use of **insertion sort**.

Lastly, we **return the values into the original array** and they would be sorted, assuming that we start from the first bucket up to the last.

Note: Bucket Sort is only <u>O(n)</u> if the number of buckets is equal to the length of the array being sorted.

The Pseudocode

```
function bucketSort(array, n)
buckets ← new array of n empty lists
for i = 0 to (length(array)-1) do
  insert array[i] into buckets[assigned_bucket]
for i = 0 to n - 1 do
  nextSort(buckets[i]);
return the concatenation of buckets[0], ...., buckets[n-1]
```

Step-by-step

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

Let's call this array, A.

Step 1: Get size of array

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

Array A has a size (N) of 10.

Step 2: Get the max value

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

The *max* value in the array is 90.

N = 10 max = 90

Step 3: Get amount of buckets

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

Then we get a **right amount of buckets** where we can put and sort the values into.

In this example, we can have 10 buckets (to satisfy O(n)).

N = 10 max = 90 buckets = 10

Step 4: Get the divider

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

Then we calculate the **divider** which is required in order to put the values into their respective buckets/bins.

With the formula: divider = ceil((max + 1) / bucket size)

N = 10 max = 90 buckets = 10

Step 4: Get the divider

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

In this case our divider would be:

Step 5: Assigning to buckets

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

Formula for getting the assigned bucket for each value:

assigned_bucket = floor(A[i] / divider)

Now let's solve everything on the board.

Step 6: Insertion Sort

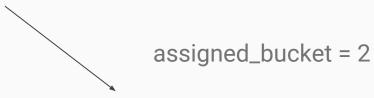
N = 10 max = 90 buckets = 10 divider = 10

When a new value is inserted into the bucket, **insertion sort** is then performed in order to sort the newly inserted value efficiently.

0	1	2	3	4	5	6	7	8	9
24	32	5	76	90	12	62	8	18	22

0	1	2	3	4	5	6	7	8	9
1	1	1	1	1	1	1	1	1	1

0	1	2	3	4	5	6	7	8	9
1	32	5	76	90	12	62	8	18	22



0	1	2	3	4	5	6	7	8	9
1	1	24	1	1	1	1	1	1	/
		1							

0	1	2	3	4	5	6	7	8	9
1	32	5	76	90	12	62	8	18	22

0	1	2	3	4	5	6	7	8	9
1	1	24	1	1	1	1	1	1	1
		1							

0	1	2	3	4	5	6	7	8	9
1	1	5	76	90	12	62	8	18	22
			assig	ned_	bucke	et = 3			
		`							
0	1	2	3	4	5	6	7	8	9
1	1	24	32	1	1	1	1	1	1
		,	1						

0	1	2	3	4	5	6	7	8	9
1	1	5	76	90	12	62	8	18	22

assigned_bucket = floor
$$(5 / 10) = 0$$

0	1	2	3	4	5	6	7	8	9
1	1	24	32	1	1	1	1	1	1
		1	1						

Bucket Sort Algorithm

0	1	2	3	4	5	6	7	8	9
1	1	1	76	90	12	62	8	18	22



assigned_bucket = **0**

0	1	2	3	4	5	6	7	8	9
5	1	24	32	1	1	1	1	1	1
1		1	1						

0	1	2	3	4	5	6	7	8	9
1	/	1	76	90	12	62	8	18	22

assigned_bucket = floor
$$(76 / 10) = 7$$

0	1	2	3	4	5	6	7	8	9
5	1	24	32	1	1	1	1	1	/
1		1	1						

Bucket Sort Algorithm

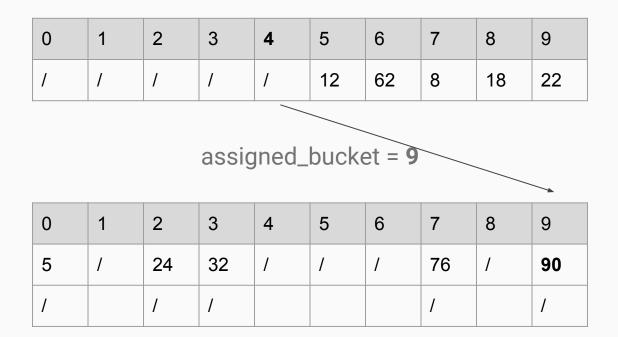
0	1	2	3	4	5	6	7	8	9
1	1	1	1	90	12	62	8	18	22

