

Prefix monoids of groups and right units of special inverse monoids

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The “driving engine” (part I)

H. H. Wilhelm Magnus (1930/31):

*The word problem for every
one-relator group $\text{Gp}\langle A \mid r = 1 \rangle$ is decidable.*

Reason (the Magnus-Moldavansky hierarchy):

- ▶ $G = \text{Gp}\langle A \mid r = 1 \rangle$ embeds into an HNN-extension of its (f.g.) subgroup $L = \text{Gp}\langle A' \mid r' = 1 \rangle$ w.r.t. a pair of free (“Magnus”) subgroups of L , where $|r'| < |r|$;
- ▶ This suffices to reduce the WP for G to that of L ;
- ▶ Eventually, we end up with a free group of finite rank, where we trivially solve the WP.
- ▶ NB. There is an older approach (Magnus’ original) using amalgamated free products.

The “driving engine” (part II)

Open problem (as of 20 June 2024):

*Does every one-relator monoid
 $\text{Mon}\langle A \mid u = v \rangle$ have a decidable WP?*

S.I.Adian (1966) – The word problem for $\text{Mon}\langle A \mid u = v \rangle$ is decidable for:

- ▶ *special monoids* – the def. relation is of the form $u = 1$,
- ▶ *the case when both u, v are non-empty, and have different initial letters and different terminal letters.*

Adian & Oganessian (1987) – The general problem reduces to two particular cases:

- ▶ $\text{Mon}\langle a, b \mid aUb = aVb \rangle$,
- ▶ $\text{Mon}\langle a, b \mid aUb = a \rangle$ (the “monadic” case).

NB. These presentations define right cancellative monoids.

The Lead Role #1: Prefix monoids (in groups)

Let $G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$ be a group.

The **prefix monoid** of this group (presentation) = the submonoid of G generated by the elements represented by all prefixes of all w_i 's

The prefix monoid is **dependent on the concrete presentation** of G
– one fixed (isomorphism type of a) group can have many presentations, leading to many prefix monoids.

Prefix Membership Problem (PMP): Given a word over $A \cup A^{-1}$, decide whether it represents an element of the prefix monoid (w.r.t. the given group presentation)

IgD & RDG (TrAMS, 2021): A kaleidoscope of sufficient conditions (via amalgamated products and HNN extensions) ensuring **decidability** for the PMP

The Lead Role #2: Right units (in inverse monoids)

Let M be an **inverse** monoid.

$$r \in M \text{ is a right unit} \iff r \mathcal{R} 1 \iff rr^{-1} = 1$$

Fun facts:

- ▶ Right units of M form a **right cancellative submonoid** R of M .
- ▶ If $M = \text{Inv}\langle A \mid w_i = 1 \ (i \in I) \rangle$ (i.e. M is a **special** inverse monoid) then R is generated by elements represented by all prefixes of all w_i 's.
- ▶ So, in the natural map $M \rightarrow G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$, the RU-monoid R of M is mapped **onto** the prefix monoid of G .
- ▶ If M happens to be E -unitary, the restriction of this map to R is a monoid **isomorphism**.
- ▶ Consequently, the RU-monoid of any E -unitary special inverse monoid (SIM) is **group-embeddable**.

The “driving engine” (part III)

Ivanov, Margolis & Meakin (JPAA, 2001):

The (right cancellative) monoid $\text{Mon}\langle A \mid aUb = aVc \rangle$ ($b \neq c$) embeds (as the monoid of right units) into

$$\text{Inv}\langle A \mid aUbc^{-1}V^{-1}a^{-1} = 1 \rangle.$$

Similarly, $\text{Mon}\langle A \mid aUb = a \rangle$ embeds into $\text{Inv}\langle A \mid aUba^{-1} = 1 \rangle$.

Hence, the WP for one-relator monoids reduces to the WP for **one-relator inverse monoids**.

Fun facts: when w is *cyclically reduced* then

- ▶ $\text{Inv}\langle A \mid w = 1 \rangle$ is *E*-unitary;
- ▶ the WP for $\text{Inv}\langle A \mid w = 1 \rangle$ reduces to the PMP for $\text{Gp}\langle A \mid w = 1 \rangle$.

Surprise, surprise...!

