NoteHack++

WHEN CONVENIENCE MEETS SECURITY

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What will be covered in this presentation

- DLL Injection
- API Hooking
- Multi-tiered Key Management (like what Ransomware uses)
- Key Derivation PBKDF2
- Crypto Algorithms AES CTR, GCM, RSA

What will **NOT** be covered in this presentation

How to create Ransomware

- Taking a snapshot of the current running processes.
- Finding the process who's name fits the name of the program we want ("NotePad++").
 - This gives us its PID

- Allocating memory in the target process that will contain the path of the DLL to be injected
 - We need to use it as an argument in a function called LoadLibrary

- LoadLibrary is in kernel32.dll
 - And therefore in every process
- We'll find the address of LoadLibrary in the target process.
 - This function can allocate memory for, and load a DLL into the program

- Creating a new thread in the target process which starts at the LoadLibrary function's address, and has the allocated DLL path address as an argument.
 - This will initiate the loading of the DLL.

- The main function in the DLL will be invoked by the operating system.
 - From there we can force the target process to do anything we want.

- Finding the functions that are optimal hooking candidates
- Messing with them

Mission: Adding a jmp command to the original code

- We don't want to mess with the offsets of the entire program
 - Not fun. It's better to overwrite something of the same size
- Instead, we take note of the relative addresses of the commands we want to replace, as well as their respective size in bytes.
 - We will use this information when we overwrite them

- In the injected DLL, we find the base address of the EXE that called it – our target program
- The absolute address is calculated by adding the relative address to the base address.
- The return address is calculated by adding the number of bytes we have overwritten to the absolute address.

- Now that we have hooked the desired API's, we create functions for the hook spots to jump to.
- Problem: registers may be changed due to our actions.

- Now that we have hooked the desired API's, we create functions for the hook spots to jump to.
- **Solution:** use naked inline assembly, so the compiler won't add potentially destructive "fixes".
 - Any registers that are directly affected by our actions need to be protected with push and pop

Now what?

- Now we can inject & hook without any crash or malfunction.
- With the technical part done, we need to start the actual securing of the editor
- We'll start with the design

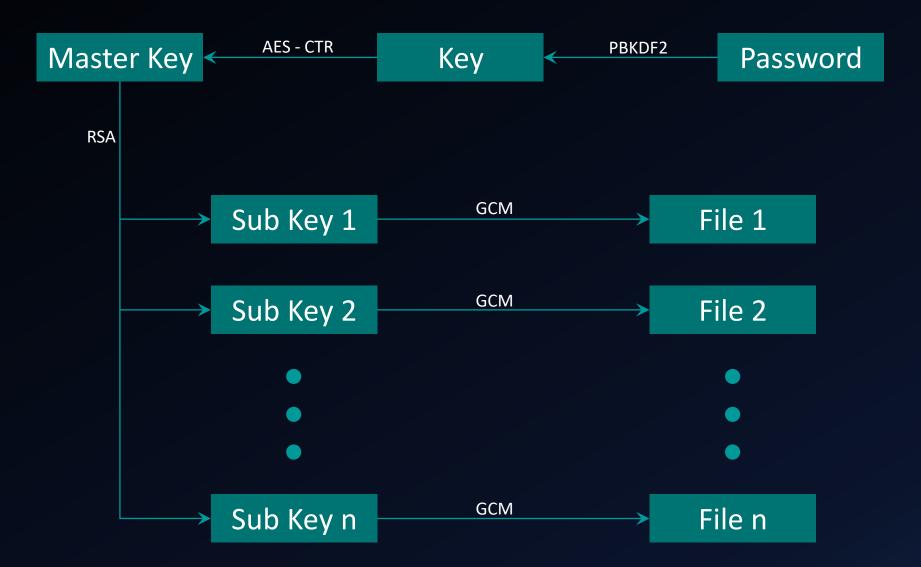
Design

- How do we manage the keys for all the files?
 - We could give each file its own password and trust that the user remembers them all
 - We could use an online authentication server
 - We could save the keys in a secure enclave
 - We can have the user call us for his password whenever he needs it

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- All of the above options are ok (except maybe the last one)
- But isn't there a better way?

- The user enters a password.
- The password "magically" becomes a key (discussed later)
 - for now, just remember the name PBKDF2
- This key protects the "master key" a 2048-bit RSA key.
- The "master key" protects the "file keys" randomly generated keys.
- Each key, encrypts a single file



Is that confused screaming I hear? Let's break it down

Key Derivation

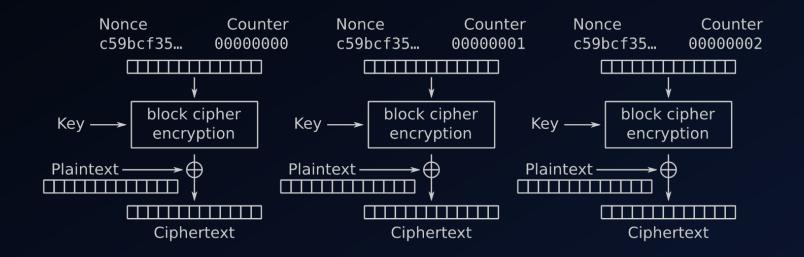
- Makes short passwords longer.
- Frustrate Brute-Force attacks
- As well as *Dictionary* attack

Key Derivation – PBKDF2

- Pros
 - Approved by NIST
 - Deliberately slow
 - Time tested
 - Is included in MbedTLS (our lib of choice)
- Cons
 - A good GPU will greatly reduce calculation time.
- SHA-256 consumes more memory then SHA-512.
- A good iteration number for the average PC: 100,000.

