

Duality in Linear Programming:

Given a LPP in canonical form

$$\text{Maximize } Z = \sum_{j=1}^n c_j x_j$$

$$\text{Subject to } \sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i = 1, 2, \dots, m \quad \text{_____} \quad (I)$$
$$x_j \geq 0, \quad j = 1, 2, \dots, n$$

We can write its corresponding dual problem as

$$\text{Minimize } W = \sum_{i=1}^m b_i y_i$$

$$\text{Subject to } \sum_{i=1}^m a_{ij} y_i \geq c_j, \quad j = 1, 2, \dots, n \quad \text{_____} \quad (II)$$
$$y_i \geq 0, \quad i = 1, 2, \dots, m$$

The given problem (I) is called Primal problem

And problem (II) is called its Dual problem

Thus problem (I) and (II) are dual of each other.

Note:

1. If given primal problem contains large numbers of constraints than numbers of variables then computational steps can be reduced by converting it into dual and then solving it.
2. It gives additional information as to how the optimal solution changes as a result of the changes in the coefficients formulation of the problem known as sensitivity analysis.
3. An unrestricted variable in sign in primal problem will result in an equality constrain in dual problem and vice versa.

Correspondence between Primal and Dual Optimal solution:

If the dual problem has optimal solution then primal problem also has optimal solution.

The value under the slack variables (if any neglecting the negative sign) and under the artificial variables (if any and neglecting the negative sign and deleting constant M) in the net evaluation row $(c_j - z_j)$ of the optimal table of the dual problem gives the values of basic variables of the primal problem and vice versa. The optimal value of dual is equal to optimal value of primal and vice versa i.e $Z_{\max} = W_{\min}$

Dual LPP :-

E-x. Write the dual of the following problems.

(1)
$$\begin{aligned} \text{Max } Z &= 5x_1 - 4x_2 + 4x_3 \\ \text{sub to } 6x_1 + 5x_2 + 10x_3 &\leq 7 \\ 2x_1 + x_2 - 6x_3 &= 20 \\ 8x_1 - 3x_2 + 6x_3 &\geq 50 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$

Soln

canonical form is

$$\text{Max } Z = 5x_1 - 4x_2 + 4x_3$$

$$\begin{aligned} \text{sub to } 6x_1 + 5x_2 + 10x_3 &\leq 7 \\ 2x_1 + x_2 - 6x_3 &\leq 20 \\ -2x_1 - x_2 + 6x_3 &\leq -20 \\ -8x_1 + 3x_2 + 6x_3 &\leq -50 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$

\therefore Dual problem is

$$\text{Min } W = 7y_1 + 20y_2' - 20y_2'' - 50y_3$$

$$\begin{aligned} \text{sub to } 6y_1 + 2y_2' - 2y_2'' - 8y_3 &\geq 5 \\ 5y_1 + y_2' - y_2'' + 3y_3 &\geq -4 \\ 10y_1 - 6y_2' + 6y_2'' - 6y_3 &\geq 4 \\ y_1, y_2', y_2'', y_3 &\geq 0 \end{aligned}$$

put $y_2' - y_2'' = y_2$;

$$\text{Min } W = 7y_1 + 20y_2 - 50y_3$$

$$\text{sub to } 6y_1 + 2y_2 - 8y_3 \geq 5$$

$$5y_1 + y_2 + 3y_3 \geq -4$$

$$10y_1 - 6y_2 - 6y_3 \geq 4$$

$$y_1, y_3 \geq 0$$

and y_2 is unrestricted in sign.

E.x. solve the LPP using Dual problem.

(1)

