

## Lecture 3: Solving linear systems of equations

Reading:



Ch. 1,  
2.1-2.2



Ch. 2-3  
(skip 3.8)

Outline: Systems of linear equations

- Examples, non-examples
- General problem, correspondence to matrices
- Solving by Gaussian elimination
  - complexity analysis
- Geometry & stability of solutions
- Homogeneous & non-homogeneous systems
- Solving with a computer: Matlab/Mathematica
  - timing
  - sparse matrices
  - 1D and 2D lattices
- When is there a solution?
- Solving repeatedly: LU decomposition
- Compressed sensing

## Asymptotic complexity of matrix multiplication

[http://en.wikipedia.org/wiki/Computational\\_complexity\\_of\\_mathematical\\_operations#Matrix\\_algebra](http://en.wikipedia.org/wiki/Computational_complexity_of_mathematical_operations#Matrix_algebra)

Matrix-vector multiplication

$$\begin{matrix} n \\ m \end{matrix} \boxed{A}$$

$$(A \vec{v})_i = \sum_{j=1}^n A_{ij} v_j$$

↑  
m entries  
to calculate

each term  
takes  $O(n)$  time

$\Rightarrow O(m \cdot n)$  time total

e.g.  $O(n^2)$  if  $m=n$

→ doubling n, time increases by  $\sim \times 4$

```
>> n = 10000;
>> A = rand(n,n); v = rand(n,1);
>> start = cputime; for i = 1:10 A * v; end; cputime - start
```

ans =

1.7813

```
>> n = 20000;
>> A = rand(n,n); v = rand(n,1);
>> start = cputime; for i = 1:10 A * v; end; cputime - start
```

ans =

6.7969

×2.82

## Sparse matrix-vector mult.

$$(\vec{A} \vec{v})_i = \sum_{j=1}^n A_{ij} v_j$$

↑  
each term takes

$\mathcal{O}(\# \text{ nonzero entries in row } i)$

$\Rightarrow \mathcal{O}(\# \text{ nonzero entries in } A)$

```
>> n = 100000; s = 10*n;
>> i = randi(n,s,1); j = randi(n,s,1); d = rand(s,1);
>> A = sparse(i,j,d,n,n); v = rand(n,1);
>> t = cputime; for i=1:10 A*v; end; cputime-t
ans =
0.0625

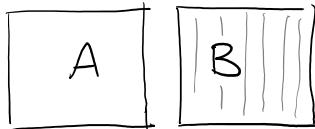
>> n = 200000; s = 10*n;
>> i = randi(n,s,1); j = randi(n,s,1); d = rand(s,1);
>> A = sparse(i,j,d,n,n); v = rand(n,1);
>> t = cputime; for i=1:10 A*v; end; cputime-t
ans =
0.1250
```

$\times 2.0$

## Matrix-matrix multiplication

$$(AB)_{ij} = \sum_{k=1}^n A_{ik} B_{kj}$$

↑  
 $n^2$  terms to calculate    each term takes  $\mathcal{O}(n)$  time     $\Rightarrow \mathcal{O}(n^3)$



But Matlab is faster! (see HW1)

## Strassen's algorithm [http://en.wikipedia.org/wiki/Strassen\\_algorithm](http://en.wikipedia.org/wiki/Strassen_algorithm)

$$A = \left( \begin{array}{c|c} A_{11} & A_{12} \\ \hline A_{21} & A_{22} \end{array} \right)_{n/2} \quad B = \left( \begin{array}{c|c} B_{11} & B_{12} \\ \hline B_{21} & B_{22} \end{array} \right)_{n/2}$$

$$AB = \left( \begin{array}{c|c} A_{11}B_{11} + A_{12}B_{21} & A_{11}B_{12} + A_{12}B_{22} \\ \hline A_{21}B_{11} + A_{22}B_{21} & A_{21}B_{12} + A_{22}B_{22} \end{array} \right)$$

8  $n/2 \times n/2$  matrix multiplications

$$\text{time } T(n) = 8T(n/2) + \mathcal{O}(n^2)$$

$$\begin{aligned} &\sim n^2 + 8\left(\frac{n}{2}\right)^2 + 64\left(\frac{n}{4}\right)^2 + \dots \\ &\sim n^2 (1+2+4+8+\dots+2^{\log_2 n}) \\ &= \Theta(n^3) \end{aligned}$$

Try computing

$$M_1 = (A_{11} + A_{12})(B_{11} + B_{22})$$

$$M_5 = (A_{11} + A_{12})B_{22}$$

$$M_2 = (A_{21} + A_{22})B_{11}$$

$$M_6 = (A_{21} - A_{11})(B_{11} + B_{12})$$

$$M_3 = A_{11}(B_{12} - B_{22})$$

$$M_7 = (A_{12} - A_{22})(B_{21} + B_{22})$$

$$M_4 = A_{22}(B_{21} - B_{11})$$

$$\Rightarrow AB = \begin{pmatrix} M_1 + M_4 - M_5 + M_7 & M_3 + M_5 \\ M_2 + M_4 & M_1 - M_2 + M_3 + M_6 \end{pmatrix}$$

$(A_{21}B_{11} + A_{22}B_{11}) + (A_{22}B_{21} - A_{21}B_{11})$

$\nexists \frac{n}{2} \times \frac{n}{2}$  matrix multiplications

$$\begin{aligned} \text{time } T(n) &= 7T\left(\frac{n}{2}\right) + O(n^2) \\ &\sim n^2 \left(1 + 7\frac{1}{2^2} + 7^2\frac{1}{4^2} + \dots + \left(\frac{7}{4}\right)^{\log_2 n}\right) \\ &\sim n^{2+\log_2(7/4)} \\ &= O(n^{\log_2 7} = 2.81\dots) \end{aligned}$$

$O(n^{2.376})$  [Coppersmith-Winograd '90]  
but the constant factor is impractical

Stothers 2010: 2.374

Williams 2012: 2.3728642

Le Gall 2014: 2.3728639

Question: Is the correct exponent 2?

Remark: The asymptotic complexity of solving systems of linear equations is the same as matrix multiplication.

## LINEAR EQUATIONS

Examples:

$$\left\{ \begin{array}{l} x = 3 \\ 2x - y = 3 \end{array} \right. \quad \checkmark$$

$$\left\{ \begin{array}{l} x^2 + xy = 2 \\ x + y = 3 \end{array} \right. \quad \times \quad \text{quadratic equation}$$

$$\left\{ \begin{array}{l} 2x - y = 3 \\ x - y = 2 \\ x + y = 0 \end{array} \right. \quad \checkmark$$

$$\left\{ \begin{array}{l} x + y + z = 0 \\ x + \cos(y) + z = 2 \end{array} \right. \quad \times \quad \text{transcendental equation}$$

General problem:

Given: Linear equations

$n$  unknowns  $x_1, \dots, x_n$

equations

$$\left\{ \begin{array}{l} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + a_{23}x_3 + \dots + a_{2n}x_n = b_2 \\ \vdots \end{array} \right.$$

$$\begin{matrix} \text{for} \\ \text{eq} \\ \text{M} \end{matrix} \left\{ \begin{array}{l} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \dots + a_{1n}x_n = b_1 \\ \vdots \\ a_{m1}x_1 + a_{m2}x_2 + a_{m3}x_3 + \dots + a_{mn}x_n = b_m \end{array} \right.$$



$$\begin{pmatrix} a_{11} & a_{12} & a_{13} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & & a_{2n} \\ \vdots & & & & \vdots \\ a_{m1} & a_{m2} & \cdots & & a_{mn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{pmatrix}$$

$$\begin{matrix} \uparrow \\ A \vec{x} = \vec{b} \end{matrix}$$

$m \times n \quad n \times 1 \quad m \times 1$

Goal: Solve for  $\vec{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$

- (- If there isn't a solution, then say so.  
 - If there are multiple solutions, find them all!)

