

Pretzel Monoids

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Left Adequate Monoids

Definition

A monoid M is *left adequate* if:

- ① Idempotents of M commute;
- ② For all $a \in M$, there exists a unique idempotent $a^+ \in E(M)$ such that

$$\forall x, y \in M \quad xa = ya \iff xa^+ = ya^+.$$

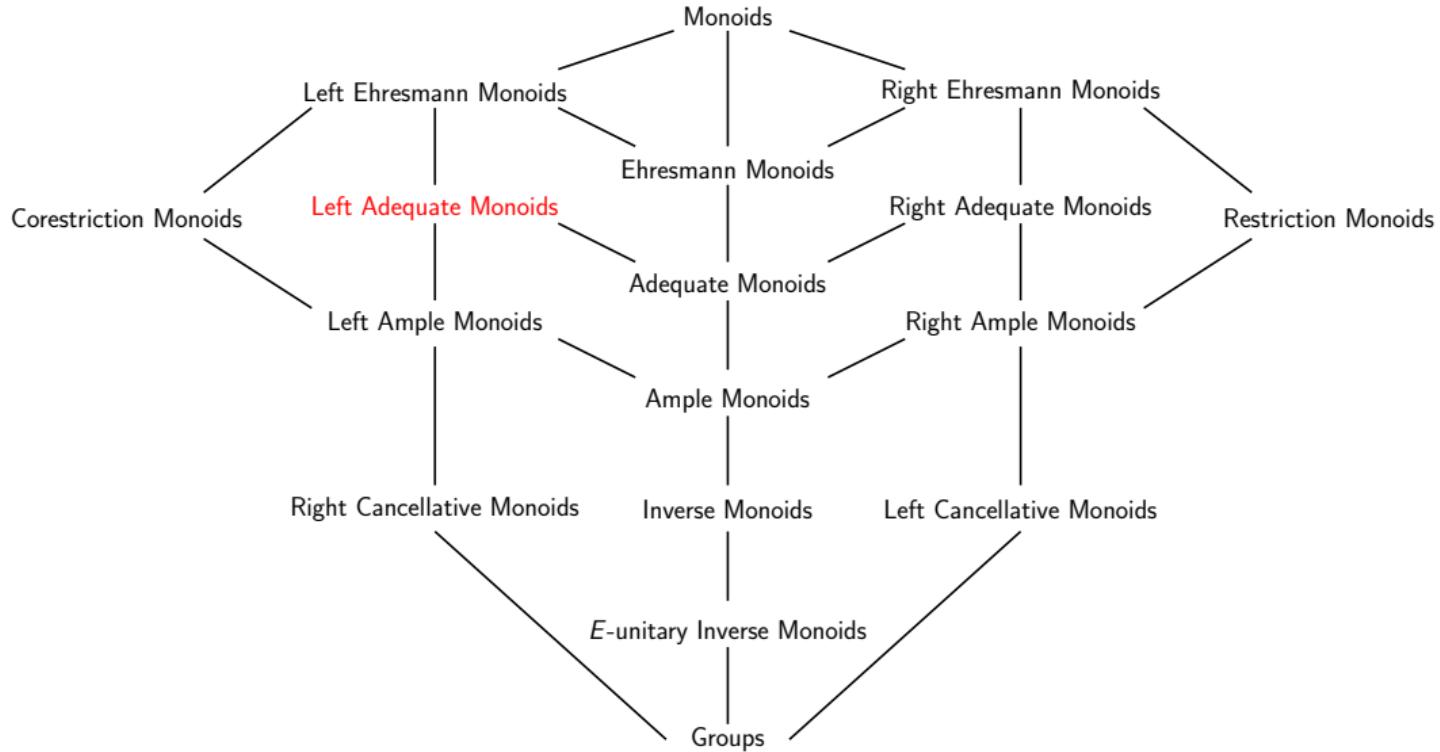
Definition

Equivalently, a monoid equipped with a unary operation $+$ is called *left adequate* if it satisfies the defining identities:

$$a^+a = a, \quad (a^+b^+)^+ = a^+b^+, \quad a^+b^+ = b^+a^+, \quad (ab)^+ = (ab^+)^+,$$

$$a^2 = a \rightarrow a = a^+ \quad \text{and} \quad ac = bc \rightarrow ac^+ = bc^+.$$

A Big Diagram



A Potted History

- Kilp 1973, *Commutative monoids all of whose principal ideals are projective.*
- Fountain 1976, *Right PP monoids with central idempotents.*

Remark. A monoid M is *right PP* (modernly known as *right abundant*) if and only if every \mathcal{L}^* -class contains an idempotent. Very reminiscent of regular and inverse semigroups!

