Report of Database Lab3 Name: Nayun Xu

Student ID: 515030910635

1.Design decisions

1.1 HashEquiJoin

I use a *HashMap* to store the elements of the child that has more tuples. And iterate through the other child. Use an *ArrayList* to store the elements with same the key field. I didn't add limit to the size of *HashMap*, so it cannot be applied to large scale data.

2.Changes on API

No changes.

3.Incomplete elements

I still don't know how to implement the lock. The size limit of *HashMap* in *HashEquiJoin*.

4.Experience

I spent about 10 to 15 hours(not exact numbers) on this lab. I first face difficulties when implementing *HashEquiJoin*. I didn't know what to do in this part, and I used google to solve this problem. And the second difficult part is the insert and delete in *HeapFile*. I first tried the same structure as what I did in BTreeFile. But I found it didn't work when adding new pages into the file. I just add new pages into the dirty cache and return it back to the *BufferPool*, which means nobody is taking care of extending the heapfile length. This is because in *BTreeFile*, the *createEmptyPage* method has done this part for me. And also, the adding all pages that are accessed into the dirty cache is not necessary in *HeapFile*, because only one page will be modified, which is different to *BTreeFile* where the leafpage's ancestors may be modified. So finally I modified the structure into a much simpler one. Another problem I encountered is after inserting a tuple, I should modify its *recordId* to match its location, which was missed when I first implemented this part.