

# DM – Introduction to graphs

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# Outline

- 1 Graph – Definition and examples
  - (Undirected) Graph
  - Directed graph
- 2 Basic properties of graphs
  - Undirected case
  - Directed case
- 3 Some special simple graphs
- 4 Subgraphs and graph operations
- 5 Graph representation
  - Adjacency list
  - Adjacency matrix
  - Incidence matrix

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## Some examples



Figure: Regional routes of VNA (as of 2017).<sup>1</sup>

<sup>1</sup>Source: <https://www.viags.vn/tin-tuc/tin-chuyen-nganh/vietnam-airlines-routemap-wordwide>

# Some examples

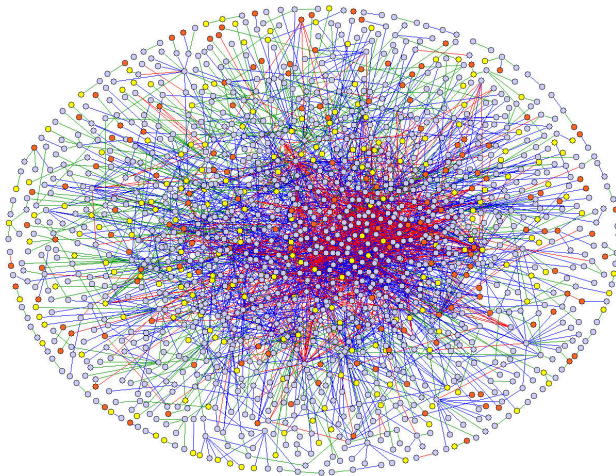
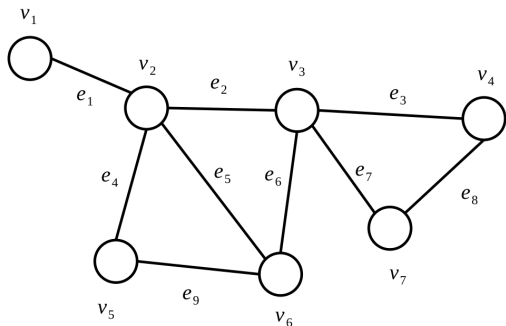


Figure: Map of human protein interaction.<sup>2</sup>

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<sup>2</sup>Source: [https://www.mdc-berlin.de/news/archive/2008/20080910-erwin\\_schr\\_dinger\\_prize\\_2008\\_goes\\_to\\_resea](https://www.mdc-berlin.de/news/archive/2008/20080910-erwin_schr_dinger_prize_2008_goes_to_resea)

# Graph



- Each vertex is represented by a dot or a circle
- Each edge is represented by a line connecting its endpoints

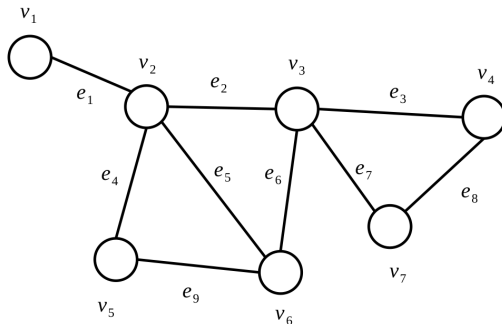
## Definition

A (undirected) **graph**  $G = (V, E)$  is defined by:

- 1 A non-empty set  $V$  of **vertices**;
- 2 A set  $E$  of **edges**, which are unordered pairs of vertices.

**Note:** In this course, we only consider *finite* graphs.

# Graphs



- $V = \{v_1, v_2, v_3, v_4, v_5, v_6, v_7\}$
- $E = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7, e_8, e_9\}$

# Multigraph and loop

## Definition

- A graph is **simple** if there is at most one edge between any pair of vertices.
- Edges connecting the same pair of vertices are called **multiple edges**. A graph having multiple edges is also called a **multigraph**.
- A **loop** is an edge connecting a vertex to itself.