- Fountain 1977, *A class of right PP monoids.*
- Fountain 1979, *Adequate semigroups.*

Definition

A right adequate semigroup is called *right ample* if $eM \cap aM = eaM$ for all $e \in E(M)$, $a \in M$.

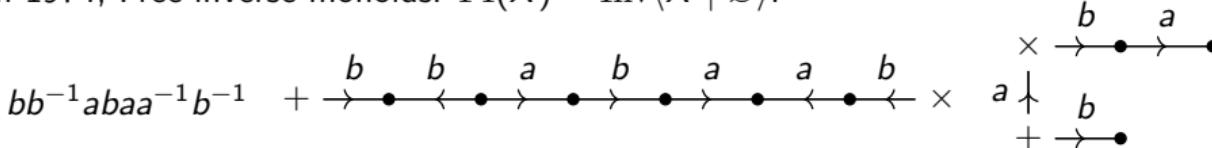
Theorem (Fountain 1977)

Let M be a right ample monoid. Then S is the image of some “proper” right ample monoid under some morphism θ where $a\theta = b\theta \implies a\mathcal{L}^*b$.

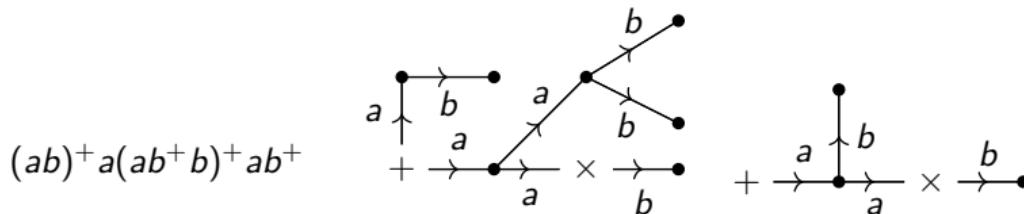
- See also Hollings 2009, *From right PP monoids to restriction semigroups: a survey.*

A Potted History 2

- Munn 1974, Free inverse monoids. $\text{FI}(X) = \text{Inv}\langle X \mid \emptyset \rangle$.



- Fountain 1991, *Free right type A semigroups*.
 - Fountain, Gomes, Gould 2009, *The free ample monoid*. $\text{FLAm}(X) = \text{LAm}\langle X \mid \emptyset \rangle$.



- Kambites 2009-11, *Free adequate semigroups* $\text{FLAd}(X) = \text{LAd}\langle X \mid \emptyset \rangle$.

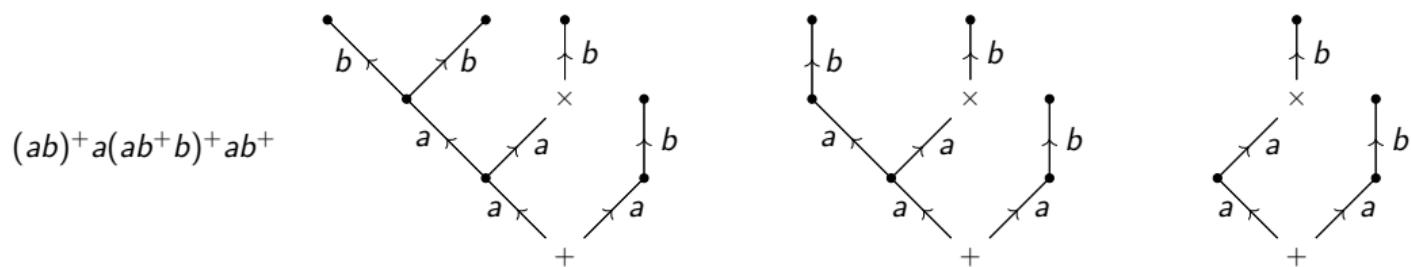
Free Left Adequate Monoids

For trees in $\text{FLAd}(X)$, first draw the tree as in $\text{FLAm}(X)$...

Definition

The *retraction* of a tree is the smallest image under an idempotent endomorphism which fixes $+$ and \times .

Fact: The retraction of a tree is unique up to isomorphism.



The Goal

Question: What results in inverse or cancellative semigroups can we generalise to (left) adequate or ample monoids?

