### 10.2 1s Complement (2/3)

Technique to negate a value: invert all the bits.

Largest value: 01111111 = +127<sub>10</sub>

Smallest value: 10000000 = -127<sub>10</sub>

 $\blacksquare$  Zeros:  $00000000 = +0_{10}$ 

 $111111111 = -0_{10}$ 

- Range (for 8 bits): -127<sub>10</sub> to +127<sub>10</sub>
- Range (for *n* bits):  $-(2^{n-1}-1)$  to  $2^{n-1}-1$
- The most significant bit (MSB) still represents the sign: 0 for positive, 1 for negative.

# 10.3 2s Complement (2/3)

 Technique to negate a value: invert all the bits, then add 1.

Largest value: 01111111 = +127<sub>10</sub>

Smallest value: 10000000 = -128<sub>10</sub>

**Zero:**  $00000000 = +0_{10}$ 

- Range (for 8 bits): -128<sub>10</sub> to +127<sub>10</sub>
- Range (for n bits): -2<sup>n-1</sup> to 2<sup>n-1</sup> 1
- The most significant bit (MSB) still represents the sign: 0 for positive, 1 for negative.

# 10.8 Excess Representation (1/2)

- Besides sign-and-magnitude and complement schemes, the excess representation is another scheme.
- It allows the range of values to be distributed <u>evenly</u> between the positive and negative values, by a simple translation (addition/subtraction).
- Example: Excess-4 representation on 3-bit numbers. See table on the right.

Excess-4 Representation	Value
000	-4
001	-3
010	-2
011	-1
100	0
101	1
110	2
111	3



#### 11.2 IEEE 754 Floating-Point Rep. (3/4)

3 components: sign, exponent and mantissa (fraction)

sign
------

- Sign bit: 0 for positive, 1 for negative.
- Mantissa is normalised with an implicit leading bit 1
  - 110.1<sub>2</sub>  $\rightarrow$  normalised  $\rightarrow$  1.101<sub>2</sub> × 2<sup>2</sup>  $\rightarrow$  only **101** is stored in the mantissa field
  - $0.00101101_2 \rightarrow \text{normalised} \rightarrow 1.01101_2 \times 2^{-3} \rightarrow \text{only } 01101 \text{ is stored in the mantissa field}$



#### 11.2 IEEE 754 Floating-Point Rep. (2/4)

3 components: sign, exponent and mantissa (fraction)

sign	exponent	mantissa
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- The base (radix) is assumed to be 2.
- Two formats:
  - Single-precision (32 bits): 1-bit sign, 8-bit exponent with bias 127 (excess-127), 23-bit mantissa
  - Double-precision (64 bits): 1-bit sign, 11-bit exponent with bias 1023 (excess-1023), and 52-bit mantissa
- We will focus on the single-precision format
- Reading
  - DLD pages 32 33
  - IEEE standard 754 floating point numbers:
    <a href="http://steve.hollasch.net/cgindex/coding/ieeefloat.html">http://steve.hollasch.net/cgindex/coding/ieeefloat.html</a>



#### 11.2 IEEE 754 Floating-Point Rep. (3/4)

3 components: sign, exponent and mantissa (fraction)

sign
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  - $0.00101101_2 \rightarrow \text{normalised} \rightarrow 1.01101_2 \times 2^{-3} \rightarrow \text{only } 01101 \text{ is stored in the mantissa field}$



### **Function Prototype**

- It is a good practice to put function prototypes at the top of the program, <u>before</u> the main() function, to inform the compiler of the functions that your program may use and their return types and parameter types.
- A function prototype includes only the function's return type, the function's name, and the data types of the parameters (names of parameters are optional).
- Function definitions to follow <u>after</u> the main() function.
- Without function prototypes, you will get error/warning messages from the compiler.