Note: We will see later that there is always either

- \* one unique solution,
- \* no solution, or
- \* infinitely many solutions

Example: Solve

$$\begin{aligned} 2x - y &= 3 \\ x - y &= 2 \\ x + y &= 0 \end{aligned}$$

$$\Leftrightarrow \begin{pmatrix} 2 & -1 \\ 1 & -1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 3 \\ 2 \\ 0 \end{pmatrix}$$

Answer:

$$\text{Adding } \begin{cases} x - y = 2 \\ x + y = 0 \end{cases} \Rightarrow 2x = 2$$

$$\begin{matrix} \Downarrow \\ x = 1 \\ \Downarrow \\ y = -1 \end{matrix}$$

$$\text{shorthand notation:} \quad \left( \begin{array}{cc|c} 2 & -1 & 3 \\ 1 & -1 & 2 \\ 1 & 1 & 0 \end{array} \right) \rightarrow \left( \begin{array}{cc|c} 2 & -1 & 3 \\ 1 & -1 & 2 \\ 2 & 0 & 2 \end{array} \right)$$

$$\boxed{\begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 1 \\ -1 \end{pmatrix}}$$

Check your answer!

Key idea: This works by cancelling the coefficient of  $y$ .

Matlab

```
>> A = [2 -1; 1 -1; 1 1];
>> b = [3 2 0]';
>> x = A \ b;
x =
```

Python

```
import numpy as np
A = [[2,-1],[1,-1],[1,1]]
b = [3,2,0]
x = np.linalg.solve(A[:2],b[:2]) ← using only
print(x)
```

```

>> x = A \ b
x =
    1.0000
   -1.0000
ans =
    1.0e-15 *
      0
      0
     -0.2220

```

```

A = [[2,-1],[1,-1],[1,1]]
b = [3,2,0]
x = np.linalg.solve(A[:2],b[:2])  
using only  
first two equations
print(x)
print(np.dot(A,x) - b) #check

```

```

[ 1. -1.]
[0. 0. 0.]
or:
x = np.linalg.lstsq(A, b, rcond=None)[0]
x
array([ 1., -1.])

```

Example: Polynomial interpolation

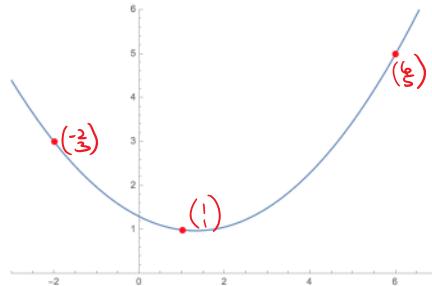
2 points determine a line

3 points determine a quadratic curve (parabola)

any  $n+1$  points  $(x_1, y_1), \dots, (x_{n+1}, y_{n+1})$   
determine a polynomial of degree  $\leq n$ .

Find the equation  $y = ax^2 + bx + c$  for the parabola  
that goes through

$$\begin{aligned} &(1, 1) \\ &(6, 5) \\ &(-2, 3) \end{aligned}$$



Answer: Set up three equations:

$$(1, 1): 1 = a \cdot 1^2 + b \cdot 1 + c = a + b + c$$

$$(6, 5): 5 = a \cdot 6^2 + b \cdot 6 + c = 36a + 6b + c \quad \text{---}^2$$

$$(-2, 3): 3 = a \cdot (-2)^2 + b \cdot (-2) + c = 4a - 2b + c \quad \text{---}^3$$

$$\left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 36 & 6 & 1 & 5 \\ 4 & -2 & 1 & 3 \end{array} \right) \xrightarrow{-36} \left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & -30 & -35 & -31 \\ 0 & -6 & -3 & -1 \end{array} \right) \xrightarrow{-4} \left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & -30 & -35 & -31 \\ 0 & -6 & -3 & -1 \end{array} \right) \xrightarrow{-5} \left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & -30 & -35 & -31 \\ 0 & 0 & 0 & 1 \end{array} \right) \xrightarrow{-6b-3c=-1}$$

$$\rightarrow \left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & 6 & 3 & 1 \\ 0 & 0 & 20 & 26 \end{array} \right) \xrightarrow{\downarrow} c = \frac{13}{10}$$

$$\begin{aligned} &\text{Now work backward:} \\ &\left( \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & 6 & 3 & 1 \\ 0 & 0 & 1 & \frac{13}{10} \end{array} \right) \xrightarrow{-1} \left( \begin{array}{ccc|c} 1 & 1 & 0 & -\frac{3}{10} \\ 0 & 6 & 0 & -2\frac{9}{10} \\ 0 & 0 & 1 & \frac{13}{10} \end{array} \right) \\ &\xrightarrow{-1} \left( \begin{array}{ccc|c} 1 & 0 & 0 & \frac{11}{60} \\ 0 & 1 & 0 & -\frac{29}{60} \\ 0 & 0 & 1 & \frac{13}{10} \end{array} \right) \end{aligned}$$

$$\rightarrow \left( \begin{array}{ccc|c} 0 & 1 & 0 & -\frac{29}{60} \\ 0 & 0 & 1 & \frac{13}{10} \end{array} \right) \rightarrow \left( \begin{array}{cc|c} 0 & 0 & -\frac{29}{60} \\ 0 & 1 & \frac{13}{10} \end{array} \right)$$

$$(a, b, c) = \left( \frac{11}{60}, -\frac{29}{60}, \frac{13}{10} \right)$$

Check this!