$$\text{Max } Z = 5x_1 - 2x_2 + 3x_3$$

$$\text{sub to } 2x_1 + 2x_2 - x_3 \geq 2$$

$$3x_1 - 4x_2 \leq 3$$

$$x_2 + 3x_3 \leq 5$$

$$x_1, x_2, x_3 \geq 0$$

soln

canonical form:-

$$\text{Max } Z = 5x_1 - 2x_2 + 3x_3$$

$$\text{sub to } -2x_1 - 2x_2 + x_3 \leq -2$$

$$3x_1 - 4x_2 \leq 3$$

$$x_2 + 3x_3 \leq 5$$

$$x_1, x_2, x_3 \geq 0$$

∴ The corresponding dual problem is

$$\text{Min } W = -2y_1 + 3y_2 + 5y_3$$

$$\text{sub to } -2y_1 + 3y_2 \geq 5$$

$$-2y_1 - 4y_2 + y_3 \geq -2$$

$$y_1 + 3y_3 \geq 3$$

$$y_1, y_2, y_3 \geq 0$$

Introducing the slacks and Artificial variables; standard form is

$$\text{Min } W = -2y_1 + 3y_2 + 5y_3 + MA_1 + MA_2$$

$$\text{sub to } -2y_1 + 3y_2 - s_1 + A_1 = 5$$

$$2y_1 + 4y_2 - y_3 + s_2 = 2$$

$$y_1 + 3y_3 - s_3 + A_2 = 3$$

$$y_1, y_2, y_3, s_1, s_2, s_3, A_1, A_2 \geq 0$$

$$\text{put } y_1 = 0, y_2 = 0, y_3 = 0, s_1 = 0, s_3 = 0$$

$$\Rightarrow A_1 = 5, s_2 = 2, A_2 = 3$$

C_j CB Basis		-2	3	5	0	0	0	M	M	So/h	Min Ratio
		y_1	y_2	y_3	s_1	s_2	s_3	A_1	A_2		
M	A_1	-2	3	0	-1	0	0	1	0	5	$5/3$
0	s_2	2	(4)	-1	0	1	0	0	0	2	$1/2 \rightarrow$
M	A_2	1	0	3	0	0	-1	0	1	3	∞
Z_j		-M	3M	3M	-M	0	-M	M	M		
$C_j - Z_j$		$2+M$	$3-3M$	$5-3M$	M	0	M	0	0		
		\uparrow									
M	A_1	$-7/2$	0	$3/4$	-1	$-3/4$	0	1	0	$7/2$	$14/3$
3	y_2	$1/2$	1	$-1/4$	0	$1/4$	0	0	0	$1/2$	-2
M	A_2	1	0	(3)	0	0	-1	0	1	3	$1 \rightarrow$
Z_j		$\frac{3}{2} - \frac{5}{2}M$	3	$\frac{-3}{4} + \frac{15}{4}M$	-M	$\frac{3}{4} - \frac{3}{4}M$	-M	M	M		
$C_j - Z_j$		$\frac{7}{2} + \frac{5}{2}M$	0	$\frac{23}{4} - \frac{15}{4}M$	M	$\frac{-3}{4} + \frac{3}{4}M$	M	0	0		
		\uparrow									
M	A_1	-2	0	0	-1	$\frac{-3}{4} (\frac{1}{4})$	1	$-\frac{1}{4}$		$\frac{11}{4}$	$11 \rightarrow$
3	y_2	$7/12$	1	0	0	$1/4$	$-1/12$	0	$1/12$	$3/4$	-9
5	y_3	$1/3$	0	1	0	0	$-1/3$	0	$1/3$	1	-3
Z_j		$\frac{41}{12} - 2M$	3	5	-M	$\frac{3}{4} - \frac{3}{4}M$	$\frac{-23}{12} + \frac{M}{4}$	$\frac{23}{12} - \frac{M}{4}$			
$C_j - Z_j$		$\frac{65}{12} + 2M$	0	0	M	$\frac{-3}{4} + \frac{3}{4}M$	$\frac{23}{12} - \frac{M}{4}$	$0 - \frac{23}{12} + \frac{M}{4}$			
		\uparrow									
0	s_3	-8	0	0	-4	-3	1	4	-1	11	
3	y_2	$-2/3$	1	0	$-1/3$	0	0	$1/3$	0	$5/3$	
5	y_3	$-14/3$	0	1	$-4/3$	-1	0	$4/3$	0	$14/3$	
Z_j		$-\frac{76}{3}$	3	5	$-\frac{23}{3}$	-5	0	$\frac{23}{3}$	0		
$C_j - Z_j$		$\frac{70}{3}$	0	0	$\frac{23}{3}$	5	0	$-\frac{23}{3} + M$	M		

\therefore optimal solution is

$$y_1 = 0, \quad y_2 = \frac{5}{3}, \quad y_3 = \frac{14}{3}$$

$$* \quad W_{\min} = \frac{85}{3}$$

The optimal solution of primal problem

$$\text{is } x_1 = \frac{23}{5}, \quad x_2 = 5, \quad x_3 = 0$$

$$* \quad Z_{\max} = \frac{85}{3}$$

Dual Simplex Method

In this method the solution starts with a basic infeasible optimal solution and work towards feasibility.

Step1. Convert the problem into maximization type.

Step2. Write all the constraints into \leq type.

Step3. Convert all the constraints into equal to type by adding slacks.

Step4. Find the initial solution and express it in the form of a table known as **dual simplex table**.

Step4. Compute $c_j - z_j$ row.

Case1. If any $c_j - z_j > 0$ then **method fails**.

Case2. If all $c_j - z_j \leq 0$ and all solution values are ≥ 0 then the solution obtained is a **optimal feasible solution**.

Case3. if all $c_j - z_j \leq 0$ and any solution value is negative then the solution obtained is **optimal but infeasible** and therefore we can go to **step5** to improve it towards feasibility.

Step5. The row containing least negative solution value is the **key row** and contains the out going variable.

Case1. If all elements of key row are ≥ 0 then the problem **does not have a feasible solution**.