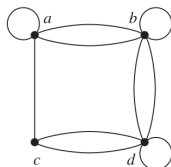


Figure: Rosen, p. 650

**Note:** By convention, “graphs” refer to simple graphs without loops unless otherwise stated.



# Examples

## Social networks:

- Each vertex represents a person
- Two friends are connected by an edge

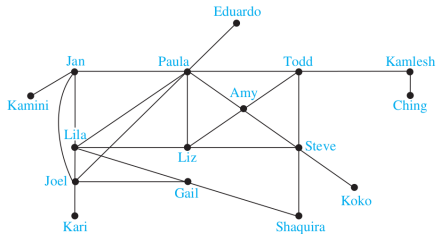


Figure: Rosen, p. 645

# Examples

## Niche overlap graphs (ecology):

- Each vertex represents a species
- Two competing species are connected by an edge

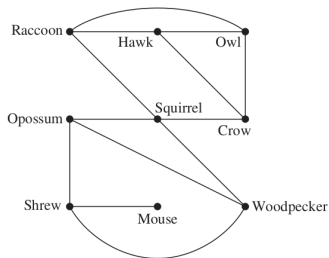


Figure: Rosen, p. 648

# Examples

## Protein interaction graphs (biology):

- Each vertex represents a protein
- Two proteins that interact are connected by an edge

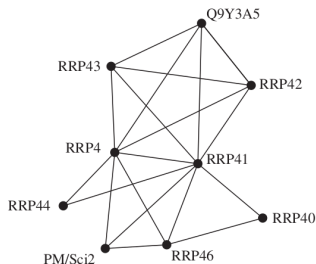


Figure: Rosen, p. 648

# Examples

## Computer networks:

- Each vertex represents a server
- Each edge represents a link

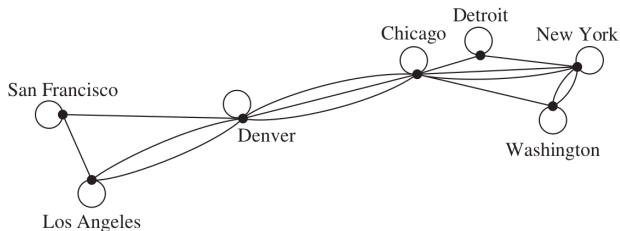


Figure: Rosen, p. 642

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# An example

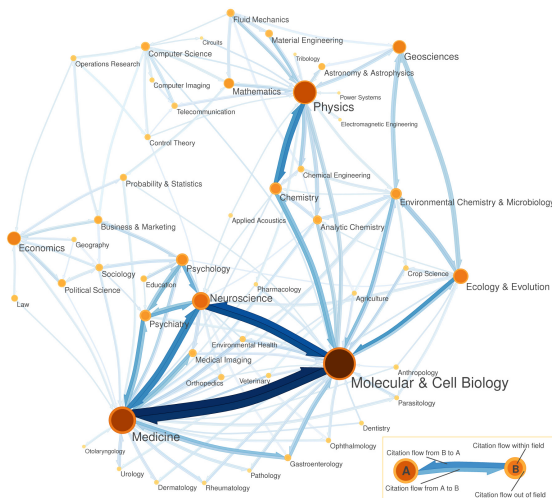


Figure: A map of science disciplines.<sup>3</sup>

<sup>3</sup>Source: <http://images.math.cnrs.fr/Representer-les-mondes.html>

# Directed graph

## Definition

A **directed graph** (or **digraph**)  $G = (V, A)$  is defined by:

- 1 A non-empty set  $V$  of **vertices** (or **nodes**);
- 2 A set  $A$  of **arcs**, which are ordered pairs of vertices.

The notions of *simple directed graphs*, *multiple arcs*, *loops* and *directed multigraphs* are defined similarly to the undirected case.

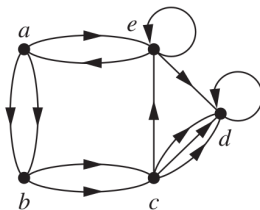


Figure: Rosen, p. 650

# Examples

## Round-robin tournament:

- Each node represents a team
- Each team play against every other team, an arc goes from the winning team to the losing team (assume that there are no ties)

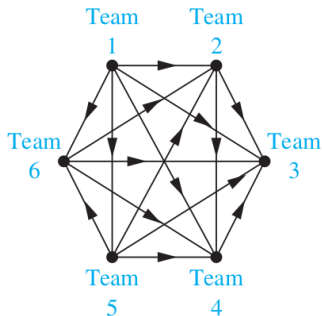


Figure: Rosen, p. 649



# Examples

## Dependency (software):

- Each node represents a module
- An arc  $(a, b)$  is added if module  $b$  depends on module  $a$