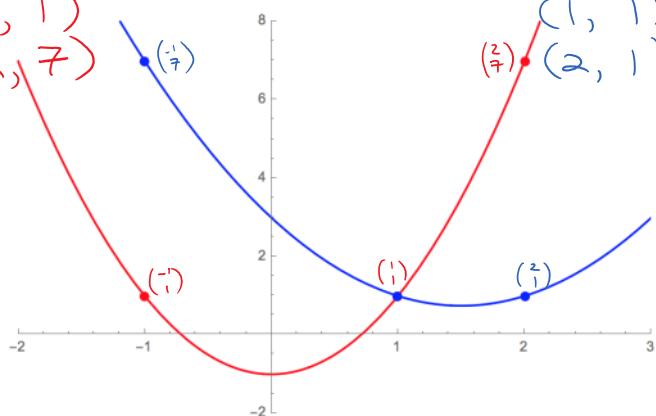
Matlab:

```
>> x = [1 6 -2]';           >> abc = A \ y
>> A = [x.^2 x ones(3,1)]    abc =
A =
          1      1      1
         36      6      1
          4     -2      1
>> y = [1 5 3]';           >> format rat; abc
abc =
          11/60
         -29/60
         13/10
```

Example:

Find the parabola through

(-1, 1)  
(1, 1)  
(2, 7)



AND Find the parabola through

(-1, 7)  
(1, 1)  
(2, 1)

Answer: Set up the equations  $x^2a + xa + 1 = y$

$$\begin{pmatrix} (-1)^2 & -1 & 1 \\ 1^2 & 1 & 1 \\ 2^2 & 2 & 1 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 7 \end{pmatrix} \quad \begin{pmatrix} (-1)^2 & -1 & 1 \\ 1^2 & 1 & 1 \\ 2^2 & 2 & 1 \end{pmatrix} \begin{pmatrix} a' \\ b' \\ c' \end{pmatrix} = \begin{pmatrix} 7 \\ 1 \\ 1 \end{pmatrix}$$

↑ same matrix!

$$\Rightarrow \begin{pmatrix} 1 & -1 & 1 \\ 1 & 1 & 1 \\ 4 & 2 & 1 \end{pmatrix} \begin{pmatrix} a & a' \\ b & b' \\ c & c' \end{pmatrix} = \begin{pmatrix} 1 & 7 \\ 1 & 1 \\ 7 & 1 \end{pmatrix}$$

We can solve them together

$$\begin{pmatrix} 1 & -1 & 1 & | & 7 \end{pmatrix} \quad \begin{pmatrix} 1 & -1 & 1 & | & 1 & 7 \end{pmatrix}$$

We can solve them together

$$\begin{array}{c} \left( \begin{array}{ccc|cc} 1 & -1 & 1 & 1 & 7 \\ 1 & 1 & 1 & 1 & 1 \\ 4 & 2 & 1 & 7 & 1 \end{array} \right) \xrightarrow{\text{R}_2 - \text{R}_1} \left( \begin{array}{ccc|cc} 1 & -1 & 1 & 1 & 7 \\ 0 & 2 & 0 & 0 & -6 \\ 0 & 6 & -3 & 3 & -27 \end{array} \right) \\ \xrightarrow{\text{R}_3 - 3\text{R}_2} \left( \begin{array}{ccc|cc} 1 & -1 & 1 & 1 & 7 \\ 0 & 1 & 0 & 0 & -3 \\ 0 & 2 & -1 & 1 & -9 \end{array} \right) \xrightarrow{\text{R}_3 - 2\text{R}_2} \left( \begin{array}{ccc|cc} 1 & 0 & 1 & 1 & 4 \\ 0 & 1 & 0 & 0 & -3 \\ 0 & 0 & 1 & -1 & 3 \end{array} \right) \xrightarrow{\text{R}_1 - \text{R}_3} \\ \Rightarrow I \left( \begin{array}{cc} a & a' \\ b & b' \\ c & c' \end{array} \right) = \boxed{\left( \begin{array}{cc} a & a' \\ b & b' \\ c & c' \end{array} \right)} = \boxed{\left( \begin{array}{cc} 2 & 1 \\ 0 & -3 \\ -1 & 3 \end{array} \right)} \end{array}$$

Matlab:

```
>> x = [-1 1 2]';
>> ys = [1 1 7; 7 1 1]';
>> A = [x.^2 x x.^0]
```

Observe: Given equations

$$A \vec{x} = \vec{y}$$

$$A \vec{x}' = \vec{y}'$$

$$A \vec{x}'' = \vec{y}''$$

with the same A

A =

$$\begin{matrix} 1 & & -1 & & 1 \\ 1 & & 1 & & 1 \\ 4 & & 2 & & 1 \end{matrix}$$

`>> abcs = A \ ys`

abcs =

$$\begin{matrix} 2 & & 1 & & 1 \\ 0 & & -3 & & 1 \\ -1 & & 3 & & 1 \end{matrix}$$

they are equivalent to

$$A \left( \begin{array}{c|c|c} 1 & | & | \\ \vec{x} & | & \vec{x}' & | & \vec{x}'' \\ 1 & | & 1 & | & 1 \end{array} \right) = \left( \begin{array}{c|c|c} 1 & | & | \\ \vec{y} & | & \vec{y}' & | & \vec{y}'' \\ 1 & | & 1 & | & 1 \end{array} \right)$$

and can be solved together.

Example: Invert the matrix  $A = \begin{pmatrix} 1 & 2 & 0 \\ 3 & 2 & 1 \\ 0 & 1 & 0 \end{pmatrix}$

Answer:  $A^{-1}$  satisfies

$AA^{-1} = I$ , so we can use Gaussian elimination to solve

$$\left( \begin{array}{ccc|cc} 1 & 2 & 0 & 1 & I_1 \\ 3 & 2 & 1 & 1 & I_2 \\ 0 & 1 & 0 & 1 & I_3 \end{array} \right) \dots$$

Matlab:

```
>> inv([1 2 0; 3 2 1; 0 1 0])
ans =
1.0000 0.0000 -2.0000
0 0 1.0000
-3.0000 1.0000 4.0000
```

Python:

```
A = [[1,2,0],[3,2,1],[0,1,0]]
np.linalg.inv(A)
array([[ 1.,  0., -2.],
       [ 0.,  0.,  1.],
       [-3.,  1.,  4.]])
```

## Solving linear equations on your computer

### ① Matlab/Octave:

```
>> A = [1 1 1; 36 6 1; 4 -2 1]
```

### ② Python

```
>>> import numpy as np
>>> A = np.array([[1,1,1],[36,6,1],[4,-2,1]])
```

## (1) MATLAB/Octave:

```

>> A = [1 1 1; 36 6 1; 4 -2 1]
A =
 1   1   1
 36  6   1
  4  -2   1
>> b = [1; 5; 3]
b =
 1
 5
 3
>> x = A \ b
x =
 0.1833
-0.4833
 1.7000
(See also "format rat")
>> A*x - b
ans =
 1.0e-15 *
 0
 0.8882
 0.4441

```

## (2) Mathematica:

```

A = {{1, 1, 1}, {36, 6, 1}, {4, -2, 1}};
b = {1, 5, 3};
LinearSolve[A, b]
{11/60, -29/60, 13/10}

```

## Gaussian elimination

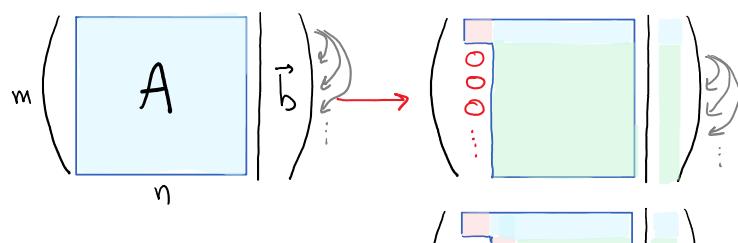
Observe:

- 1) Multiplying an equation by  $c \neq 0$   
 2) Adding one equation to another } doesn't change the solutions

$\left( \vec{x} \text{ solves eq. 1} \Leftrightarrow \vec{x} \text{ solves eq. 1 and eq. 2} \quad \text{and} \quad \vec{x} \text{ solves eq. 1 + eq. 2} \right)$

Therefore: Repeat:

Add multiples of one equation to the others  
 to eliminate one variable from them.



