

Chapter 5: Relations

CS1231 Discrete Structures

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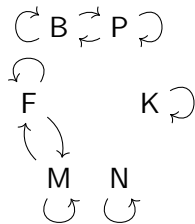
2022/23 Semester 2

Professionals in any discipline need to know the foundations of their field. So if you're a database professional, you need to know the relational model, because the relational model is the foundation (or a huge part of the foundation, at any rate) of the database field in particular. [Date \(2009\)](#)

Plan

Start using sets to represent mathematical objects.

- ▶ relations representing *predicates*
 - ordered tuples and Cartesian products
 - composition
 - inverse
- ▶ graphs representing *graphs*
 - directed graphs
 - undirected graphs



How one can represent a predicate by a set

belongs to	
Student ID	name
001R	Gates
012B	Brin
062E	Bezos
126N	Ma
254E	Zuckerberg

{ (001R, Gates),
(012B, Brin),
(062E, Bezos),
(126N, Ma),
(254E, Zuckerberg) }

Ordered pairs and Cartesian products

Definition 5.1.1

An *ordered pair* is an expression of the form (x, y) . Let (x_1, y_1) and (x_2, y_2) be ordered pairs. Then $(x_1, y_1) = (x_2, y_2)$ if

$$x_1 = x_2 \quad \text{and} \quad y_1 = y_2.$$

Example 5.1.2

(1) $(1, 2) \neq (2, 1)$, although $\{1, 2\} = \{2, 1\}$.

(2) $(3, 0.5) = (\sqrt{9}, \frac{1}{2})$.

read as "*A cross B*"

Definition 5.1.3

Let A, B be sets. The *Cartesian product* of A and B , denoted $A \times B$, is defined to be

$$\{(x, y) : x \in A \text{ and } y \in B\}.$$

Example 5.1.4

$$\{a, b\} \times \{1, 2, 3\} = \{(a, 1), (a, 2), (a, 3), (b, 1), (b, 2), (b, 3)\}.$$

$(a, 3) - (b, 3)$

| |

$(a, 2) - (b, 2)$

| |

$(a, 1) - (b, 1)$

Predicates as sets

Definition 5.1.5

Let A, B be sets.

- (1) A *relation* from A to B is a subset of $A \times B$.
- (2) Let R be a relation from A to B and $(x, y) \in A \times B$. Then we may write

$$x R y \text{ for } (x, y) \in R \quad \text{and} \quad x \not R y \text{ for } (x, y) \notin R.$$

We read “ $x R y$ ” as “ x is *R -related* to y ” or simply “ x is *related* to y ”.

Example 5.1.6

Let $\Gamma = \{A, B, \dots, Z, 0, 1, 2, \dots, 9\}$ and $\Phi = \{A, B, \dots, Z, a, b, \dots, z\}$. Then

$$\left\{ \begin{array}{l} (001R, \text{ Gates} \quad \quad), \\ (012B, \text{ Brin} \quad \quad), \\ (062E, \text{ Bezos} \quad \quad), \\ (126N, \text{ Ma} \quad \quad), \\ (254E, \text{ Zuckerberg}) \end{array} \right\}$$

is a relation from Γ^* to Φ^* .

Arrow diagrams for binary relations

Example 5.1.7

Let $A = \{0, 1, 2\}$ and $B = \{1, 2, 3, 4\}$. Define the relation R from A to B by setting

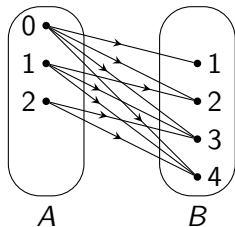
$$x R y \iff x < y.$$

Then $0 R 1$ and $0 R 2$, but $2 \not R 1$. Thus

$$R = \{(0, 1), (0, 2), (0, 3), (0, 4), \\ (1, 2), (1, 3), (1, 4), (2, 3), (2, 4)\}.$$

Definition 5.1.8 (arrow diagrams)

An arrow from x to y indicates x is related to y .



How one can represent a predicate of higher arity by a set

teaching			
module	department	faculty	instructor
CS3234	CS	Computing	Turing
MA2001	Mathematics	Science	Gauss
MU2109	Music	Arts	Mozart
PC2130	Physics	Science	Newton
PL3103	Psychology	Arts	Freud

{ (CS3234, CS, Computing, Turing),
 (MA2001, Mathematics, Science, Gauss),
 (MU2109, Music, Arts, Mozart),
 (PC2130, Physics, Science, Newton),
 (PL3103, Psychology, Arts, Freud) }

Ordered n -tuples and Cartesian products

Let $n \in \mathbb{Z}_{\geq 2}$.

