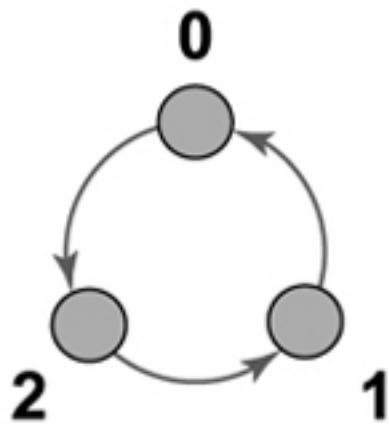
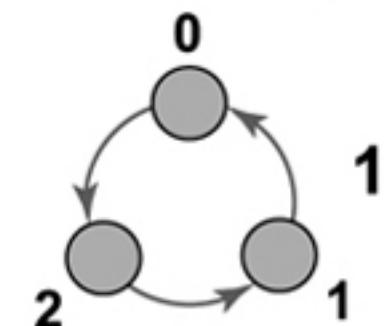


$T1 =$  $T2 =$  $Tk = T2 * Tk-1$  $T3 =$  $0$  $1$ 

**FIGURE 2.13** The precedence graph for a bounded timestamping system. Consider an initial situation in which there is a token *A* on node 12 (node 2 in subgraph 1) and tokens *B* and *C* on nodes 21 and 22 (nodes 1 and 2 in subgraph 2). Token *B* will move to node 20 to dominate the others. Token *C* will then move to node 21 to dominate the others, and *B* and *C* can continue to cycle in the  $T^2$  subgraph 2 forever. If *A* is to move to dominate *B* and *C*, it cannot pick a node in subgraph 2 since it is full (any  $T^k$  subgraph can accommodate at most  $k$  tokens). Instead, token *A* moves to node 00. If *B* now moves, it will choose node 01, *C* will choose node 10, and so on.