## (2) Python:

```

>>> import numpy as np
>>> A = np.array([[1,1,1], [36,6,1], [4,-2,1]])
>>> b = np.array([1,5,3])
>>> x = np.linalg.solve(A,b)
>>> x
array([ 0.18333333, -0.48333333,  1.3        ])
>>> np.dot(A,x) - b
array([ 0.00000000e+00,  0.00000000e+00,  4.44089210e-16])

```

```

>>> n = 800; A = np.random.randn(n,n); b = np.dot(A, np.random.randn(n));
>>> start = time.time(); x = np.linalg.solve(A,b); end = time.time(); end-start
0.029742002487182617
>>> n = 1600; A = np.random.randn(n,n); b = np.dot(A, np.random.randn(n));
>>> start = time.time(); x = np.linalg.solve(A,b); end = time.time(); end-start
0.17376208305358887
>>> n = 3200; A = np.random.randn(n,n); b = np.dot(A, np.random.randn(n));
>>> start = time.time(); x = np.linalg.solve(A,b); end = time.time(); end-start
1.0949180126190186
>>> n = 6400; A = np.random.randn(n,n); b = np.dot(A, np.random.randn(n));
>>> start = time.time(); x = np.linalg.solve(A,b); end = time.time(); end-start
7.7281270027160645

```

You still need to check  
 your answer!

(what if there's no solution?)

0

1.0e-15 \*

0.8882

0.4441

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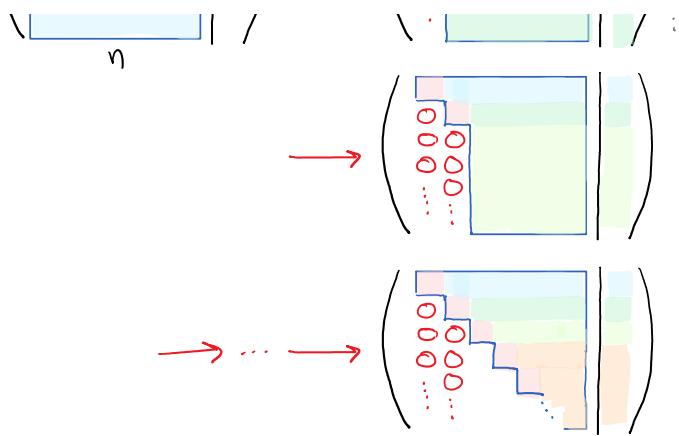
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Notes:

- You might have to rearrange the rows/equations

e.g.,  $\begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \xrightarrow{\text{row swap}} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$

- Rearranging the columns/variables can be helpful

e.g.  $\begin{pmatrix} a & b & 1 \\ c & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}$

- These are called "pivots"

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \dots + a_{1n}x_n = b_1$$

$$a'_{22}x_2 + a'_{23}x_3 + \dots + a'_{2n}x_n = b'_2$$

$$a'_{33}x_3 + \dots + a'_{3n}x_n = b'_3$$

$\vdots$

$$a'_{nn}x_n = b'_n$$

Now solve for the last variable, backsubstitute.

last equation  $\Rightarrow x_n = \frac{b'_n}{a'_{nn}}$   
 substitute into  $\Rightarrow a'_{n-1,n-1}x_{n-1} = b'_{n-1} - a'_{n-1,n} \frac{b'_n}{a'_{nn}}$   
 e.g.  $n-1$

Keep on substituting backwards, solving for variables

$x_n, x_{n-1}, x_{n-2}, \dots, x_2, x_1$

Complexity: How many basic operations are needed?

$\boxed{\mathcal{O}(n^3)}$  Why?

$\mathcal{O}(n)$  time to add a multiple of a row to another

- $\overbrace{\quad}^{\text{O}(n)}$  time to add a multiple of a row to another
- $\times \text{ O}(n)$  rows you add into
  - $\times \text{ O}(n)$  pivots (rows you add from)

Then backsubstitution takes  $\text{O}(n^2)$  steps. (Check this!)

$\therefore$  doubling  $n \Rightarrow 4 \times$  the memory (to store A)  
 $8 \times$  the running time!

### Examples:

① Use Gaussian elimination to solve

$$\begin{cases} -x_1 + 3x_2 - 2x_3 = 1 \\ -x_1 + 4x_2 - 3x_3 = 0 \\ -x_1 + 5x_2 - 4x_3 = 0 \end{cases}$$

Answer:

$$\begin{pmatrix} -1 & 3 & -2 \\ -1 & 4 & -3 \\ -1 & 5 & -4 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$$

$$\begin{array}{c} \xrightarrow{\text{pivot}} \\ \left( \begin{array}{ccc|c} -1 & 3 & -2 & 1 \\ -1 & 4 & -3 & 0 \\ -1 & 5 & -4 & 0 \end{array} \right) \xrightarrow{-1} \left( \begin{array}{ccc|c} 1 & 3 & -2 & 1 \\ 0 & 1 & -1 & -1 \\ 0 & 2 & -2 & -1 \end{array} \right) \xrightarrow{2-2} \\ \xrightarrow{\quad} \left( \begin{array}{ccc|c} 1 & 3 & -2 & 1 \\ 0 & 1 & -1 & -1 \\ 0 & 0 & 0 & 1 \end{array} \right) \end{array}$$

Stuck! 3<sup>rd</sup> equation is  $0x_1 + 0x_2 + 0x_3 = 1$   
 $\Rightarrow$  no solution.

② Solve

$$\begin{aligned} 2x_1 + x_2 &= 5 \\ x_1 + 2x_2 + 3x_3 &= 4 \\ -2x_1 + 3x_2 &= 7 \end{aligned}$$

Answer:

$$\left( \begin{array}{ccc|c} 2 & 1 & 0 & 5 \\ 1 & 2 & 3 & 4 \\ -2 & 3 & 0 & 7 \end{array} \right)$$

It is easiest to start with  $x_3$ !

$$\left( \begin{array}{ccc|c} 1 & 2 & 3 & 4 \\ 2 & 1 & 0 & 5 \\ -2 & 3 & 0 & 7 \end{array} \right) \xrightarrow{\text{Row } 3 + 2 \times \text{Row } 1} \left( \begin{array}{ccc|c} 1 & 2 & 3 & 4 \\ 2 & 1 & 0 & 5 \\ 0 & 4 & 0 & 12 \end{array} \right)$$

Now  $x_1$ ?