Definition 5.1.10

- (1) An expression of the form (x_1, x_2, \dots, x_n) is called an *ordered n -tuple*.
- (2) Ordered n -tuples are defined recursively by setting, for all objects x_1, x_2, \dots, x_{n+1} ,

$$(x_1, x_2, \dots, x_{n+1}) = ((x_1, x_2, \dots, x_n), x_{n+1}).$$

Exercise 5.1.11

Prove by induction on n that for all (x_1, x_2, \dots, x_n) and (y_1, y_2, \dots, y_n) ,

 5a

$$(x_1, x_2, \dots, x_n) = (y_1, y_2, \dots, y_n) \iff x_1 = y_1 \text{ and } x_2 = y_2 \text{ and } \dots \text{ and } x_n = y_n.$$

Definition 5.1.13

Let A_1, A_2, \dots, A_n be sets. The *Cartesian product* of A_1, A_2, \dots, A_n , denoted $A_1 \times A_2 \times \dots \times A_n$, is defined to be

$$\{(x_1, x_2, \dots, x_n) : x_1 \in A_1 \text{ and } x_2 \in A_2 \text{ and } \dots \text{ and } x_n \in A_n\}.$$

If A is a set, then $A^n = \underbrace{A \times A \times \dots \times A}_{n\text{-many } A\text{'s}}$.

Example 5.1.14

$$\{0, 1\} \times \{0, 1\} \times \{a, b\} = \{(0, 0, a), (0, 0, b), (0, 1, a), (0, 1, b), (1, 0, a), (1, 0, b), (1, 1, a), (1, 1, b)\}.$$

Let A_1, A_2, \dots, A_n be sets.

Definition 5.1.15

A *n -ary relation over A_1, A_2, \dots, A_n* is a subset of $A_1 \times A_2 \times \dots \times A_n$.

Example 5.1.16

Let $\Gamma = \{A, B, \dots, Z, 0, 1, 2, \dots, 9\}$ and $\Phi = \{A, B, \dots, Z, a, b, \dots, z\}$. Then

$\{(CS3234, CS, Computing, Turing), (MA2001, Mathematics, Science, Gauss),$
 $(MU2109, Music, Arts, Mozart), (PC2130, Physics, Science, Newton),$
 $(PL3103, Psychology, Arts, Freud)\}.$

is a 4-ary relation over $\Gamma^*, \Phi^*, \Phi^*, \Phi^*$.

A fictitious miniature university database and its representation (Figure 5.1)

identity	
Student ID	name
001R	Gates
012B	Brin
062E	Bezos
126N	Ma
254E	Zuckerberg

$$SN = \{ (001R, \text{Gates}), (012B, \text{Brin}), (062E, \text{Bezos}), (126N, \text{Ma}), (254E, \text{Zuckerberg}) \}$$

progress		
Student ID	faculty	year
062E	Arts	1
254E	Arts	2
012B	Science	2
001R	Science	1
126N	Science	3

teaching			
module	department	faculty	instructor
CS3234	CS	Computing	Turing
MA2001	Mathematics	Science	Gauss
MU2109	Music	Arts	Mozart
PC2130	Physics	Science	Newton
PL3103	Psychology	Arts	Freud

is enrolled in	
Student ID	module
126N	CS3234
254E	CS3234
001R	MA2001
012B	MA2001
062E	MA2001
126N	MA2001
012B	MU2109
001R	PC2130
062E	PL3103
254E	PL3103

$$SM = \{ (126N, \text{CS3234}), (254E, \text{CS3234}), (001R, \text{MA2001}), (012B, \text{MA2001}), (062E, \text{MA2001}), (126N, \text{MA2001}), (012B, \text{MU2109}), (001R, \text{PC2130}), (062E, \text{PL3103}), (254E, \text{PL3103}) \}$$

$$SFY = \{ (062E, \text{Arts}, 1), (254E, \text{Arts}, 2), (012B, \text{Science}, 2), (001R, \text{Science}, 1), (126N, \text{Science}, 3) \}$$

$$MDFI = \{ (\text{CS3234}, \text{CS}, \text{Computing}, \text{Turing}), (\text{MA2001}, \text{Mathematics}, \text{Science}, \text{Gauss}), (\text{MU2109}, \text{Music}, \text{Arts}, \text{Mozart}), (\text{PC2130}, \text{Physics}, \text{Science}, \text{Newton}), (\text{PL3103}, \text{Psychology}, \text{Arts}, \text{Freud}) \}$$

The set $\{SM, SN, SFY, MDFI\}$ represents the relational database.

