LET'S PLAY SOME WATER (NOT SO) FUN GAME!

Problem:

We study a simple game consisting in a grid of tiles, separated or not by walls, forming different groups of tiles. Numbers figure on the end of some lines/columns. They correspond to a number of flood tiles in this column/line. The goal of the game is to find a solution, consisting of a combination of tiles flooded satisfying the number on each end of line/column and natural rules of physic.

R_side_numbers:

- -Each line/column must contain the exact number of flooded tiles indicated on the end of the said line/column.
- -Water must follow the following rules:

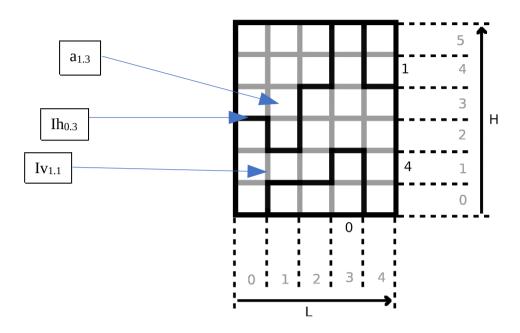
R_water_down:

A tile without water cannot be under a tile of water if there is no wall to stop said water.

R_water_level:

In the same group of tiles, if the group contains water, the level must be the same in all the group even if it is U shaped (or n shaped, for "simplicity").

Modelling:

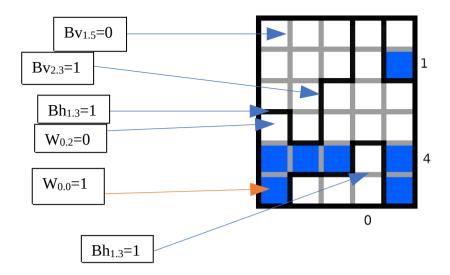


Definitions:

Let L being the horizontal size of the grid. Let H being the vertical size of the grid.

 $\forall 0 \le i < L, 0 \le j < H$ we name $a_{i,j}$ the tile of coordinates i,j $\forall 0 < i < L, 0 \le j < H$ we name $Iv_{i,j}$ the (vertical) intersection of the tiles $a_{i,j}$ and $a_{i-1,j}$ $\forall 0 \le i < L, 0 < j < H$ we name $Ih_{i,j}$ the (horizontal) intersection of the tiles $a_{i,j}$ and $a_{i,j-1}$ $\forall 0 \le i < L$ we name Zc_i the number aside to the column i. $\forall 0 < j < H$ we name Zl_j the number aside to the line j.

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\forall i,j,\ W_{i,j} = 1 <=> \text{ water in } a_{i,j}.
\forall i,j,\ Bh_{i,j} = 1 <=> \text{ horizontal barrier at } Ih_{i,j}
\forall i,j,\ Bv_{i,j} = 1 <=> \text{ vertical barrier at } Iv_{i,j}
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Modelling the "water's physics":

[schema]

We attribute a number to each tile such that each tile inside a same group of tiles have the same value, in order to be able to say if two tiles are in the same group (i.e. if there's a path between these tiles). For each tile $a_{i,j}$, we can access this number at $a_{i,j}$.val.

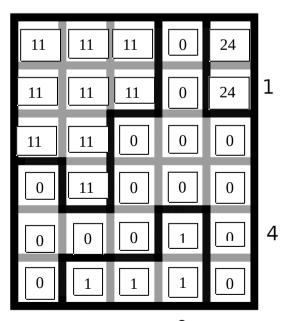
To attribute such values, a way to proceed is to initialize the values as the following:

∀i,j a_{i,j}.val=j*h+i

And then use the following rule to get the same values in each group:

 $\forall i,j \neg Bh_{i,j} \Rightarrow a_{i,j}.val = min(a_{i,j}.val; a_{i,j-1}.val)$

 $\forall i,j \neg Bv_{i,j} \Rightarrow a_{i,j}.val = min(a_{i,j}.val; a_{i-1,j}.val)$



Once a_{i,i}.val is defined, the following rules model the fact that water must follow the rules of physics:

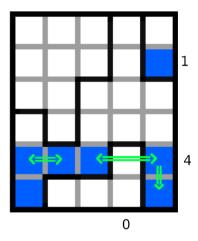
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R_water_down:

\forall i \text{ in } \{0,...,L-1\}, j \text{ in } \{0,...,H-1\}:

(ai, j.val=ai, j-1.val) \Rightarrow (Wi, j \Rightarrow Wi, j-1)
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R water level

$$\forall$$
i1,i2 in {0,...,L-1},j in {0,...,H-1}:
 $(ai1, j.val = ai2, j.val) \Rightarrow (Wi1, j \Leftrightarrow Wi2, j)$



Modelling the fact that Zc_i and Zl_j represents the number of tiles of water on the line or column (application of R_side_numbers):

Since there happens the same thing on each line with ZI_i and each column with Zc_j , let's say that S is either L or H, and N is the other.

We will then define Z and Wz_i as follows:

if S = H:

for any 0≤i<S, Z is Zci

for any $0 \le j \le N$, Wz_j is $W_{i,j}$ (representing the fact of water being on the tiles of the same column as Z) if S = L:

for any 0≤i<S, Z is Zl_i

for any 0≤j<N, Wz_j is W_{j,i} (representing the fact of water being on the tiles of the same line as Z)

Since we want exactly Z tiles to be filled of water, a way to model it is by making a disjunction of each model, but transforming it to a conjunction normal form by developing is long and grindy. Instead, we can model the fact that any assignment with at least Z+1 tiles filled of water should be unsatisfiable, and that any assignment with at most Z-1 tiles filled of water should be unsatisfiable.

To model the fact that any assignment with at most Z-1 tiles filled of water should be unsatisfiable, we do a conjunction of all possible clauses containing exactly N-(Z-1) elements of the form Wz_i . (like this, any assignment of at most Z-1 elements that are true have a clause that they can't validate). If Z = 0, no need to add a formula for this rule. Otherwise, the formula is the following:

$$\prod_{x_0=0}^{Z-1} \big(\prod_{x_1=x_0}^{Z} \prod_{x_2=x_1}^{Z+1} \dots \prod_{x_{N-Z}=x_{N-Z-1}}^{N-1} \big(\sum_{i \in \{x_0, \dots, x_{N-Z}\}} W z_i \big) \big)$$

(The product symbolizes a succession of conjunctions and the sum, a succession of disjunctions.)

To model the fact that any assignment with at least Z+1 tiles filled of water should be unsatisfiable, it is the same reasoning, but with a conjunction of all possible clauses containing exactly Z+1 elements of the form $\neg Wz_i$ (note the negation, and so it is the same as verifying that at most N-Z elements of the form Wz_i are false, which is directly equivalent to saying that there is at most Z tiles filled)

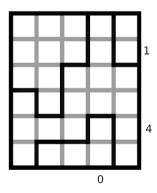
If Z = N, no need to add a formula for this rule. Otherwise, the formula is the following:

$$\prod_{x_0=0}^{N-Z-1} \left(\prod_{x_1=x_0}^{N-Z} \prod_{x_2=x_1}^{N-Z+1} \dots \prod_{x_z=x_{z-1}}^{N-1} \left(\sum_{i \in \{x_0,\dots,x_z\}} \neg W z_i \right) \right)$$

And so, the conjunction of these two formulas models the fact that exactly Z of the Wz_i are true, and the other false.

Example:

Here we take the following grid:



In this example we need a model satisfying all the above-mentioned rules:

R_water_down:

$$(\neg W_{0.5} V W_{0.4}) \Lambda (\neg W_{0.4} V W_{0.3}) \Lambda (\neg W_{0.2} V W_{0.1}) \Lambda (\neg W_{0.1} V W_{0.0})$$

$$\Lambda(\neg W_{1.5} \ V \ W_{1.4}) \Lambda(\neg W_{1.4} \ V \ W_{1.3}) \Lambda(\neg W_{1.3} \ V \ W_{1.2})$$

$$\Lambda(\neg W_{2.5} V W_{2.4})\Lambda(\neg W_{2.3} V W_{2.2})\Lambda(\neg W_{2.2} V W_{2.1})$$

$$\Lambda(\neg W_{3.5} V W_{3.4})\Lambda(\neg W_{3.4} V W_{3.3})\Lambda(\neg W_{3.3} V W_{3.2})\Lambda(\neg W_{3.1} V W_{3.0})$$

$$\Lambda(\neg W_{4.5} V W_{4.4})\Lambda(\neg W_{4.3} V W_{4.2})\Lambda(\neg W_{4.2} V W_{4.1})\Lambda(\neg W_{4.1} V W_{4.0})$$