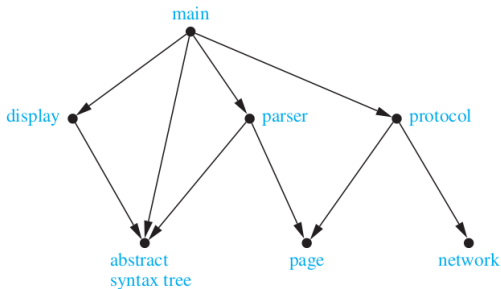


Figure: Rosen, p. 647

# Examples

## Precedence graph:

- Each node represents a statement
- An arc  $(S_i, S_j)$  is added if statement  $S_i$  must wait for statement  $S_j$  to be executed

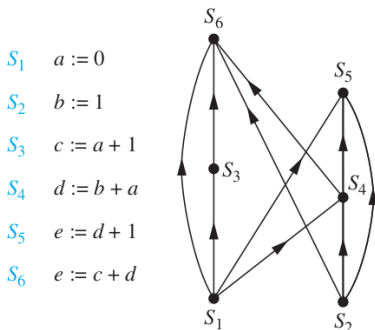


Figure: Rosen, p. 647

# Examples

**... and many other examples:**

- citation graphs (academic)
- co-author graphs (academic)
- genealogical graphs
- road networks
- the Web
- etc.

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# Adjacency, neighborhood

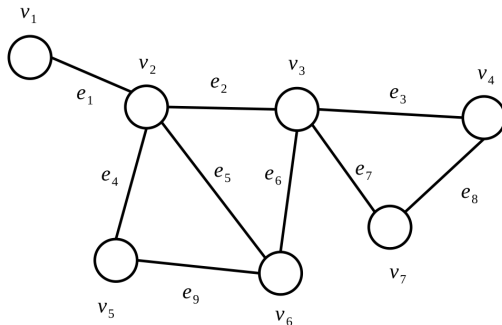
Let  $G = (V, E)$  be a graph.

## Definition

- Two vertices  $u$  and  $v$  are **adjacent** if they are endpoints of the same edge  $e$ . In this case,  $u$  is a **neighbor** of  $v$  and the edge  $e$  is **incident** with the vertices  $u$  and  $v$ .
- The set of all neighbors of a vertex  $v$  is the **neighborhood** of  $v$  and is denoted by  $N(v)$ . The neighborhood of a set  $A$  of vertices is the union of the neighborhoods of the vertices in  $A$ :  $N(A) = \bigcup_{v \in A} N(v)$ .

# Adjacency, neighborhood

Example:

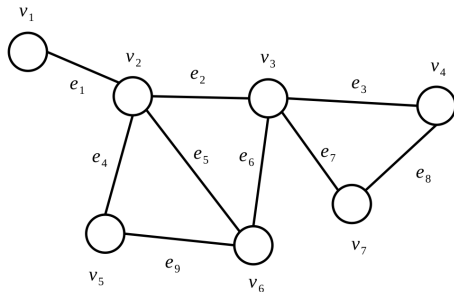


- $v_1$  is adjacent to  $v_2$ ;
- $v_4$  is not adjacent to  $v_2$ ;
- $N(v_2) = \{v_1, v_3, v_5, v_6\}$ ;
- $N(\{v_1, v_2\}) = \{v_1, v_2, v_3, v_5, v_6\}$ ;
- $N(\{v_2, v_3\}) = V$ .

# Degree

## Definition

The **degree** of a vertex  $v$ , denoted by  $\deg(v)$  or simply  $d(v)$ , is the number of edges incident with it. A loop contributes 2 to the degree of its endpoint.



$$\deg(v_1) = 1, \deg(v_2) = 4.$$

# Degree

## Note:

- A vertex is *isolated* if it has degree 0; it is *pendant* if it has degree 1.
- In a multigraph, the degree of a vertex may be greater than the number of its neighbors!