Composition: the teacher of the father of ...

Definition 5.2.1

Let R be a relation from A to B , and S be a relation from B to C . Then $S \circ R$ is the relation from A to C defined by

$$S \circ R = \{(x, z) \in A \times C : (x, y) \in R \text{ and } (y, z) \in S \text{ for some } y \in B\}.$$

We read $S \circ R$ as “ S composed with R ” or “ S circle R ”.

Example 5.2.4

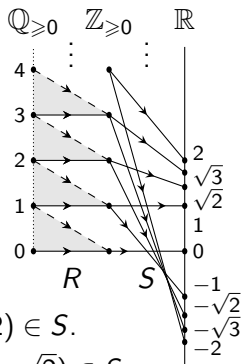
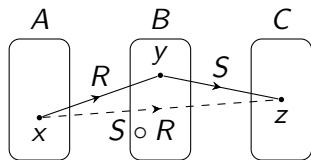
Define relations R from $\mathbb{Q}_{\geq 0}$ to $\mathbb{Z}_{\geq 0}$ and S from $\mathbb{Z}_{\geq 0}$ to \mathbb{R} by:

$$R = \{(x, y) \in \mathbb{Q}_{\geq 0} \times \mathbb{Z}_{\geq 0} : \lfloor x \rfloor = y\},$$

$$S = \{(y, z) \in \mathbb{Z}_{\geq 0} \times \mathbb{R} : y = z^2\}.$$

- ▶ $(4.8, 2) \in S \circ R$ because $4 \in \mathbb{Z}_{\geq 0}$ such that $(4.8, 4) \in R$ and $(4, 2) \in S$.
- ▶ $(5/2, -\sqrt{2}) \in S \circ R$ because $2 \in \mathbb{Z}_{\geq 0}$ such that $(5/2, 2) \in R$ and $(2, -\sqrt{2}) \in S$.

In general, we have $S \circ R = \{(x, z) \in \mathbb{Q}_{\geq 0} \times \mathbb{R} : \lfloor x \rfloor = z^2\}.$

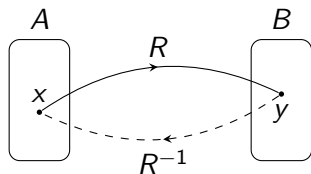


Inverse: from parents to children, and back

Definition 5.2.5

Let R be a relation from A to B . Then the *inverse of R* is the relation R^{-1} from B to A defined by

$$R^{-1} = \{(y, x) \in B \times A : (x, y) \in R\}.$$



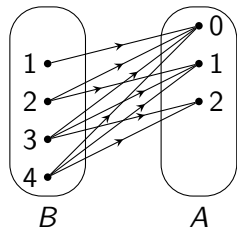
Example 5.2.6

Let R be a relation from A to B where

$$A = \{0, 1, 2\}, \quad B = \{1, 2, 3, 4\},$$

$$R = \{(0, 1), (0, 2), (0, 3), (0, 4), (1, 2), (1, 3), (1, 4), (2, 3), (2, 4)\}.$$

Then $R^{-1} = \{(1, 0), (2, 0), (3, 0), (4, 0), (2, 1), (3, 1), (4, 1), (3, 2), (4, 2)\}.$



Composition interacts with inversion

Proposition 5.2.7

Let R be a relation from A to B , and S be a relation from B to C . Then $(S \circ R)^{-1} = R^{-1} \circ S^{-1}$.

Proof

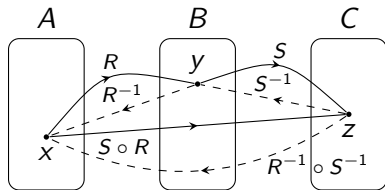
Both $(S \circ R)^{-1}$ and $R^{-1} \circ S^{-1}$ are relations from C to A . For all $(z, x) \in C \times A$,

$$\begin{aligned} (z, x) \in (S \circ R)^{-1} &\Leftrightarrow (x, z) \in S \circ R && \text{by the definition of inverses;} \\ \Leftrightarrow (x, y) \in R \text{ and } (y, z) \in S \text{ for some } y \in B &&& \text{by the definition of composition;} \\ \Leftrightarrow (y, x) \in R^{-1} \text{ and } (z, y) \in S^{-1} \text{ for some } y \in B &&& \text{by the definition of inverses;} \\ \Leftrightarrow (z, x) \in R^{-1} \circ S^{-1} &&& \text{by the definition of composition.} \end{aligned}$$

So $(S \circ R)^{-1} = R^{-1} \circ S^{-1}$. □

Real-life example. Peter is Mary's parent's elder sibling.