Next we backsubstitute

$$\left( \begin{array}{ccc|c} 1 & 2 & 3 & 4 \\ 2 & 1 & 0 & 5 \\ 0 & 1 & 0 & 3 \end{array} \right) \xrightarrow{\text{Row } 2 - 2 \times \text{Row } 1} \left( \begin{array}{ccc|c} 1 & 0 & 3 & 2 \\ 0 & 0 & 0 & 2 \\ 0 & 1 & 0 & 3 \end{array} \right) \xrightarrow{\text{Row } 3 - \text{Row } 1} \left( \begin{array}{ccc|c} 0 & 0 & 3 & -3 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 3 \end{array} \right)$$

$$\Rightarrow \vec{x} = \boxed{\begin{pmatrix} 1 \\ 3 \\ -1 \end{pmatrix}}$$

Check the answer!

$$\left( \begin{array}{ccc} 2 & 1 & 0 \\ 1 & 2 & 3 \\ -2 & 3 & 0 \end{array} \right) \begin{pmatrix} 1 \\ 3 \\ -1 \end{pmatrix} = \begin{pmatrix} 5 \\ 4 \\ 7 \end{pmatrix} \checkmark$$

## LU DECOMPOSITION OF A MATRIX

Theorem: [Turing 1948]

Any  $m \times n$  matrix  $A$  can be factored as

$$\left( \begin{array}{c} A \\ \hline m \times n \end{array} \right) = \left( \begin{array}{c} P \\ \hline m \times m \end{array} \right) \text{permutation} \left( \begin{array}{cc} L & 0 \\ \hline m \times m & m \times m \end{array} \right) \text{lower triangular} \left( \begin{array}{cc} U & 0 \\ 0 & \hline m \times m & m \times n \end{array} \right) \text{upper triangular}$$

if  $m \geq n$

Why? It is just Gaussian elimination!

e.g.,

$$\left( \begin{array}{ccc|c} 3 & 2 & 1 & 1 \\ 1.5 & 2 & 1.5 & 1.5 \\ -1 & 0 & 1 & 1 \end{array} \right) \xrightarrow{\text{Row } 2 - \frac{1}{2} \text{Row } 1} \left( \begin{array}{ccc|c} 3 & 2 & 1 & 1 \\ 0 & 1 & 1 & 1.5 \\ -1 & 0 & 1 & 1 \end{array} \right) \xrightarrow{\text{Row } 3 + \frac{1}{3} \text{Row } 1} \left( \begin{array}{ccc|c} 3 & 2 & 1 & 1 \\ 0 & 1 & 1 & 1.5 \\ 0 & 0 & \frac{4}{3} & \frac{4}{3} \end{array} \right)$$

$\gg [L, U, P] = lu([3 2 1; 1.5 2 1.5; -1 0 1])$

$L =$

$$\begin{pmatrix} 1.0000 & 0 & 0 \\ 0.5000 & 1.0000 & 0 \\ 0.3333 & 0.6667 & 1.0000 \end{pmatrix}$$

$$L = \begin{pmatrix} 1.0000 & 0 & 0 \\ 0.5000 & 1.0000 & 0 \\ -0.3333 & 0.6667 & 1.0000 \end{pmatrix}$$

$$U = \begin{pmatrix} 3.0000 & 2.0000 & 1.0000 \\ 0 & 1.0000 & 1.0000 \\ 0 & 0 & 0.6667 \end{pmatrix}$$

$$P = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

$U$  is the result of Gaussian elimination  
 $L$  and  $P$  keep track of the steps

Proof sketch:

Claim 1: Multiplying on the left by

$$\begin{pmatrix} 1 & & 0 \\ 0 & 1 & 0 \\ i & c & 1 \\ & \uparrow & \\ & j & \end{pmatrix} \text{ adds } c \text{ times row } j \text{ to row } i \quad (j < i)$$

Claim 2: The product of lower-triangular matrices is also lower triangular.

$$( \begin{array}{ccc|c} & & & 0 \\ & & & 0 \\ & & & 0 \\ \hline & & & 0 \\ A & & & \end{array} ) \left( \begin{array}{ccc|c} & & & 0 \\ & & & 0 \\ & & & 0 \\ \hline & & & 0 \\ B & & & \end{array} \right) = \left( \begin{array}{ccc|c} & & & 0 \\ & & & 0 \\ & & & 0 \\ \hline & & & 0 \\ C & & & \end{array} \right)$$

Proof:

$$A_{ik} = 0 \text{ if } i < k$$

$$B_{kj} = 0 \text{ if } k < j$$

$$C_{ij} = \sum_{k=1}^n A_{ik} B_{kj} = \sum_{k=j}^i A_{ik} B_{kj} = 0 \text{ if } i < j \quad \square$$

Matlab

```
>> rand('twister', 1)
>> A = rand(3, 5)
```

A =

$$\begin{pmatrix} 0.4170 & 0.3023 & 0.1863 & 0.5388 & 0.2045 \\ 0.7203 & 0.1468 & 0.3456 & 0.4192 & 0.8781 \\ 0.0001 & 0.0923 & 0.3968 & 0.6852 & 0.0274 \end{pmatrix}$$

```
>> [L, U, P] = lu(A);
```

>> P

>> L

>> U

P =

L =

U =

$$\begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \left( \begin{array}{ccc|c} 1.0000 & & & 0 \\ 0.5789 & 1.0000 & & 0 \\ 0 & 0.4026 & 1.0000 & 0 \end{array} \right) \begin{pmatrix} 0.7203 & 0.1468 & 0.3456 & 0.4192 & 0.8781 \\ 0 & 0.2174 & -0.0138 & 0.2961 & -0.3039 \\ 0 & 0 & 0.4026 & 0.5594 & 0.1563 \end{pmatrix}$$

Python

```
import numpy as np

# doing it like this to get the same
# randoms as Matlab
np.random.seed(1)
A = np.random.rand(5,3).T
print(A)
from scipy.linalg import lu
lu(A)
```

$$\begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1.0000 & 0 & 0 \\ 0.5789 & 1.0000 & 0 \\ 0.0002 & 0.4247 & 1.0000 \end{pmatrix} \begin{pmatrix} 0.7203 & 0.1468 & 0.3456 & 0.4192 & 0.8781 \\ 0 & 0.2174 & -0.0138 & 0.2961 & -0.3039 \\ 0 & 0 & 0.4026 & 0.5594 & 0.1563 \end{pmatrix}$$

>>  $P^T * L * U - A$

ans = *note the transpose!*

$1.0e-16 *$

$$\begin{pmatrix} 0 & 0 & 0 & 0 & -0.5551 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

How to solve the SAME system of linear equations repeatedly (with different right-hand sides)

Problem: How can we solve the equations faster?

Given a matrix  $A$ ,

and vectors  $\vec{b}, \vec{b}', \vec{b}'', \dots$ ,

Goal: Solve quickly

$$\begin{aligned} A \vec{x} &= \vec{b} \\ A \vec{x}' &= \vec{b}' \\ A \vec{x}'' &= \vec{b}'' \\ &\vdots \\ \text{same matrix } A & \quad \text{different vectors } \vec{b} \end{aligned}$$

Idea: Combine them, then use Gaussian elim.

$$A \begin{pmatrix} | & | & | \\ \vec{x} & \vec{x}' & \vec{x}'' \\ | & | & | \end{pmatrix} = \begin{pmatrix} | & | & | \\ \vec{b} & \vec{b}' & \vec{b}'' \\ | & | & | \end{pmatrix}$$

This only works if we know all  $\vec{b}$ 's in advance

Bad idea: Precompute  $A^{-1}$  (if it exists).