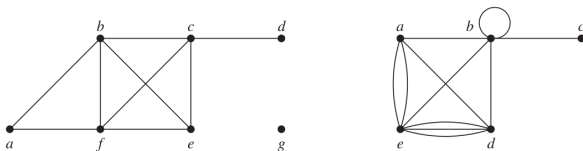


Figure: Rosen, p. 652



# The handshaking theorem

## Theorem

*Let  $G = (V, E)$  be an undirected graph. Then*

$$\sum_{v \in V} \deg(v) = 2 |E|.$$

**Proof:** Each edge contributes 2 to the sum of degrees of all the vertices.

## Corollary

*There are an even number of vertices of odd degrees.*

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# Adjacency

Let  $G = (V, A)$  be a digraph.

## Definition

If  $a = (u, v)$  is an arc, then  $u$  is **adjacent to**  $v$  and  $v$  is **adjacent from**  $u$ ;  $u$  is the **initial node** and  $v$  is the **end node** of  $a$ ;  $a$  is an **out-going arc** of  $u$  and an **in-going arc** of  $v$ .

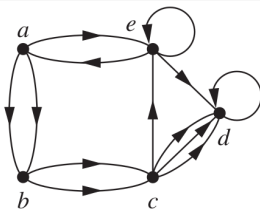


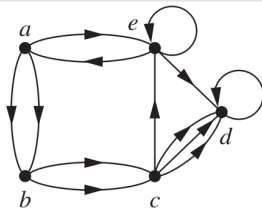
Figure: Rosen, p. 650

- $a$  is adjacent to  $b$ ,  $b$  is adjacent from  $a$ ;
- $b$  is *not* adjacent to  $a$ ,  $a$  is *not* adjacent from  $b$ .
- $e$  is adjacent to and from itself.

# Degrees

## Definition

The **out-degree** (resp. **in-degree**) of a node  $v$ , denoted by  $\deg^+(v)$  or  $d^+(v)$  (resp.  $\deg^-(v)$  or  $d^-(v)$ ), is the number of out-going (resp. in-going) arcs of  $v$ .



- $d^+(a) = 3$ ,  $d^-(a) = 1$ ;
- $d^+(d) = 1$ ,  $d^-(d) = 5$ .

Figure: Rosen, p. 650

## Theorem

Let  $G = (V, A)$  be a directed graph. Then

$$\sum_{v \in V} \deg^+(v) = \sum_{v \in V} \deg^-(v) = |A|.$$

# Complete graphs

## Definition

A **complete graph** of  $n$  vertices, denoted by  $K_n$ , is a simple graph having  $n$  vertices and such that every pair of vertices is connected by an edge.

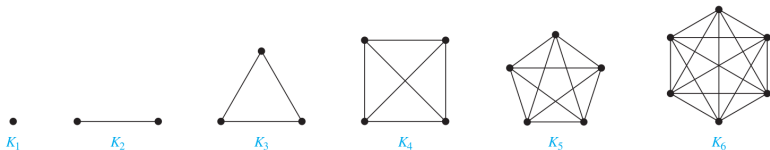


Figure: Rosen, p. 655

# Paths

## Definition

A **path** on  $n$  vertices, denoted by  $P_n$ , is a simple graph having  $n$  vertices  $\{v_1, v_2, \dots, v_n\}$  and  $n - 1$  edges  $\{\{v_i, v_{i+1}\} \mid 1 \leq i \leq n - 1\}$ .



Figure:  $P_6$

# Cycles

## Definition

A **cycle** on  $n \geq 3$  vertices, denoted by  $C_n$ , is a simple graph having  $n$  vertices  $\{v_1, v_2, \dots, v_n\}$  and  $n$  edges  $\{\{v_i, v_{i+1}\} \mid 1 \leq i \leq n\}$  ( $v_{n+1} \equiv v_1$ ).

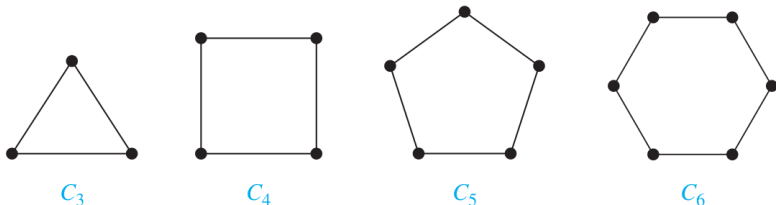


Figure: Rosen, p. 655

# Wheel graphs

## Definition

A **wheel graph** of  $n + 1$  vertices ( $n \geq 3$ ), denoted by  $W_n$ , is a simple graph obtained from  $C_n$  by adding a new vertex and connecting it to every vertex of  $C_n$ .