\Leftrightarrow Mary is Peter's younger sibling's child.



Commutativity of relation composition

Exercise 5.2.8

Let $A = \{0, 1, 2\}$. Define two relations R, S from A to A by:

$$R = \{(x, y) \in A^2 : x < y\} \quad \text{and} \quad S = \{(0, 1), (1, 2), (2, 0)\}.$$

Is $R \circ S = S \circ R$? Prove that your answer is correct.



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Summary

Let A, B, C be sets.

Definition 5.1.5. R is a *relation* from A to $B \iff R \subseteq A \times B$.

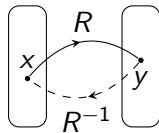
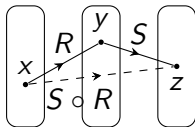
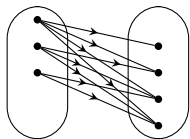
Sometimes we write $x R y$ for $(x, y) \in R$.

Let R be a relation from A to B and S be a relation from B to C .

Definitions 5.2.1 and 5.2.5. For all $x \in A$, all $y \in B$ and all $z \in C$,

$$\begin{array}{lll} (x, z) \in S \circ R & \iff & \exists y \in B ((x, y) \in R \wedge (y, z) \in S) \\ \text{or} & & \\ x (S \circ R) z & \iff & \exists y \in B (x R y \wedge y S z); \\ (y, x) \in R^{-1} & \iff & (x, y) \in R \\ \text{or} & & \\ y R^{-1} x & \iff & x R y. \end{array}$$

Proposition 5.2.7. $(S \circ R)^{-1} = R^{-1} \circ S^{-1}$.



Representing directed graphs using sets

Definition 5.3.1 and Remark 5.3.2

A *(binary) relation on a set A* is a relation from A to A , i.e., a subset of $A \times A$.

Definition 5.3.5

A *directed graph* is an ordered pair (V, D) where V is a set and D is a binary relation on V . In the case when (V, D) is a directed graph,

- (1) the *vertices* or the *nodes* are the elements of V ;
- (2) the *edges* are the elements of D ;
- (3) an edge *from x to y* is the element $(x, y) \in D$;
- (4) a *loop* is an edge from a vertex to itself;
- (5) a *drawing* of (V, D) is an arrow diagram for D as a relation on A .

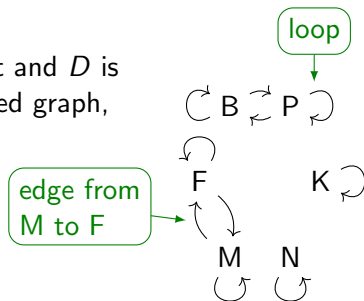
Examples 5.3.6 and 5.3.4

Let $V = \{B, P, F, M, K, N\}$, and

$$D = \{(B, P), (P, B), (F, M), (M, F), (B, B), (P, P), (F, F), (M, M), (K, K), (N, N)\}.$$

Then (V, D) is a directed graph. The figure above is a drawing of (V, D) .

Arrow diagram. An arrow from x to y indicates x is related to y .



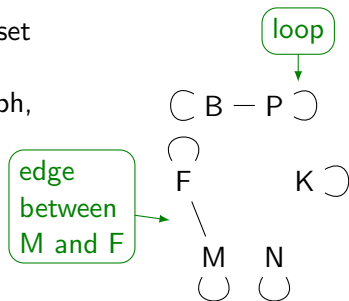
Representing undirected graphs using sets

Drawing. A line between x and y indicates an edge between x and y .

Definition 5.3.7

An **undirected graph** is an ordered pair (V, E) where V is a set and E is a set all of whose elements are of the form $\{x, y\}$ with $x, y \in V$. In the case when (V, E) is an undirected graph,

- (1) the **vertices** or the **nodes** are the elements of V ;
- (2) the **edges** are the elements of E ;
- (3) an edge **between x and y** is the element $\{x, y\} \in E$;
- (4) a **loop** is an edge between a vertex and itself.



Examples 5.3.6 and 5.3.4

Let $V = \{B, P, F, M, K, N\}$, and

$$E = \{\{B, P\}, \{P, B\}, \{F, M\}, \{M, F\}, \{B, B\}, \{P, P\}, \{F, F\}, \{M, M\}, \{K, K\}, \{N, N\}\}.$$

Then (V, E) is an undirected graph. The figure above is a drawing of (V, E) .