Then just multiply

$$\vec{x} = A^{-1} \vec{b}, \vec{x}' = A^{-1} \vec{b}', \vec{x}'' = A^{-1} \vec{b}'', \dots$$

*Don't compute inverses!*

Slow

Numerically unstable

(sparse matrix) can be dense

Application: Quickly solving multiple sets of equations with the same  $A$

Goal: Solve  $A\vec{x} = \vec{b}$   
 $\text{PLU}$

① Precompute  $P, L, U$

$\mathcal{O}(n^3)$

② Let  $\vec{y} = U\vec{x}$ .

Solve  $L\vec{y} = P^{-1}\vec{b}$  for  $\vec{y}$ .

③ Solve  $U\vec{x} = \vec{y}$  for  $\vec{x}$ .

For an  $n \times n$  matrix, this is fast!  $\mathcal{O}(n^2)$

$$\left( \begin{array}{ccc|c} 1 & 2 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right) \vec{y} = P^{-1}\vec{b} \quad \text{by forward substitution} \\ (\text{get } y_1, \text{ substitute, } y_2, y_3, \dots)$$

$$\left( \begin{array}{ccc|c} 1 & 2 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right) \vec{x} = \vec{y} \quad \text{by back substitution} \\ (\text{get } x_n, x_{n-1}, x_{n-2}, \dots)$$

∴ Precomputing the  $LL^T$  decomposition of  $A$  allows for solving  $A\vec{x} = \vec{b}$ ,  $A\vec{x}' = \vec{b}'$ ,  $A\vec{x}'' = \vec{b}''$ , ... faster.

In fact, if  $A$  is tridiagonal,  $L$  and  $U$  will be sparse while  $A^{-1}$  can be dense → even given  $A^{-1}$  for free, it is faster to solve  $A\vec{x} = \vec{b}$  using the  $LL^T$  decomposition than to multiply by  $A^{-1}$ !

## More interesting Gaussian elimination examples

Example: Solve

$$\begin{aligned} x + 3y - z &= 4 \\ 3x + 10y &= 8 \end{aligned}$$

Answer:

$$\left( \begin{array}{ccc|c} 1 & 3 & -1 & 4 \\ 3 & 10 & 0 & 8 \end{array} \right)$$

The first step is already done!

$$\left( \begin{array}{ccc|c} 1 & 3 & -1 & 4 \\ 0 & 1 & 0 & -2 \end{array} \right)$$

Now backsubstitute.

$$\left( \begin{array}{ccc|c} 1 & 3 & -1 & 4 \\ 0 & 1 & 0 & -2 \end{array} \right) \xrightarrow{-3} \left( \begin{array}{ccc|c} 0 & -2 & -1 & 4 \\ 0 & 1 & 0 & -2 \end{array} \right)$$

$$\Rightarrow x = \frac{8}{3} - \frac{10}{3}y \quad \dots \quad \dots \quad \dots$$

$$z = -\frac{4}{3} - \frac{1}{3}y \quad \text{where } y \text{ is arbitrary:}$$

$$\Rightarrow \begin{pmatrix} x \\ y \\ z \end{pmatrix} = \frac{1}{3} \begin{pmatrix} 8 \\ 0 \\ -4 \end{pmatrix} + \begin{pmatrix} -\frac{10}{3} \\ -\frac{1}{3} \\ 1 \end{pmatrix} y, \quad y \in \mathbb{R}$$

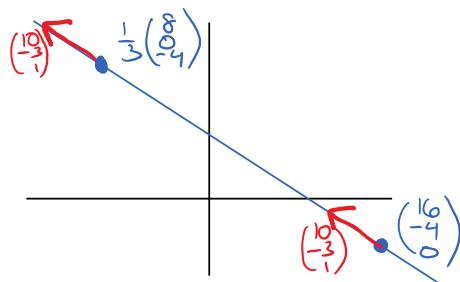
or  $\begin{pmatrix} x \\ y \\ z \end{pmatrix} = \frac{1}{3} \begin{pmatrix} 8 \\ 0 \\ -4 \end{pmatrix} + \begin{pmatrix} 10 \\ -3 \\ 1 \end{pmatrix} t, \quad t \in \mathbb{R}$

↑ "free variable"  
t = -3y

Note: You can also solve the equations with x or z free

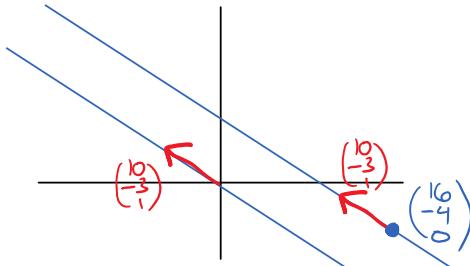
e.g.,

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} = \begin{pmatrix} 16 \\ -4 \\ 0 \end{pmatrix} + \begin{pmatrix} 10 \\ -3 \\ 1 \end{pmatrix} z, \quad z \in \mathbb{R}$$



Observe:

- ① The set  $\{(10z, -3z, z) \mid z \in \mathbb{R}\}$  is a line through  $\vec{0} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$ .  
Adding  $(16, -4, 0)$  shifts this line.



②  $\begin{pmatrix} 16 \\ -4 \\ 0 \end{pmatrix}$  solves  $\begin{cases} x + 3y - z = 4 \\ 3x + 10y = 8 \end{cases}$

$\begin{pmatrix} 10 \\ -3 \\ 1 \end{pmatrix}$  solves  $\begin{cases} x + 3y - z = 0 \\ 3x + 10y = 0 \end{cases}$  "homogeneous, equations"

Definition: A system of equations  $A\vec{x} = \vec{b}$  is

- homogeneous if  $\vec{b} = \vec{0}$
- non-homogeneous if  $\vec{b} \neq \vec{0}$

$\Rightarrow$  The general solution for a nonhomogeneous equation  $A\vec{x} = \vec{b}$

$\Rightarrow$  The general solution for a nonhomogeneous equation  $A\vec{x} = \vec{b}$   
 is a particular solution like  $\begin{pmatrix} 16 \\ -4 \end{pmatrix}$  or  $\frac{1}{3}\begin{pmatrix} 8 \\ -4 \end{pmatrix}$   
 plus the general solution to  $A\vec{x} = \vec{0}$ .  
 called the nullspace of A

This makes sense:

- If  $A\vec{x} = \vec{b}$  and  $A\vec{y} = \vec{b}$   $\Rightarrow A(\vec{x} - \vec{y}) = \vec{b} - \vec{b} = \vec{0}$
- If  $A\vec{x} = \vec{b}$  and  $A\vec{y} = \vec{0}$   $\Rightarrow A(\vec{x} + \vec{y}) = \vec{b} + \vec{0} = \vec{b}$