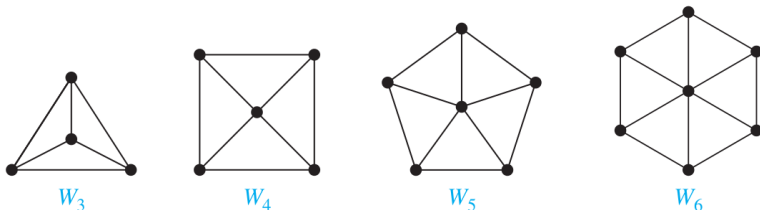


Figure: Rosen, p. 655



# Hypercubes

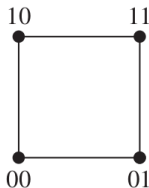
## Definition

A  *$n$ -dimensional hypercube*, or an  *$n$ -cube*, denoted by  $Q_n$ , is a simple graph in which

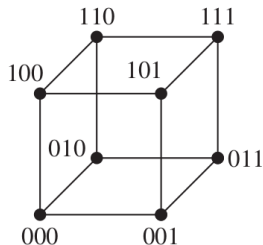
- the vertices are the binary strings of length  $n$
- two strings are adjacent iff. they differ at only one position.



$Q_1$



$Q_2$



$Q_3$

Figure: Rosen, p. 655

# Hypercubes

## Construction of $Q_{n+1}$ from $Q_n$ :

- ① Consider two copies  $Q_n$ :  $Q_n$  and  $Q'_n$
- ② Change the labels of  $Q_n$  and  $Q'_n$ :
  - $Q_n$ :  $s \rightarrow 0s$
  - $Q'_n$ :  $s \rightarrow 1s$
- ③ Add the edges  $\{0s, 1s\}$

# Bipartite graphs

## Definition

A graph  $G = (V, E)$  is a **bipartite graph** if  $V$  can be partitioned into two sets  $V_1$  and  $V_2$  so that every edge connects a vertex in  $V_1$  and a vertex in  $V_2$ .

The pair  $(V_1, V_2)$  is a **bipartition** of the vertices of  $G$ .

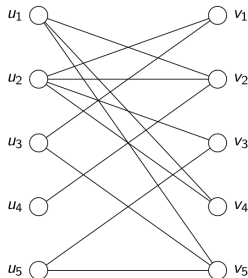


Figure: A bipartite graph.

# Bipartite graphs

## Examples:

- The path  $P_5$  is a bipartite graph.
- The cycle  $C_6$  is a bipartite graph.
- The cycle  $C_3$  is *not* a bipartite graph.
- The complete graph  $K_4$  is *not* a bipartite graph.
- The hypercubes  $Q_1, Q_2, Q_3$  are bipartite graphs.

# Recognizing bipartite graphs

Is this graph bipartite?

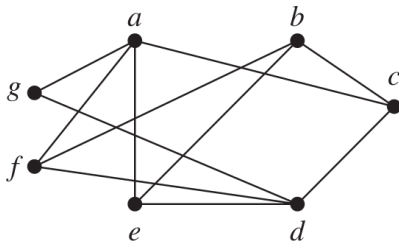


Figure: Rosen, p. 656

It is.

# Recognizing bipartite graphs

## Theorem

*A simple graph is bipartite iff. its vertices can be colored by two different colors so that no two adjacent vertices have the same color.*

## Examples:

- $C_n$  is bipartite iff.  $n$  is even.
- $P_n$  is bipartite for all  $n$ .
- $Q_n$  is bipartite for all  $n$ .

# Complete bipartite graphs

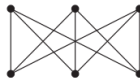
## Definition

A **complete bipartite graph**  $K_{m,n}$  is a bipartite graph with bipartition  $V = V_1 \cup V_2$  such that:

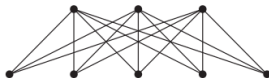
- $V_1$  has  $m$  vertices,  $V_2$  has  $n$  vertices;
- Two vertices are connected by an edge iff. one vertex is in  $V_1$ , the other is in  $V_2$ .



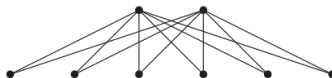
$K_{2,3}$



$K_{3,3}$



$K_{3,5}$



$K_{2,6}$

Figure: Rosen, p. 658

# Subgraphs

## Definition

A **subgraph** of a graph  $G = (V, E)$  is a graph  $H = (W, F)$  where  $W \subset V, F \subset E$ .

