50 Monoids, Functors and Monads

Monoids, Functors and Monads

- These terms come from category theory. But really, their definitions as far as Scala is concerned are quite simple:
 - for instance, a monoid is just the type of thing we were getting at with our List.x2 method (a.k.a sum)—where List's underlying type must be a monoid:
 - something that can be operated on dyadic-ally and something which has a "zero" (identity) value:

```
def sum: A = {
    @tailrec def inner(as: List[A], x: A): A = as match {
    case Nil => x
    case Cons(hd,tl) => internal(tl,x++hd)
    }
    inner(this,0)
}
```

 You don't really need to remember all this stuff. I will review what's important at the end.

Monoids

- A monoid consists of the following:
 - a type A
 - an associative binary operation, op, that takes two values of type A and combines them into one such that:

```
op(op(x,y),z) == op(x,op(y,z))
for any x: A, y: A, z: A.
```

• an "identity" value: zero: A that is an identity for op, i.e.

```
op(x,zero) == x \text{ for any } x: A.
```

• For example:

```
trait Monoid[A] {
  def op(a1: A, a2: A): A
  def zero: A
}
object Monoid {
  val stringMonoid = new Monoid[String] {
    def op(a1: String, a2: String) = a1 + a2
    val zero = ""
  }
  def listMonoid[A] = new Monoid[List[A]] {
    def op(a1: List[A], a2: List[A]) = a1 ++ a2
    val zero = Nil
  }
}
```

Monoids (2)

• There are two important functions on collections that we met briefly in Recursion: *foldRight* and *foldLeft*:

```
def foldLeft[A,B](z: B)(fl: (B, A) => B): B
def foldRight[A,B](z: B)(fr: (A, B) => B): B
```

- These methods are equivalent providing that fl and fr are associative—they simply work through a container in different directions—but only foldLeft is tail-recursive for List so we'll concentrate on that.
- Here's foldLeft implemented for our own List class from week 2:

```
def foldLeft[B](z: B)(f: (B,A)=>B): B = this match {
  case Nil => z
  case Cons(hd,tl) => tl.foldLeft(f(z,hd))(f)
}
```

Let's reimplement sum on List in terms of foldLeft...

```
def sum[B]: B = foldLeft(B.zero)(B.plus(_,_))
```

 Just one snag: B.zero and B.plus are not defined. But if, for any B, zero and plus were defined we'd be all set.

Monoids (3)

 Let's go back to our stringMonoid and create an intMonoid too (we'll also make op explicit as plus):

```
object Monoid {
  val stringMonoid = new Monoid[String] {
    def plus(a1: String, a2: String) = a1 + a2
    val zero = ""
  }
  def intMonoid[A] = new Monoid[Int] {
    def plus(a1: Int, a2: Int) = a1 + a2
    val zero = 0
  }
}
```

So we would have:

```
def sum[B]: B = foldLeft(intMonoid.zero)(intMonoid.plus(_,_))
```

Therefore, in general, a *Monoid* is something that is *foldable*. We could
make this explicit by defining the following type constructor:

```
trait Foldable[F[_]] extends Functor[F] {
    def foldLeft[A,B](z: B)(f: (B, A) => B): B
    def foldRight[A,B](z: B)(f: (A, B) => B): B
}

def foldableList[A] = new Foldable[List] {
    def map[A,B](as: List[A])(f: A => B): List[B] = as map f
    def foldLeft[A,B](z: B)(f: (B, A) => B): B = ??? // tail recursive
    def foldRight[A,B](z: B)(f: (A, B) => B): B = ??? // NOT tail recursive
}
```

Functors

- Functor is another term from category theory:
 - a *functor* is just a mapping between two categories (or types in Scala):

```
trait List[+A] { def map[B](f: A=>B): List[B] }
trait Option[+A] { def map[B](f: A=>B): Option[B] }
trait Stream[+A] { def map[B](f: A=>B): Stream[B] }
etc. etc. etc.
```

- Notice anything?
 - All the definitions are identical—only the implementations differ.
 - Repetitive code like this is anathema to functional programmers!
 - Let's define a functor trait:

```
trait Functor[F[_]] {
   def map[A, B](fa: F[A])(f: A => B): F[B]
}
object Functor {
   def listFunctor = new Functor[List] {
     def map[A,B](as: List[A])(f: A => B): List[B] = as map f
   }
}
```

Monads

- Monad also comes from category theory. You've heard me mention monads already:
 - A Monad <u>isa</u> Functor (i.e. it implements map):

```
trait List[+A] { def map[B](f: A=>B): List[B] }
trait Option[+A] { def map[B](f: A=>B): Option[B] }
trait Stream[+A] { def map[B](f: A=>B): Stream[B] }
etc. etc. etc.
```

 Remember how map2 was basically identical for several different container types? These were all monads! So let's define map2 once and for all...

```
trait Monad[F[_]] extends Functor[F] {
  def map2[A,B,C](ma: F[A], mb: F[B])(f: (A,B)=>C): F[C] =
     flatMap(ma)(a => map(mb)(b => f(a,b)))
}
```

 But flatMap isn't defined: (Except that since this definition is the canonical definition of map2 and we are essentially defining monad as things which support map2, ergo Monad must define flatMap too:

```
trait Monad[F[_]] extends Functor[F] {
  def flatMap[A,B](ma: F[A])(f: A=>F[B]): F[B]
  def map2[A,B,C](ma: F[A], mb: F[B])(f: (A,B)=>C): F[C] =
    flatMap(ma)(a => map(mb)(b => f(a,b)))
}
```

Monads (2)

Actually, there's a certain flexibility in exactly how we
define the primitive methods of *Monad*. Let's think about
map. If we had a method unit which took a value and
simply wrapped it (like a single-element List) then we
could actually write map in terms of flatMap:

```
trait Monad[F[_]] extends Functor[F] {
  def unit[A](a: => A): F[A] // abstract
  def flatMap[A,B](ma: F[A])(f: A=>F[B]): F[B] // abstract
  def map[A,B](ma: F[A])(f: A=>B): F[B] = flatMap(ma)(a => unit(f(a)))
  def map2[A,B,C](ma: F[A], mb: F[B])(f: (A,B)=>C): F[C] =
     flatMap(ma)(a => map(mb)(b => f(a,b)))
```

In this definition of *Monad*, *unit* and *flatMap* are abstract methods—they must be defined by implementers of *Monad*. The other two methods are concrete methods defined in terms of the first two.

- So, unit+flatMap is one possible form.
- But why do we care about all this stuff? Because of composability.

Monads (3)

- Articles which explain monads:
 - https://medium.com/@sinisalouc/demystifying-the-monadin-scala-cc716bb6f534

Review: what do you need to remember?

- A type that implements map is a functor.
- A monad is a functor but not all functors are monads.
- For-comprehensions work on monads.
- The reduce method works only for monoids.