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R water level:
(\neg W_{0.5} \ V \ W_{1.5}) \land (W_{0.5} \ V \ \neg W_{1.5}) \land (\neg W_{1.5} \ V \ W_{2.5}) \land (W_{1.5} \ V \ \neg W_{2.5})
\Lambda( \neg W_{0.4} \ V \ W_{1.4}) \Lambda(W_{0.4} \ V \ \neg W_{1.4}) \Lambda( \neg W_{1.4} \ V \ W_{2.4}) \Lambda(W_{1.4} \ V \ \neg W_{2.4})
\neg W_{4.3})
\Lambda(\neg W_{2.2} \ V \ W_{3.2})\Lambda(W_{2.2} \ V \ \neg W_{3.2})\Lambda(\neg W_{3.2} \ V \ W_{4.2})\Lambda(W_{3.2} \ V \ \neg W_{4.2})
\Lambda(\neg W_{0.1} \ V \ W_{1.1})\Lambda(W_{0.1} \ V \ \neg W_{1.1})\Lambda(\neg W_{1.1} \ V \ W_{2.1})\Lambda(W_{1.1} \ V \ \neg W_{2.1})
\Lambda( \neg W_{1.0} \ V \ W_{2.0}) \Lambda(W_{1.0} \ V \ \neg W_{2.0}) \Lambda( \neg W_{2.0} \ V \ W_{3.0}) \Lambda(W_{2.0} \ V \ \neg W_{3.0})
R side numbers:
For the line 1 (Zl_1 = 4):
S = 6, Z = Zl_1 = 4, N = 5
N-Z = 1
(W_{0,1} \vee W_{1,1}) \wedge (W_{0,1} \vee W_{2,1}) \wedge (W_{0,1} \vee W_{3,1}) \wedge (W_{0,1} \vee W_{4,1})
\Lambda(W_{1,1} \vee W_{2,1})\Lambda(W_{1,1} \vee W_{3,1})\Lambda(W_{1,1} \vee W_{4,1})
\Lambda(W_{2,1} \vee W_{3,1})\Lambda(W_{2,1} \vee W_{4,1})
\Lambda(W_{3.1}V W_{4.1})
\Lambda (\neg W_{0,1} \lor \neg W_{1,1} \lor \neg W_{2,1} \lor \neg W_{3,1} \lor \neg W_{4,1})
For the line 4 (Zl_4 = 1):
S = 6, Z = Zl_4 = 1, N = 5
N-Z=4
(W_{0,4} \vee W_{1,4} \vee W_{2,4} \vee W_{3,4} \vee W_{4,4})
\Lambda(\neg W_{0,4} \lor \neg W_{1,4})\Lambda(\neg W_{0,1} \lor \neg W_{2,4}) \Lambda(\neg W_{0,1} \lor \neg W_{3,4}) \Lambda(\neg W_{0,1} \lor \neg W_{4,4})
\Lambda(\neg W_{1,4} \lor \neg W_{2,4}) \Lambda(\neg W_{1,1} \lor \neg W_{3,4}) \Lambda(\neg W_{1,1} \lor \neg W_{4,4})
\Lambda(\neg W_{2,4} \vee \neg W_{3,4}) \Lambda(\neg W_{2,1} \vee \neg W_{4,4})
\Lambda(\neg W_{3,4} \vee \neg W_{4,4})
For the column 3 (Zc_3 = 0):
S = 5, Z = Zc_3 = 0, N = 6
N-Z=6
 \neg W_{3,0} \land \neg W_{3,1} \land \neg W_{3,2} \land \neg W_{3,3} \land \neg W_{3,4} \land \neg W_{3,5}
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