

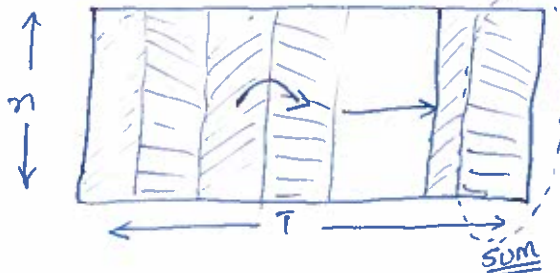


## (1) Computing likelihood $P(o_1, o_2, \dots, o_T)$

Efficient recursion

$$\begin{aligned}
 P(o_1, o_2, \dots, o_{t+1}, s_{t+1}=j) &= \sum_{i=1}^n P(o_1, o_2, \dots, o_{t+1}, s_t=i, s_{t+1}=j) \\
 &= \sum_{i=1}^n P(o_1, o_2, \dots, o_t, s_t=i) \cdot P(s_{t+1}=j | o_1, o_2, \dots, o_t, s_t=i) \quad \text{marginalization} \\
 &\quad \cdot P(o_{t+1} | o_1, \dots, o_t, s_t=i, s_{t+1}=j) \quad \text{Product rule} \\
 &= \sum_{i=1}^n \underbrace{P(o_1, o_2, \dots, o_t, s_t=i)}_{\text{recursive instance}} \cdot \underbrace{P(s_{t+1}=j | s_t=i)}_{\text{CPT}_S} \cdot \underbrace{P(o_{t+1} | s_{t+1}=j)}_{\text{C.I.}}
 \end{aligned}$$

Shorthand Notation:  $\alpha_{it} = P(o_1, o_2, \dots, o_t, s_t=i)$



$$\alpha_{j,t+1} = \sum_{i=1}^n \alpha_{it} \cdot a_{ij} \cdot b_j(o_{t+1})$$

FORWARD ALGORITHM IN HMMs

Base Case:  $\alpha_{i1} = P(o_1, s_1=i) = P(s_1=i) \cdot P(o_1 | s_1=i)$  P.R.

$$\Rightarrow \alpha_{i1} = \pi_i b_i(o_1)$$

$$\begin{aligned}
 * \text{Likelihood } P(o_1, o_2, \dots, o_T) &= \sum_{i=1}^n P(o_1, \dots, o_T, s_T=i) \quad \text{marginalization} \\
 &= \sum_{i=1}^n \alpha_{iT} \quad (\text{sum of the last column of matrix})
 \end{aligned}$$

\* warning: for long sequences, watch out for underflow...

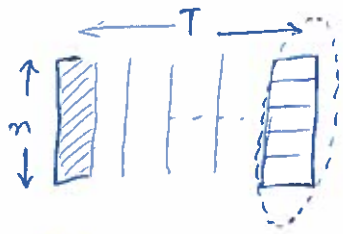
## (2) Computing most likely state sequence

$$\{s_1^*, s_2^*, \dots, s_T^*\} = \arg \max_{s_1, s_2, \dots, s_T} P(s_1, s_2, \dots, s_T | o_1, o_2, \dots, o_T)$$

$$= \arg \max_{s_1, s_2, \dots, s_T} \left[ \frac{P(s_1, s_2, \dots, s_T, o_1, o_2, \dots, o_T)}{P(o_1, o_2, \dots, o_T)} \right] \quad \text{Product rule}$$

$$= \arg \max_{s_1, s_2, \dots, s_T} \left[ P(s_1, s_2, \dots, s_T, o_1, o_2, \dots, o_T) \right] \quad \text{Constant w.r.t } s_1, s_2, \dots, s_T$$

Define:  $l_{i,t}^* = \max_{s_1, s_2, \dots, s_{t-1}} \log P(o_1, o_2, \dots, o_t, s_1, s_2, \dots, s_{t-1}, s_t = i)$



log-prob. of most likely "t-step path" of hidden states  $s_1, s_2, \dots, s_t$  that explains  $o_1, o_2, \dots, o_t$  and ends at  $s_t = i$  at time  $t$ .

\* Base Case: when  $t = 1$

$$l_{i,1}^* = \log P(o_1, s_1 = i) = \log [P(s_1 = i) P(o_1 | s_1 = i)] = \log \pi_i + \log b_i(o_1)$$

use this to fill first column of the matrix

\* Recursive Step from time step  $t$  to  $t+1$

$$l_{j,t+1}^* = \max_{s_1, s_2, \dots, s_t} \log P(o_1, o_2, \dots, o_{t+1}, s_1, s_2, \dots, s_t, s_{t+1} = j)$$

$$= \max_{s_1, s_2, \dots, s_{t-1}} \max_{\substack{i \\ \text{value of } s_t}} \log P(o_1, o_2, \dots, o_{t+1}, s_1, s_2, \dots, s_t = i, s_{t+1} = j)$$

$$= \max_{s_1, \dots, s_{t-1}} \max_i \log \left[ P(o_1, o_2, \dots, o_t, s_1, \dots, s_{t-1}, s_t = i) \cdot P(s_{t+1} = j | o_1, \dots, o_t, s_1, \dots, s_{t-1}, s_t = i) \cdot P(o_{t+1} | o_1, \dots, o_t, s_1, \dots, s_{t-1}, s_t = i, s_{t+1} = j) \right]$$