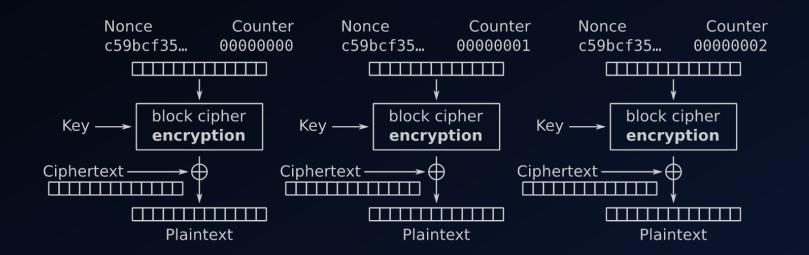
AES – Counter Mode

- Easy implementation.
- Faster than CBC Mode.
- Useful when there is no need for authentication



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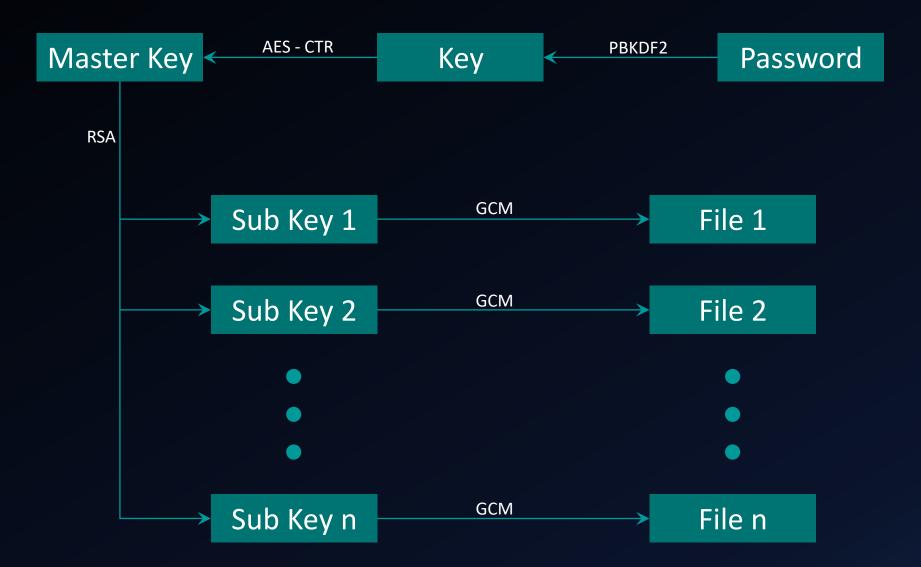
GCM – Galois Counter Mode

- High efficiency & performance
- Fast authentication
- If the tag is invalid, nothing will be decrypted



Breaking it down

- PBKDF2
 - Slow
 - Helps against Brute-force and Dictionary attacks
- AES-CTR
 - Fast and secure
 - No authentication
- GCM
 - Fast authentication



* screaming stops * Isn't that better?

Why use three keys instead of one?

- Easy password changing
- Secure collaboration

Vulnerabilities

That sounds perfect! A system nobody can hack?

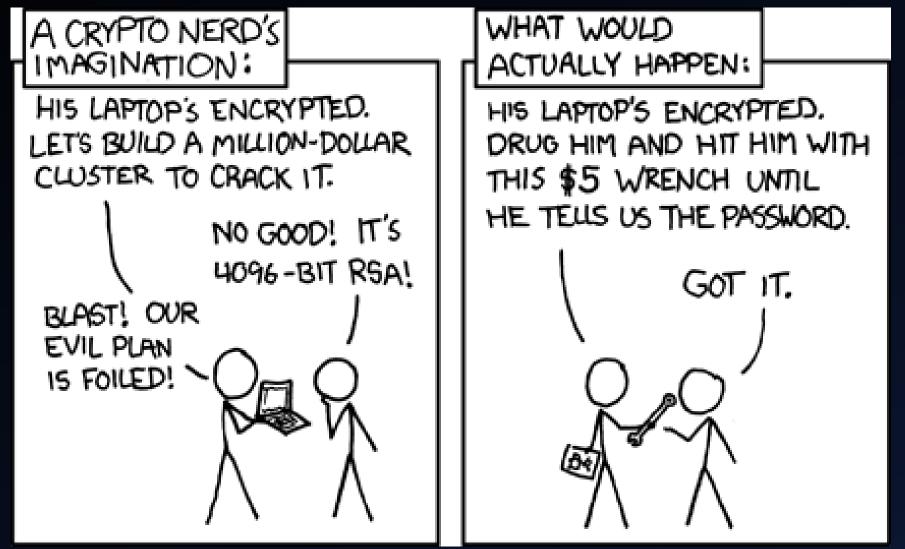
Well...

Vulnerabilities

- Memory dump.
- Another injection.
- A person behind your back.

And...

Vulnerabilities



Future development

- Collaboration
 - Our multi-tiered key management system uses an asymmetric key to encrypt the file keys
 - Using a colleague's public key, we can securely allow him to collaborate without revealing our key or password