RDG (Inventiones, 2020):

*There exists a one-relator special inverse monoid
with an undecidable WP. [!!!]*

Fun facts:

- ▶ the counterexample(s) is/are even *E-unitary*;
- ▶ at the heart of the proof is *Lohrey-Steinberg*'s result (JAlg, 2008) that the *RAAG $A(P_4)$* has a fixed f.g. submonoid with undecidable membership;
- ▶ then, $A(P_4)$ *embeds* into a one-relator group $G = \text{Gp}\langle a, b \mid \dots \rangle$;
- ▶ finally, a one-relator SIM $M = \text{Inv}\langle a, b, t \mid \dots \rangle$ is constructed so that $u \in \{a, b, a^{-1}, b^{-1}\}^*$ represents an element of the “critical” undecidable f.g. submonoid of $G \iff tut^{-1}$ is a *right unit* in M .

Still, this *does not* invalidate the IMM approach.

Know your limits

Guba (1997):

For any monadic $M = \text{Mon}\langle a, b \mid aUb = a \rangle$ constructs $G_M = \text{Gp}\langle x, y, A \mid xWyx^{-1} = 1 \rangle$ (for some $W \in (A \cup \{x, y\})^*$) such that the WP for M reduces to PMP for G_M .

However, there are groups $G = \text{Gp}\langle A \mid w = 1 \rangle$ with:

- ▶ w reduced and undecidable PMP for G (IgD, RDG, 2021);
- ▶ $w = uv^{-1}$ reduced ($u, v \in A^+$) and undecidable PMP for G (Foniqi, RDG, CFNB, to appear);
- ▶ $w \in A^+$ and undecidable submonoid membership problem for G (again, FGNB).

Mon vs Inv

Obviously (imagine *Snape's* voice here), one-relator/f.p. special monoids and special inverse monoids are very different creatures. For example:

- ▶ the group of units U of a $M = \text{Mon}\langle A \mid w = 1 \rangle$ is a one-relator/f.p. group;
- ▶ the RU-monoid of M is a free product of U and a free monoid of finite rank;
- ▶ all other maximal subgroups of M are $\cong U$.

In contrast:

- ▶ the group of units U of a $M = \text{Inv}\langle A \mid w = 1 \rangle$ can be non-one-relator (*RGD, Ruškuc, Jussieu, to appear*);
- ▶ the RU-monoid of M can be even non-f.p.;
- ▶ other maximal subgroups of M can be wildly different from U .

The questions

All of this very much justifies the study of **prefix monoids in f.p. groups** and **RU-monoids in f.p. SIMs** in their own right.

- (1) What can the prefix monoids of f.p. groups be?
- (2) What can the RU-monoids of f.p. SIMs be?
- (3) What are the possible groups of units of these monoids?
- (4) What are the possible Schützenberger groups of these monoids?

Recursive stuff

A group G is **recursively presented** if

$$G = \text{Gp}\langle A \mid w_i = 1 \ (i \in I) \rangle$$

where A is finite and $\{w_i : i \in I\}$ is a **r.e. language** over $A \cup A^{-1}$.

Similarly, a **monoid** is recursively presented if

$$M = \text{Mon}\langle A \mid u_i = v_i \ (i \in I) \rangle$$

where A is finite and $\{(u_i, v_i) : i \in I\}$ is a **r.e. subset** of $A^* \times A^*$.

The Higman Embedding Theorem: A finitely generated group embeds into a f.p. group if and only if it is recursively presented.

- ☞ A finitely generated monoid embeds into a f.p. group if and only if it is group-embeddable and recursively presented.
- ☞ Every prefix monoid (of a f.p. group) is f.g.
⇒ it is recursively presented.

The characterisation of prefix monoids (of f.p. groups)

Two (easy) facts:

- ▶ Every group-embeddable **f.p.** monoid arises as a prefix monoid.
- ▶ If a **group** arises as a prefix monoid then it is **f.p.** So, not all group-embeddable recursively presented monoids are prefix monoids.

Theorem (IgD, RDG, 2023):

For every group-embeddable recursively presented monoid M there is a natural number μ_M such that

$$M * \Sigma_k^*$$

is a prefix monoid (with $|\Sigma_k| = k$) if and only if $k \geq \mu_M$.

Also:

The class of groups of units of prefix monoids is precisely the recursively presented groups.

Recursively enumerable stuff

Let G be a f.p. group (generated by A). Let $L \subseteq (A \cup A^{-1})^*$ be a **recursively enumerable language** such that the set of all elements of G represented by words from L forms a subgroup $H \leq G$. Then H is said to be a **recursively enumerable subgroup** of G .

NB. A r.e. subgroup of G is not necessarily finitely generated.
However, all f.g. (i.e. recursively presented) subgroups of G are r.e.