assigned_bucket = 7

0	1	2	3	4	5	6	7	8	9
5	1	24	32	1	1	1	76	1	1
1		1	1				1		

0	1	2	3	4	5	6	7	8	9
1	/	1	1	90	12	62	8	18	22

0	1	2	3	4	5	6	7	8	9
5	1	24	32	1	1	1	76	1	1
1		1	1				1		



0	1	2	3	4	5	6	7	8	9
1	1	1	1	1	12	62	8	18	22

assigned_bucket = floor
$$(12 / 10) = 1$$

0	1	2	3	4	5	6	7	8	9
5	1	24	32	1	1	1	76	1	90
1		1	1				1		1

Bucket Sort Algorithm

0	1	2	3	4	5	6	7	8	9
1	1	1	1	1	1	62	8	18	22

C)	1	2	3	4	5	6	7	8	9
1		1	1	1	1	1	62	8	18	22

0	1	2	3	4	5	6	7	8	9
5	12	24	32	1	1	1	76	1	90
1	1	1	1				1		/

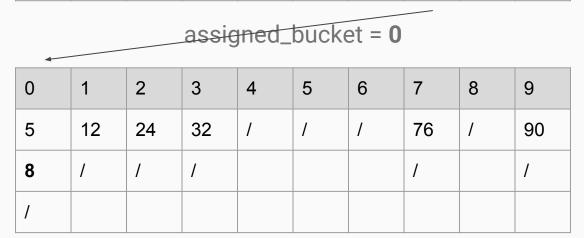
0	1	2	3	4	5	6	7	8	9
1	1	1	/	1	1	1	8	18	22

0	1	2	3	4	5	6	7	8	9
5	12	24	32	1	1	62	76	1	90
1	1	1	1			1	1		1

0	1	2	3	4	5	6	7	8	9
1	/	1	1	1	1	1	8	18	22

0	1	2	3	4	5	6	7	8	9
5	12	24	32	1	1	62	76	1	90
1	1	1	1			1	1		1

0	1	2	3	4	5	6	7	8	9
1	/	1	1	1	/	/	/	18	22



0	1	2	3	4	5	6	7	8	9
1	1	1	1	1	1	1	1	18	22

0	1	2	3	4	5	6	7	8	9
5	12	24	32	1	1	1	76	1	90
8	1	1	1				1		1
1									

Bucket Sort Algorithm

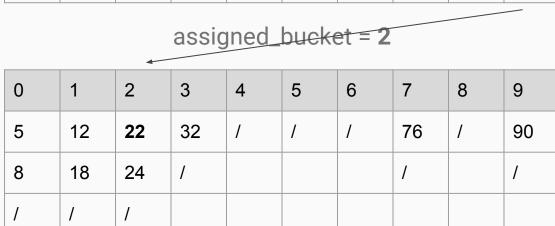
0	1	2	3	4	5	6	7	8	9
/	1	1	/	/	1	1	/	1	22

0 1 2 3 4 5 6 7 8 9
5 12 24 32 / / / 76 / 90
8 18 / / / / / / / / / / /

0	1	2	3	4	5	6	7	8	9
1	1	1	1	1	1	1	1	1	22

0	1	2	3	4	5	6	7	8	9
5	12	24	32	1	1	1	76	1	90
8	18	1	1				1		1
1	1								

0	1	2	3	4	5	6	7	8	9
1	/	1	1	1	1	1	1	1	1



Bucket Sort Algorithm

The Final Bucket

0	1	2	3	4	5	6	7	8	9
5	12	22	32	1	1	1	76	1	90
8	18	24	1				1		1
1	1	1							

Final Step: Put back EVERYTHING

When all values have been put into their respective buckets, we then **put them back** into array *A* starting from the first bucket up to the last bucket.

DONE!

0	1	2	3	4	5	6	7	8	9
5	12	22	32	1	1	1	76	1	90
8	18	24	1				1		1
1	1	1							

0	1	2	3	4	5	6	7	8	9	
5	8	12	18	22	24	32	62	76	90	

The values are then sorted!

Assumptions and Limitations

Assumes that all values belong to any range of a bucket.

O(n) is only achieved when the number of buckets is equal to the length of the array currently being sorted. If not, it can achieve $O(n^2)$ which would be like insertion sort, which can make bucket sort not necessary.