Product Rule

$$= \max_{s_1, \dots, s_{t-1}} \max_i \left\{ \log P(o_1, o_2, \dots, o_t, s_1, s_2, \dots, s_{t-1}, s_t = i) + \log P(s_{t+1} = j | s_t = i) + \log P(o_{t+1} | s_{t+1} = j) \right\}$$

$$l_{j,t+1}^* = \max_i \left[ \underbrace{\max_{s_1, s_2, \dots, s_{t-1}} \log P(o_1, \dots, o_t, s_1, \dots, s_{t-1}, s_t = i)}_{\text{recursive instance}} + \underbrace{\log P(s_{t+1} = j | s_t = i)}_{\text{CPT}} \right] + \underbrace{\log P(o_{t+1} | s_{t+1} = j)}_{\text{CPT}}$$

$$l_{j,t+1}^* = \max_i [l_{it}^* + \log a_{ij}] + \log b_j(o_{t+1})$$

(use this to recursively compute columns of matrix.)

- How to derive  $\{s_1^*, s_2^*, \dots, s_T^*\}$  from matrix  $\{l_{it}^*\}$  ?  
 $i=1 \text{ to } n$   
 $t=1 \text{ to } T$

\* Record most likely state transitions:

$$\Phi_{t+1}(j) = \arg \max_i [l_{it}^* + \log a_{ij}]$$

- what is the most likely state at time 't' given ' $s_{t+1}=j$ ' and we see  $\{o_1, o_2, \dots, o_{t+1}\}$  ?

\* Compute  $\{s_T^*, s_{T-1}^*, \dots, s_1^*\}$  by backtracking:

$$s_T^* = \arg \max_i [l_{iT}^*]$$

for  $t = T-1$  to  $1$ ,

$$s_t^* = \Phi_{t+1}(s_{t+1}^*)$$

•  $\{s_1^*, s_2^*, \dots, s_T^*\}$  called viterbi path

Viterbi algorithm: instance of dynamic programming.

### (3) Learning in HMMs

Given: sequence of observations  $\{o_1, o_2, \dots, o_T\}$   
 (just one for simplicity)

Goal: Estimate  $\{\pi_i, a_{ij}, b_{ik}\}$  to maximize  $P(o_1, o_2, \dots, o_T)$

Assume cardinality 'n' of hidden states is fixed  
 $s_t \in \{1, 2, \dots, n\}$

\* CPTs to estimate:

$$\pi_i = P(S_1 = i) \quad \text{root node}$$

$$a_{ij} = P(S_{t+1} = j \mid S_t = i) \quad \text{shared across time}$$

$$b_{ik} = P(O_t = k \mid S_t = i) \quad \text{shared across time}$$

Recall from EM algorithm:

to update  $P(X_i = x \mid P_{0i} = \pi)$ , need to compute  $P(X_i = x, P_{0i} = \pi \mid V)$

\* E-step for HMMs:

$$\text{Compute } P(S_1 = i \mid O_1, O_2, \dots, O_T)$$

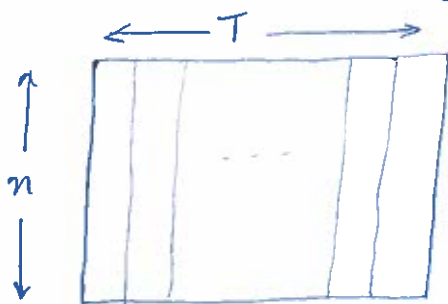
$$P(S_{t+1} = j, S_t = i \mid O_1, \dots, O_T)$$

$$P(O_t = k, S_t = i \mid O_1, \dots, O_T) = I(O_t, k) P(S_t = i \mid O_1, \dots, O_T)$$

\* How to compute these posteriors?

Analogous to  $\alpha_{it} = P(O_1, \dots, O_t, S_t = i)$  up to and including time  $t$

define  $\beta_{it} = P(\underline{O_{t+1}}, \underline{O_{t+2}}, \dots, \underline{O_T} \mid S_t = i)$  beyond time  $t$ .  
 Start at 't+1'      conditioning bar



\* Recursion for  $\beta_{it}$ :

(1) base case:  $t = T$

$$\beta_{iT} = P(\underline{\quad} \mid S_T = i) ? \quad \text{Set } \beta_{iT} = 1 \text{ for all } i$$

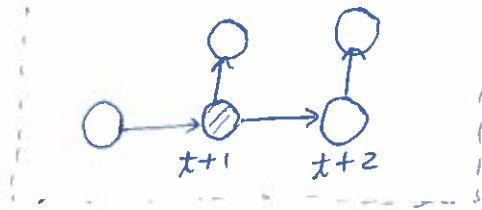
(2) Backwards step from time  $t+1$  to  $t$

$$\beta_{it} = P(o_{t+1}, \dots, o_T | s_t = i)$$

$$= \sum_{j=1}^n P(o_{t+1}, \dots, o_T, s_{t+1}=j | s_t=i) \quad \text{marginalization}$$

$$= \sum_{j=1}^n P(s_{t+1}=j | s_t=i) P(o_{t+1} | s_t=i, s_{t+1}=j) \cdot P(o_{t+2}, \dots, o_T | s_t=i, s_{t+1}=j, o_{t+1})$$

Product rule



$$= \sum_{j=1}^n P(s_{t+1}=j | s_t=i) \cdot P(o_{t+1} | s_{t+1}=j) \cdot P(o_{t+2}, \dots, o_T | s_{t+1}=j) \quad \text{C.I.}$$

$$\beta_{it} = \sum_{j=1}^n a_{ij} b_j(o_{t+1}) \beta_{j,t+1}$$