Theorem (IgD, RDG, 2023):

A group H arises as a Schützenberger group of a prefix monoid (of a f.p. group) $\iff H$ arises as a r.e. subgroup of a f.p. group.

Ingredients:

- ▶ M (left/right) cancellative \implies every Sch-group **embeds** into the group of units of M .
- ▶ For every r.e. subgroup H of a f.p. group G there is a f.p. overgroup $G_1 \geq G$ and $t \in G_1$ such that $G \cap t^{-1}Gt = H$.

RU-monoids (take 1)

Again, some (easy) facts:

- ▶ Every RU-monoid is a **right cancellative** recursively presented monoid.
- ▶ If the monoid of right units of a f.p. SIM is a group
 \implies it is f.p.

Theorem 1 (RDG, Kambites, JEMS, to appear):

*The class of groups of units of f.p. SIMs (and thus of RU-monoids)
is precisely the recursively presented groups.*

Theorem 2 (RDG, Kambites):

*A group arises as a maximal subgroup (i.e. as a group \mathcal{H} -class) of
a f.p. SIM \iff it arises as a r.e. subgroup of a f.p. group.*

RC-presentations

$$M = \text{MonRC}(A \mid \mathfrak{R})$$

$\Leftrightarrow M \cong A^*/\mathfrak{R}^{\text{RC}}$, where \mathfrak{R}^{RC} is the intersection of all congruences σ of A^* such that

- ▶ $\mathfrak{R} \subseteq \sigma$,
- ▶ A^*/σ is right cancellative.

A.J.Cain (2005) (+ Robertson, Ruškuc, 2008): A concept of formal, syntactic derivation for RC-presentations.

Theorem (IgD, RDG, 2023):

Every finitely RC-presented monoid is an RU-monoid.

In a way, this is a generalisation of the Ivanov-Margolis-Meakin result.

RU-monoids (take 2)

Theorem (IgD, RDG, 2023):

The class of Schützenberger groups of RU-monoids is exactly the class of r.e. subgroups of f.p. groups.

Open Problem: Characterise the class of all RU-monoids.

In the remainder of the talk, I'll present **two interesting phenomena** in this vein discovered by IgD+RDG during this Spring's online sessions.

The RU-monoid in the Gray-Ruškuc construction (1)

RDG, Ruškuc: For any group G (f.p. or not) and f.g. submonoid T of G , a(n E -unitary) SIM M is constructed (which is f.p. when G is) such that:

- ▶ $U(M) \cong G * U(T)$,
- ▶ if the monoid of right units of M is f.p. so must be both G and T .

With the right choice of parameters, this produces:

- ▶ a one-relator SIM whose group of units is **not one-relator**;
- ▶ a one-relator SIM whose group of units is f.p. but whose RU-monoid is **not f.p.**;
- ▶ a f.p. SIM whose group of units is **not f.p.**

The RU-monoid in the Gray-Ruškuc construction (2)

IgD, RDG (2024):

The RU-monoid of M = the greatest right cancellative image of the HNN-like Otto-Pride extension of G w.r.t. $T \hookrightarrow G$ =

$$\text{MonRC}\langle A, B, t \mid u_i = v_i \ (i \in I), \ tw_j = b_j t \ (j \in J) \rangle$$

where $G = \text{Mon}\langle A \mid u_i = v_i \ (i \in I) \rangle$ and $T = \langle w_j : j \in J \rangle_G$.

Hence:

- ▶ If G is f.p. then the RU-monoid of M is necessarily **finitely RC-presented**;
- ▶ The group of units $U(M)$ can still be **not f.p.**, and also the RU-monoid can be **not f.p.** (as a monoid!);
- ▶ There is a **finitely RC-presented** monoid S in which the complement of the group of units $S \setminus U$ is an **ideal**, and still U is **not f.p.**

Conclusion: RC-presentations are strange animals!