Example: Solve

$$\begin{cases} x_1 + 2x_2 + 2x_3 + 3x_4 = 0 \\ 2x_1 + 4x_2 + x_3 + 3x_4 = 0 \\ 3x_1 + 6x_2 + x_3 + 4x_4 = 0 \end{cases}$$

Note: Since the right-hand side is  $\vec{0}$ ,  $\vec{x} = \vec{0}$  is a solution.  
 Are there more solutions?

$$\left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 2 & 4 & 1 & 3 & 0 \\ 3 & 6 & 1 & 4 & 0 \end{array} \right) \xrightarrow{-2} \left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & -5 & -5 & 0 \end{array} \right) \xrightarrow{-\frac{5}{3}} \left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right) \xrightarrow{+1} \left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right) \xrightarrow{0=0}$$

(Note: To minimize accumulation of floating point errors, it is often best to pivot on the largest magnitude entry in the column (3))

$$x_1 + 2x_2 - x_3 = x_3 + x_4 = 0$$

$$\Rightarrow x_3 = -x_4$$

$$x_1 = -2x_2 - x_4$$

$$\Rightarrow \text{solutions} = \{(x_1, x_2, x_3, x_4) \in \mathbb{R}^4 : \begin{cases} x_1 = -2x_2 - x_4 \\ x_3 = -x_4 \end{cases}\}$$

$x_2$  &  $x_4$  are free — can be arbitrary

$$\begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} -2x_2 - x_4 \\ x_2 \\ -x_4 \\ x_4 \end{pmatrix} = x_2 \begin{pmatrix} -2 \\ 1 \\ 0 \\ 0 \end{pmatrix} + x_4 \begin{pmatrix} -1 \\ 0 \\ 1 \\ -1 \end{pmatrix}$$

— the solution set is a 2D plane in  $\mathbb{R}^4$

everything expressed

in the free variables

Example (continued): Solve

$$\left\{ \begin{array}{l} x_1 + 2x_2 + 2x_3 + 3x_4 = 1 \\ 2x_1 + 4x_2 + x_3 + 3x_4 = -1 \\ 3x_1 + 6x_2 + x_3 + 4x_4 = 0/-2 \end{array} \right. \underbrace{\quad}_{\text{same } A}$$

$$\left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 2 & 4 & 1 & 3 & -1 \\ 3 & 6 & 1 & 4 & 0 \end{array} \right) \xrightarrow{-2} \left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 0 & 0 & -3 & -3 & -3 \\ 0 & 0 & -5 & -5 & -5 \end{array} \right) \xrightarrow{-3} \left( \begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 0 & 0 & -3 & -3 & -3 \\ 0 & 0 & -5 & -5 & -5 \end{array} \right) \xrightarrow{\text{contradiction}} \Rightarrow \text{no solution!}$$

$$x_1 + 2x_2 - x_3 = -2$$

$$x_3 + x_4 = 1$$

$$\begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} -2 - 2x_2 + (1-x_4) \\ x_2 \\ 1 - x_4 \\ x_4 \end{pmatrix} = \begin{pmatrix} -1 \\ 0 \\ 1 \\ 0 \end{pmatrix} + x_2 \begin{pmatrix} -2 \\ 1 \\ 0 \\ 0 \end{pmatrix} + x_4 \begin{pmatrix} -1 \\ 0 \\ -1 \\ 1 \end{pmatrix}$$

this part is the same  
as for the homogeneous  
system

## How To SOLVE SYSTEMS OF LINEAR EQUATIONS

- Inverting the matrix  $A$  (if it is invertible)
- Gaussian elimination
  - LU decomposition
- Iterative methods
  - preconditioners
  - conjugate gradient (for positive semidefinite  $A$ )
  - multigrid methods
- Methods for solving SDD systems
  - parse diagonally dominant

More goals: \* Solve approximately instead of exactly

- \* How good is the solution? (stability)
- \* Solve the same system multiple times (different r.h.s.)
- \* Perturbed systems
- \* Systems with a special form
- \* Understand both how the algorithms work and how to use them! (and when to use them)

## ① Stability — very important

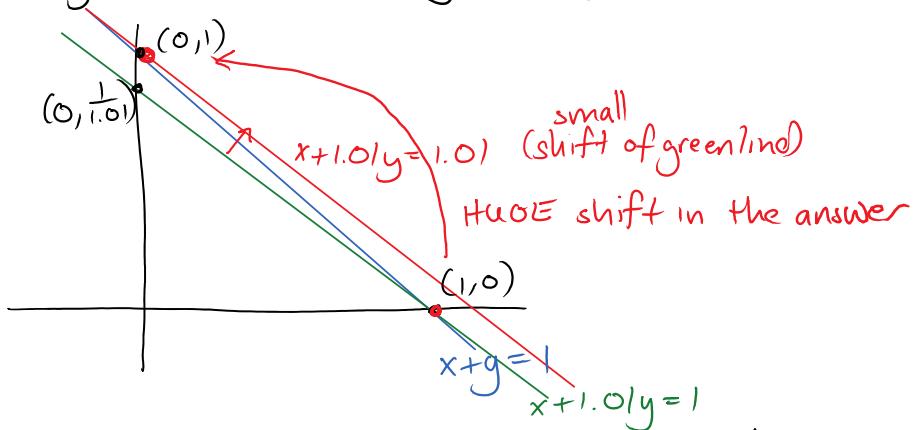
Example: Consider the equations

$$\begin{cases} x + y = 1 \\ x + 1.01y = 1 \end{cases} \Rightarrow (x, y) = (1, 0)$$

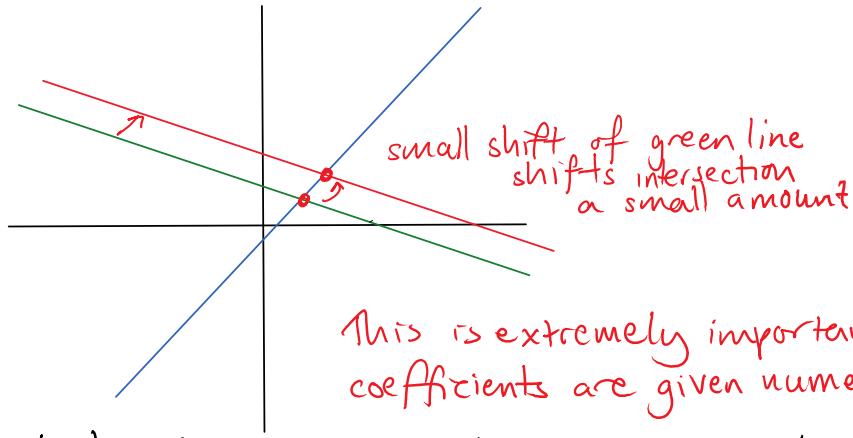
$$\begin{cases} x + y = 1 \\ x + 1.01y = 1.01 \end{cases} \Rightarrow (x, y) = (0, 1)$$

A small change to the equations led to a **HUGE** change in the solutions!

Why? Observe the geometry:



- because the lines are almost parallel  
 this wouldn't have happened if the lines were at a greater  $\angle$



This is extremely important when coefficients are given numerically.

