

The Art of Multiprocessor Programming
Solutions to Exercises
Chapter 8

July 14, 2009

```

1  public class BadCLHLock implements Lock {
2      // most recent lock holder
3      AtomicReference<Qnode> tail;
4      // thread-local variable
5      ThreadLocal<Qnode> myNode;
6      public void lock() {
7          Qnode qnode = myNode.get();
8          qnode.locked = true;           // I'm not done
9          // Make me the new tail, and find my predecessor
10         Qnode pred = tail.getAndSet(qnode);
11         // spin while predecessor holds lock
12         while (pred.locked) {}
13     }
14     public void unlock() {
15         // reuse my node next time
16         myNode.get().locked = false;
17     }
18     static class Qnode { // Queue node inner class
19         public boolean locked = false;
20     }
21 }

```

Figure 1: An incorrect attempt to implement a CLHLock.

Exercise 85. Fig. 1 shows an alternative implementation of CLHLock in which a thread reuses its own node instead of its predecessor node. Explain how this implementation can go wrong.

Solution Suppose *A* acquires and releases the lock. At this point, the same node is referenced by `tail` and `myNode`. If *A* tries to reacquire the lock, it sets the node's `locked` field to *true* (Line 8), swaps that node with itself (Line 10), implicitly making the node its own predecessor. It then deadlocks waiting for that node to become unlocked.