## Examples:

- Every graph is a subgraph of itself.
- The path  $P_n$  is a subgraph of the cycle  $C_m$  if  $n \leq m$ .
- The cycle  $C_n$  is a subgraph of the wheel  $W_n$ .
- Every graph on  $n$  vertices is a subgraph of  $K_n$ .



# Subgraphs

## Definition

Let  $G = (V, E)$  be a simple graph and let  $W$  be a subset of vertices. The **subgraph induced** by  $W$  is the subgraph of  $G$  where:

- The set of vertices is  $W$ ;
- The set of edges consists of all edges having both endpoints in  $W$ .

## Examples:

- $C_n$  is the subgraph of  $W_n$  induced by the “outer” vertices.
- If  $m \leq n$ ,  $K_m$  is the subgraph of  $K_n$  induced by a subset of  $m$  vertices.
- $P_n$  is *not* an induced subgraph of  $C_n$ .
- $Q_{n-1}$  is the subgraph of  $Q_n$  induced by all vertices whose labels starts with 0.

# Some operations on graphs

Let  $G = (V, E)$  be a simple graph.

## Edge removal:

- $G - e = (V, E \setminus \{e\})$ .
- Result: a subgraph of  $G$ .

## Vertex removal:

- $G - v = (V \setminus \{v\}, E \setminus \{e \mid e \text{ is incidence with } v\})$ .
- Result: a subgraph of  $G$ .

## Edge contraction:

- Remove an edge, then merge its endpoints.
- Result: *not* necessarily a subgraph!

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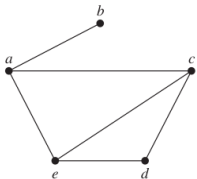
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# Adjacency list

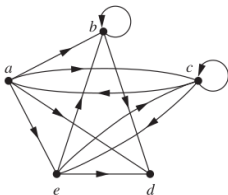


**TABLE 1** An Adjacency List  
for a Simple Graph.

<i>Vertex</i>	<i>Adjacent Vertices</i>
<i>a</i>	<i>b, c, e</i>
<i>b</i>	<i>a</i>
<i>c</i>	<i>a, d, e</i>
<i>d</i>	<i>c, e</i>
<i>e</i>	<i>a, c, d</i>

Figure: Rosen, p. 668

# Adjacency list



**TABLE 2** An Adjacency List for a Directed Graph.

<i>Initial Vertex</i>	<i>Terminal Vertices</i>
<i>a</i>	<i>b, c, d, e</i>
<i>b</i>	<i>b, d</i>
<i>c</i>	<i>a, c, e</i>
<i>d</i>	<i>b, c, d</i>
<i>e</i>	<i>b, c, d</i>

Figure: Rosen, p. 669

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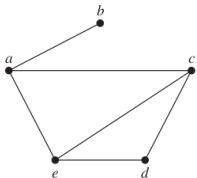
# Adjacency matrix

## Definition

The **adjacency matrix** of a graph  $G = (V, E)$  is a  $V \times V$  matrix of integer entries  $A = (a_{u,v})$  such that:

$$a_{u,v} = \begin{cases} 1 & \text{if } u, v \text{ are adjacent,} \\ 0 & \text{otherwise.} \end{cases}$$

**Example:** (with vertex ordering  $a, b, c, d, e$ )



$$A = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 \end{pmatrix}$$

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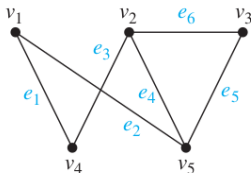
# Incidence matrix

## Definition

The **incidence matrix** of a graph  $G = (V, E)$  is a  $V \times E$  matrix of integer entries  $M = (m_{v,e})$  such that:

$$m_{v,e} = \begin{cases} 1 & \text{if } e \text{ is incident with } v, \\ 0 & \text{otherwise.} \end{cases}$$

## Example:



$$\begin{matrix} & e_1 & e_2 & e_3 & e_4 & e_5 & e_6 \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{matrix} & \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}.$$

**Note:** The notions of adjacency matrix and incidence matrix can be extended for multigraphs, for directed graphs, etc.