Case2. If at least one element is negative, then find the ratios between the corresponding values of $c_j - z_j$ and these values (ignoring positive and zero elements of key row) .

The column containing the smallest of these ratios contains the incoming variable.

Perform the usual simplex operation to improve the solution till an optimal feasible solution obtained.

Dual simplex method

E-x. Use dual simplex method to solve

① $\text{Max } Z = -3x_1 - 2x_2$

sub to $x_1 + x_2 \geq 1$

$$x_1 + x_2 \leq 7$$

$$x_1 + 2x_2 \geq 10$$

$$x_2 \leq 3$$

$$x_1, x_2 \geq 0$$

Solⁿ converting to maximization canonical form.

$$\text{Max } Z = -3x_1 - 2x_2$$

sub to $-x_1 - x_2 \leq -1$

$$x_1 + x_2 \leq 7$$

$$-x_1 - 2x_2 \leq -10$$

$$x_2 \leq 3$$

$$x_1, x_2 \geq 0$$

Introducing slacks and obtaining initial solution.

$$\text{Max } Z = -3x_1 - 2x_2$$

sub to $-x_1 - x_2 + s_1 = -1$

$$x_1 + x_2 + s_2 = 7$$

$$-x_1 - 2x_2 + s_3 = -10$$

$$x_2 + s_4 = 3$$

$$x_1, x_2, s_1, s_2, s_3, s_4 \geq 0$$

put $x_1=0, x_2=0,$

$$s_1 = -1, s_2 = 7, s_3 = -10, s_4 = 3$$

C_j C_B Basis	-3	-2	0	0	0	0	$So b$
	x_1	x_2	s_1	s_2	s_3	s_4	
0 s_1	-1	-1	1	0	0	0	-1
0 s_2	1	1	0	1	0	0	7
0 s_3	-1	(-2)	0	0	1	0	-10 \rightarrow
0 s_4	0	1	0	0	0	1	3
Z_j	0	0	0	0	0	0	
$C_j - Z_j$	-3	-2	0	0	0	0	
$C_j - Z_j$ θ_{3j}	3	1	-	-	-	-	
		\uparrow					
0 s_1	-1/2	0	1	0	-1/2	0	4
0 s_2	1/2	0	0	1	1/2	0	2
-2 x_2	1/2	1	0	0	-1/2	0	5
0 s_4	(-1/2)	0	0	0	1/2	1	-2 \rightarrow
Z_j	-1	-2	0	0	1	0	
$C_j - Z_j$	-2	0	0	0	-1	0	
$C_j - Z_j$ θ_{4j}	4	-	-	-	-	-	
	\uparrow						

		x_1	x_2	s_1	s_2	s_3	s_4	Soln
0	s_1	0	0	1	0	-1	-1	6
0	s_2	0	0	0	1	1	1	0
-2	x_2	0	1	0	0	0	1	3
-3	x_1	1	0	0	0	-1	-2	4
Z_j		-3	-2	0	0	3	4	
$C_j - Z_j$		0	0	0	0	-3	-4	

\therefore optimal feasible solution is
 $x_1 = 4, x_2 = 3$

$$\& Z_{\max} = -18$$

Note that s_2 is in the basis
 with $s_2 = 0 \Rightarrow$ It is optimal
 feasible degenerate solution.

(2) $\text{Min } Z = x_1 + x_2$
 $2x_1 + x_2 \geq 2$
 $-x_1 - x_2 \geq 1$
 $x_1, x_2 \geq 0$

Soln converting into maximization
 canonical form.

$$\text{Max } W = -Z = -x_1 - x_2$$

$$\text{sub to } \begin{aligned} -2x_1 - x_2 &\leq -2 \\ x_1 + x_2 &\leq -1 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Introducing slacks to obtain initial solution.

$$\text{Max } W = -x_1 - x_2$$

$$\text{sub to } -2x_1 - x_2 + s_1 = -2$$

$$x_1 + x_2 + s_2 = -1$$

$$x_1, x_2, s_1, s_2 \geq 0$$

C_j		-1 x_1	-1 x_2	0 s_1	0 s_2	
0	s_1	(-2)	-1	1	0	-2 →
0	s_2	1	1	0	1	-1
Z_j		0	0	0	0	
$C_j - Z_j$		-1	-1	0	0	
$\frac{C_j - Z_j}{a_{ij}}$		1/2	1	-	-	
		↑				
-1	x_1	1	1/2	-1/2	0	1
0	s_2	0	1/2	1/2	1	-2 →
Z_j		-1	-1/2	1/2	0	
$C_j - Z_j$		0	-1/2	-1/2	0	

Since all $C_j - Z_j \leq 0$, solution is optimal, but $s_2 = -2 \Rightarrow$ It is infeasible.

Here s_2 is outgoing variable but s_2 -row contains no negative coefficient \Rightarrow no other variable enters the basis \Rightarrow Problem does not have feasible optimal solution.