Note: In higher dimensions, you have the same problem with planes that are nearly parallel.

But they don't have to be nearly parallel; it is more complicated than that. (Eg, think of above green and blue lines as intersections of planes in 3D with  $\Sigma z = 0$  plane. They needn't be nearly parallel!)

We'll study this more later.

## ② Linear programming

Problem: Solve a system of linear inequalities

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n &\leq b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n &\leq b_2 \\ &\vdots \end{aligned}$$

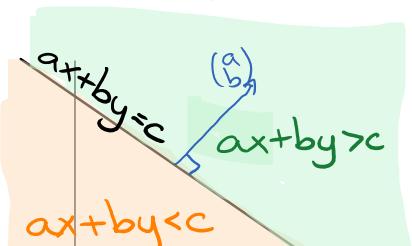
- Very useful in practice  
inequalities often express resource constraints  
eg. you have \$100  
rice is \$3/pound  
beans are \$5/pound

$$3r + 5b \leq 100$$

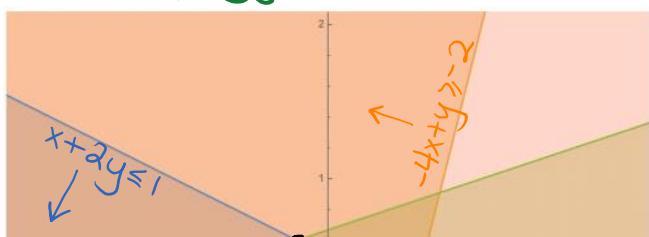
- Generalizes linear systems of equations  
 $ax+by \leq c$      $\Leftrightarrow ax+by = c$   
 $ax+by \geq c$

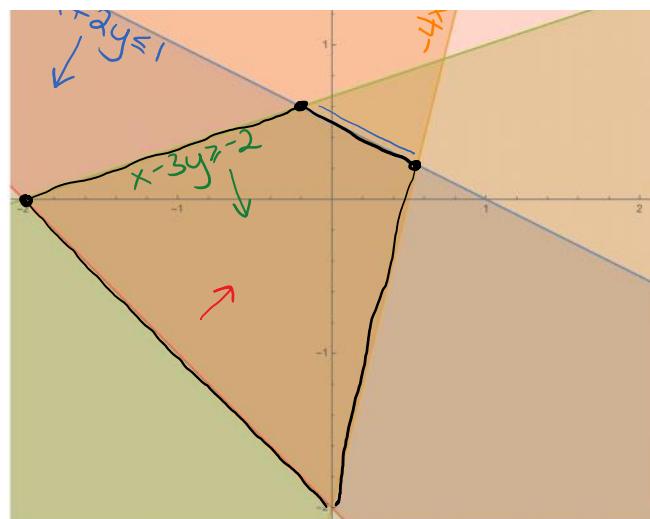
Geometry:

Halfspaces:



Convex polygon = intersection of halfspaces





Another (less interesting) application: Computing  $\text{Det}(A)$   
 $\text{Det}(A) = \text{Det}(P) \cdot \text{Det}(L) \cdot \text{Det}(U)$

Proof idea: Apply Gaussian elimination to A.

$$A = \begin{pmatrix} 2 & 2 & 2 \\ 4 & 7 & 7 \\ 6 & 18 & 22 \end{pmatrix} \xrightarrow{2 \leftrightarrow 3}$$

$$\xrightarrow{\begin{pmatrix} 2 & 2 & 2 \\ 0 & 3 & 3 \\ 6 & 18 & 22 \end{pmatrix} \xrightarrow{-3}} \xrightarrow{\begin{pmatrix} 2 & 2 & 2 \\ 0 & 3 & 3 \\ 0 & 12 & 18 \end{pmatrix} \xrightarrow{-4}}$$

$\begin{matrix} \\ \\ \text{A}_1 \\ \text{A}_2 \end{matrix}$

$$\rightarrow \begin{pmatrix} 1 & 2 & 2 \\ 0 & 3 & 3 \\ 0 & 0 & 6 \end{pmatrix} \text{ done}$$

$\overset{\text{"}}{A_1}$        $\overset{\text{"}}{A_2}$   
 $\overset{\text{"}}{A_3}$

Observe:

$$A = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} A_1$$

$$A_1 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 3 & 0 & 1 \end{pmatrix} A_2$$

$$A_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 4 & 1 \end{pmatrix} A_3$$

$$\Rightarrow A = \underbrace{\begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 3 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 4 & 1 \end{pmatrix}}_{\begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 3 & 4 & 1 \end{pmatrix} = L} \begin{pmatrix} 2 & 2 & 2 \\ 0 & 3 & 3 \\ 0 & 0 & 6 \end{pmatrix} \text{ U}$$

Why is P necessary?

$$A = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} = \begin{pmatrix} a & 0 \\ b & c \end{pmatrix} \begin{pmatrix} d & e \\ 0 & f \end{pmatrix}$$

$$\Rightarrow 0 = a \cdot d$$

$$\Rightarrow \frac{a=0}{(0 \ 0) \begin{pmatrix} d & e \\ 0 & f \end{pmatrix}} \quad \text{or} \quad \frac{d=0}{\begin{pmatrix} a & 0 \\ b & c \end{pmatrix} (0 \ 0)}$$

$$= \begin{pmatrix} 0 & 0 \\ u & u \end{pmatrix} X \quad = \begin{pmatrix} 0 & u \\ 0 & u \end{pmatrix} X$$

Problem: You can't pivot on 0.

So switch the two rows with  $P = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

$$PA = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} = \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}}_L \underbrace{\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}}_U$$

Exercise: Find the equation  $y = mx + b$  for the line through  $(1, 1)$  and  $(6, 5)$

Answer: slope  $m = \frac{5-1}{6-1} = \frac{4}{5}$   
 y-intercept  $b = \frac{1}{5}$  (check this)

Example: Solve

$$\begin{aligned} x + 3y - z &= 4 \\ 3x + 10y &= 8 \end{aligned}$$

Answer: To eliminate  $x$ , subtract 3 times the first equation from the 2<sup>nd</sup> one:

$$\Rightarrow \begin{cases} x + 3y - z = 4 \\ y + 3z = -4 \end{cases}$$

$$\Rightarrow y = -4 - 3z, \text{ where } z \text{ is arbitrary}$$

Substituting into the first equation,

$$\begin{aligned} x &= 4 - 3y + z \\ &= 4 - 3(-4 - 3z) + z \\ &= 16 + 10z \end{aligned}$$

$\Rightarrow$  Infinitely many solutions, all of the form

$$x = 16 + 10z, y = -4 - 3z, z \in \mathbb{R} \text{ arbitrary}$$

As a vector,

$$\begin{pmatrix} x \\ y \\ z \end{pmatrix} = \begin{pmatrix} 16 \\ -4 \\ 0 \end{pmatrix} + z \begin{pmatrix} 10 \\ -3 \\ 1 \end{pmatrix}, z \in \mathbb{R}$$

↑ constant coefficients      ↓ free variable