The Gray-Kambites construction (1)

Realising an arbitrary recursively presented group as the **group of units** of a f.p. SIM. Here we present a slight generalisation (by IgD & RDG).

$$T = \text{MonRC}\langle A \mid u_i = v_i \ (i = 1, \dots, k) \rangle$$

$$S = \langle B \rangle_T - \text{a f.g. submonoid}$$

$M_{T,S}$ – a f.p. SIM gen. by A and $p_0, p_1, \dots, p_k, z, d$ subject to

$$p_i a p_i^{-1} p_i a^{-1} p_i^{-1} = 1 \quad (a \in A, i = 0, 1, \dots, k)$$

$$p_i u_i d^{-1} v_i^{-1} p_i^{-1} = 1 \quad (i = 1, \dots, k)$$

$$p_0 d p_0^{-1} = 1$$

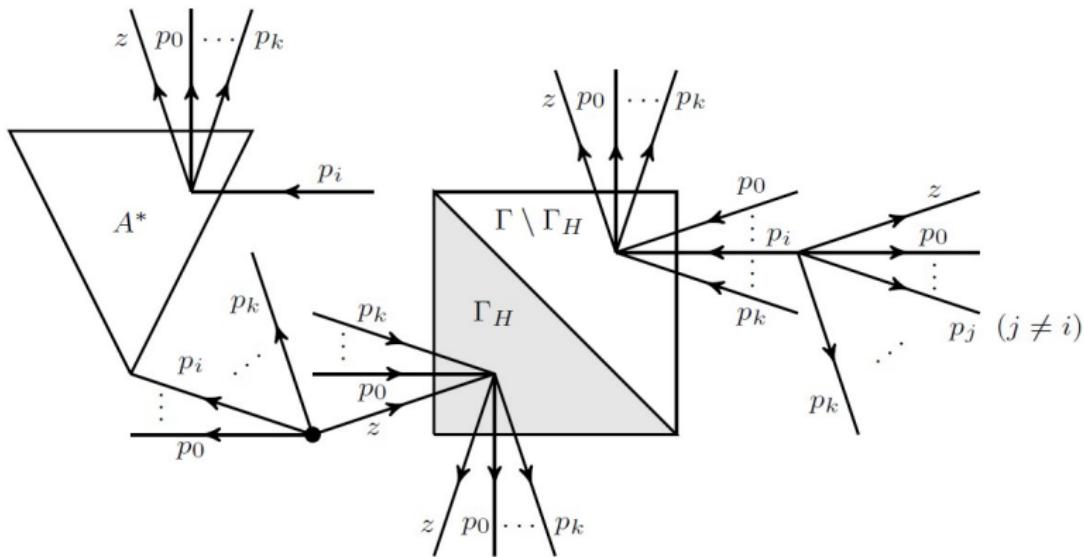
$$z b z^{-1} z b^{-1} z^{-1} = 1 \quad (b \in B)$$

$$z \left(\prod_{i=0}^k p_i^{-1} p_i \right) z^{-1} = 1.$$

The Gray-Kambites construction (2)

RDG, Kambites (JEMS, to appear): When $T = G$ (a group given by a finite special monoid pres.) and $S = H$ (a f.g. subgroup), then

$$U(M_{G,H}) \cong H.$$



The Gray-Kambites construction (3)

So, what is the RU-monoid of $M_{T,S}$?

IgD, RDG (2024): RC-presented by p_i , q_i ($= zp_i^{-1}$) ($0 \leq i \leq k$), $a^{(i)}$ ($= p_i a p_i^{-1}$) ($a \in A$, $0 \leq i \leq k$), $b^{(z)}$ ($= z b z^{-1}$) ($b \in B$), and relations

$$q_i w^{(i)} p_i = q_0 w^{(0)} p_0 \quad (w \in A^*, i = 1, \dots, k)$$

$$q_i u^{(i)} = q_i v^{(i)} \quad (u, v \in A^* \text{ s.t. } u = v \text{ holds in } T, i = 0, 1, \dots, k)$$

$$q_i b^{(i)} = b^{(z)} q_i \quad (b \in B, i = 0, 1, \dots, k)$$

NB. For all $u, v \in B^*$ s.t. $u = v$ holds in S , $u^{(z)} = v^{(z)}$ can be RC-derived. In fact, $\langle b^{(z)} : b \in B \rangle \cong S$.

The Gray-Kambites construction (4)

For example, when we take $T = \{a\}^*$ and $S = \langle \emptyset \rangle = \{1\}$ (and a silly presentation for T , say $a = a$, to have $k = 1$) we get the RU-monoid

$$\text{MonRC}\langle a_1, a_1, p_0, p_1, q_0, q_1 \mid q_0 a_0^n p_0 = q_1 a_1^n p_1 \ (n \geq 0) \rangle.$$

This can be shown to be:

- ▶ **not** finitely RC-presented,
- ▶ with a **trivial** group of units.

Conclusion: There are non-finitely RC-presented RU-monoids out there!

Thank you!