- McAlister's covering theorem and the P -theorem (Fountain 1970s).
- Munn trees and free inverse monoids (Fountain, Gomes, Gould, Kambites, 2000s).
- Stephen's procedure...? (Likely very hard!)
- *Expansions!*

Semigroup Expansions

Definition (Birget, Rhodes 1984)

Let $\mathcal{C} \subseteq \mathcal{D} \subseteq \mathbf{Sgp}$. An *expansion of \mathcal{C} to \mathcal{D}* is a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ such that there is a natural transformation $\eta : F \Rightarrow \iota_{\mathcal{C} \rightarrow \mathcal{D}}$ whose components η_S are all surjective.

i.e. for all $S \in \mathcal{C}$, there is a semigroup $F(S) \in \mathcal{D}$ and a surjective morphism $\eta_S : F(S) \rightarrow S$ such that whenever $\tau : S \rightarrow T$ is a morphism, there is a morphism $F(\tau) : F(S) \rightarrow F(T)$ making the following diagram commute:

$$\begin{array}{ccc} F(S) & \xrightarrow{F(\tau)} & F(T) \\ \eta_S \downarrow & & \downarrow \eta_T \\ S & \xrightarrow{\tau} & T \end{array}$$

Theorem (Birget, Rhodes 1984 / Szendrei 1989)

There is an expansion $\text{Sz} : \mathbf{Gp} \rightarrow \mathbf{FlInv}$ given by $\text{Sz}(G) = \{(H, g) : H \subseteq G \text{ finite and } 1, g \in H\}$.

Theorem (Szendrei 1989)

Sz is left adjoint to the maximal group image functor $\sigma^\natural : \mathbf{FlInv} \rightarrow \mathbf{Gp}$.

Expansions of other Categories

Recall that $\text{FI}(X)$ was constructed by ‘tracing’ the Cayley graph of $\text{FG}(X)$... what about other Cayley graphs?

Theorem (Margolis, Meakin 1989)

Let X be a set and let G be an X -generated group. There is an expansion $\mathcal{M} : \mathbf{XGp} \rightarrow \mathbf{XEInv}$ given by $\mathcal{M}(G) = \{(\Gamma, g) : \Gamma \text{ is a finite connected subgraph of } \text{Cay}(G), 1, g \in V(\Gamma)\}$.
 \mathcal{M} is left adjoint to the maximal group image functor $\sigma^\natural : \mathbf{XEInv} \rightarrow \mathbf{XGp}$.

Theorem (Gould 1996, + Gomes 2000)

Let X be a set and let M be an X -generated monoid. Define

$$\mathcal{G}(M) = \{(\Gamma, m) : \Gamma \text{ is a finite connected subgraph of } \text{Cay}(M), 1, m \in V(\Gamma)\}.$$

Then \mathcal{G} forms expansions $\mathbf{XRC} \rightarrow \mathbf{XPLAm}$ and $\mathbf{XU} \rightarrow \mathbf{XPWLAm}$. Moreover, \mathcal{G} is left adjoint to taking the maximal right cancellative image and maximal unipotent image respectively.

Question

Can we find an expansion $\mathbf{XRC} \rightarrow \mathbf{XLAd}$? Preferably with some graphical interpretation?

Pretzels!

Fix a set X and an X -generated right cancellative monoid C .

Definition

An *idempath* in an X -labelled digraph Γ is a path labelled by a word $x_1x_2 \cdots x_n$ which is equal to the identity in C . We take the empty path with label ϵ to have $\epsilon =_C 1$.

An *idempath identification* in Γ is the process of ‘cycling up’ an idempath.

Lemma (H., Kambites, Szakács 2024)

Given a tree $T \in \text{FLAd}(X)$, there exists a unique graph obtainable by sequentially performing all non-trivial idempath identifications (in any order) to T .

Definition

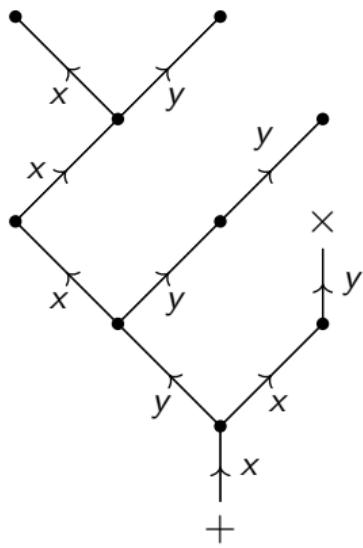
Given any tree $T \in \text{FLAd}(X)$, perform the following:

- ① Idempath identify as far as possible...
- ② ...then retract anything in the result which can retract (take minimal image under idempotent graph endomorphisms).

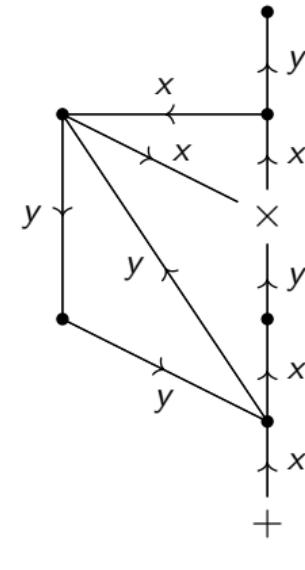
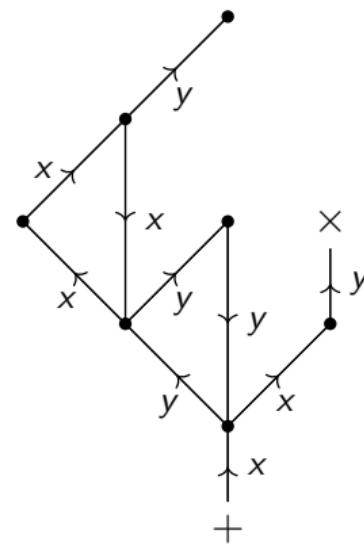
We call the (uniquely obtained) result the *pretzel* of T , denoted \widetilde{T} .

Example

Take $X = \{x, y\}$ and $C = C_3 \times C_3 = \text{Mon}\langle x, y \rangle$.



T



\tilde{T}

(2, 1, 0)-algebras

Define a multiplication $\overline{\widetilde{S}} \cdot \overline{\widetilde{T}}$ on pretzels as follows:

- ① Glue $\overline{\widetilde{T}}$ to $\overline{\widetilde{S}}$, start-to-end.
- ② Pretzel-ify the result (note that new idempaths could have been created!).

Define a unary operation $+$ on pretzels as follows:

- ① Relabel the end vertex of $\overline{\widetilde{T}}$ to be the start vertex.
- ② Pretzel-ify the result (note that new retractions might be possible!).

Theorem (H., Kambites, Szakács 2024)

The set of all pretzels $\mathcal{PT}(C; X)$ forms an X -generated left adequate monoid.

Theorem (H., Kambites, Szakács 2024)

$\mathcal{PT}(C; X) \cong \mathbf{LAd}\langle X \mid w^2 = w \text{ for } w \in X^* \text{ s.t. } w =_c 1 \rangle.$

Margolis-Meakin Expansions vs. Pretzels

Properties of $\mathcal{M}(G)$

- ① $\mathcal{M}(\text{FG}(X)) \cong \text{FI}(X)$.
- ② $\mathcal{M}(G)$ is finite $\iff G$ is finite.
- ③ Elements are subgraphs of $\text{Cay}(G)$.
- ④ $\mathcal{M}(G) \cong \mathbf{Inv}\langle X \mid w^2 = w \text{ for } w \in X^* \text{ s.t. } w =_G 1 \rangle$.
- ⑤ \mathcal{M} defines an expansion $\mathbf{XGp} \rightarrow \mathbf{XEInv}$.

Properties of $\mathcal{PT}(C)$

- ① $\mathcal{PT}(X^*) \cong \text{FLAd}(X)$.
- ② $\mathcal{PT}(C)$ is finite $\iff C$ is finite $\implies C$ is a group.
- ③ Elements are trees of strongly connected subgraphs of $\text{Cay}(C)$.
- ④ $\mathcal{PT}(C; X) \cong \mathbf{LAd}\langle X \mid w^2 = w \text{ for } w \in X^* \text{ s.t. } w =_C 1 \rangle$.

Theorem (H., Kambites, Szakács 2024)

\mathcal{PT} defines an expansion $\mathbf{XRC} \rightarrow \mathbf{XLAd}$.

Open Questions and What's Next

- What about right adequate and two-sided adequate pretzel monoids?
- Can we find geometric interpretations of other analogues of Margolis-Meakin expansions in the left adequate setting, perhaps one such that $M(C)$ has maximal right cancellative image C ?

Proposition

The maximal right cancellative image of $\mathcal{PT}(C)$ is $\mathbf{RC}\langle X \mid w = 1 \text{ for } w \in X^* \text{ s.t. } w =_C 1 \rangle$.
In particular, it is two-sided cancellative.

Proposition

\mathcal{PT} is not left adjoint to the maximal right cancellative image functor.

- Can we apply similar pretzel-style techniques in F -inverse land? In particular for the free F -inverse monoid...?
- What about other interesting presentations of (left) adequate monoids?

Thank you!