

## Lecture 23: Positive semidefinite matrices

### Eigenvalue examples (continued):

a)

```
>> n = 10;  
A = -2 * diag(ones(n,1)) + diag(ones(n-1,1),1) + diag(ones(n-1,1),-1)
```

A =

```
-2   1   0   0   0   0   0   0   0   0  
 1  -2   1   0   0   0   0   0   0   0  
 0   1  -2   1   0   0   0   0   0   0  
 0   0   1  -2   1   0   0   0   0   0  
 0   0   0   1  -2   1   0   0   0   0  
 0   0   0   0   1  -2   1   0   0   0  
 0   0   0   0   0   1  -2   1   0   0  
 0   0   0   0   0   0   1  -2   1   0  
 0   0   0   0   0   0   0   1  -2   1  
 0   0   0   0   0   0   0   0   1  -2
```

```
>> eig(A)
```

ans =

-3.9190  
-3.6825  
-3.3097  
-2.8308  
-2.2846  
-1.7154  
-1.1692  
-0.6903  
-0.3175  
-0.0810

why are these centered on -2?

b)

$$B = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

### POSITIVE SEMI-DEFINITE MATRICES

Reading: Meyer § 7.6, Strang ch. 6

Definition: A Hermitian (or real-symmetric) matrix  $A$  with all eigenvalues  $\geq 0$  is called **positive semi-definite** denoted  $A \succeq 0$ .

Key property:

Theorem: A Hermitian matrix  $A$

is positive semidefinite  
 $\iff$   
 $x^T A x \geq 0$  for all vectors  $x$ .

Proof:

$\Rightarrow$ : If  $A \succeq 0$ ,

let  $v_1, \dots, v_n$  be an orthonormal basis of eigenvectors, with corresponding e-values  $\lambda_1, \dots, \lambda_n \geq 0$ .

Any vector  $x$  can be expanded as

$$\vec{x} = \sum_{j=1}^n \alpha_j \vec{v}_j$$

where  $\alpha_j = v_j \cdot x$ .

$$\begin{aligned} \Rightarrow x^T A x &= \sum_{i,j} \alpha_i^* v_i^T \underbrace{A}_{\lambda_j v_i^T v_j} \vec{v}_j \alpha_j \\ &\quad \lambda_j v_i^T v_j = \begin{cases} \lambda_j, & \text{if } i=j \\ 0, & \text{otherwise} \end{cases} \\ &= \sum_{j=1}^n |\alpha_j|^2 \cdot \lambda_j \\ &\geq 0 \quad \checkmark \end{aligned}$$

$\Leftarrow$ : If  $A \not\succeq 0$ , ie., some eigenvalue  $\lambda < 0$ , let  $\vec{x}$  be a corresponding e-vector.

$$\begin{aligned} \Rightarrow x^T A x &= \lambda x^T x \\ &= \lambda \|x\|^2 \\ &< 0 \quad \checkmark \end{aligned}$$

□

Example 1: For any real matrix  $A$ ,  $A^T A \succeq 0$ .

Proof:

$$\begin{aligned} x^T A^T A x &= (Ax)^T (Ax) \\ &= \|Ax\|^2 \\ &\geq 0 \end{aligned}$$

$\Rightarrow$  by the theorem,  $A^T A \succeq 0$ .  $\checkmark$  □

Examples:

a)  $A = \begin{pmatrix} 3 & 4 \\ 4 & 0 \end{pmatrix}$

Claim:  $A \not\succeq 0$

Proof 1:  $\det(A) = 3 \cdot 0 - 4 \cdot 4 = -16$

$= \lambda_1, \lambda_2$   
 Since  $\det(A) < 0$  and eigenvalues are real (as  $A = A^T$ ),  
 one of the e-values is  $< 0$ .

Proof 2:

$$(1-x) A \begin{pmatrix} 1 \\ x \end{pmatrix} = 3 + \delta x$$

This is  $< 0$  if  $x = -1$  (for example).

b)  $B = \begin{pmatrix} 3 & 4 & 5 \\ 4 & 0 & 0 \\ 5 & 0 & 0 \end{pmatrix}$

Claim:  $B \not\succeq 0$

Note that  $\text{rank}(B) = 2$ , so  $\det(B) = 0 = \lambda_1 \lambda_2 \lambda_3$ .  
 One of the e-values is  $0$ , so the determinant  
 doesn't tell us anything about the others.

Proof:

$$(1, -1, 0) B \begin{pmatrix} 1 \\ -1 \\ 0 \end{pmatrix} = -5 < 0 \quad \square$$

Observe: Since  $A$  is a submatrix of  $B$  along the diagonal

$$\begin{pmatrix} 3 & 4 & 5 \\ 4 & 0 & 0 \\ 5 & 0 & 0 \end{pmatrix}$$

and  $A \not\succeq 0$ ,  $B$  can't be  $\succ 0$  either.  $\checkmark$

c)  $C = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 2 & 0 \\ 1 & 0 & 2 \end{pmatrix}$

Note:  $C = C^T$  (so e-values are real)

$\text{Rank}(C) = 2$  since  $\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} = 2 \cdot \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$

(so one e-value is  $0$ , and  $\det(C) = 0$ )

Claim:  $C \not\succeq 0$ .

`>> C = [1 1 1; 1 2 0; 1 0 2];  
>> eig(C)`

`ans =`

`- - - - -`

$\leftarrow$  this is not a normal

ans =

-0.0000  
2.0000  
3.0000

← this is not a proof!

Proof:

$$\begin{aligned} C &= \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} + \begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & -1 \\ 0 & -1 & 1 \end{pmatrix} \\ &= \underbrace{\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}}_{\vec{x}} \underbrace{(1 \ 1 \ 1)}_{\vec{x}^T} + \underbrace{\begin{pmatrix} 0 \\ -1 \\ 1 \end{pmatrix}}_{\vec{y}} \underbrace{(0 \ -1 \ 1)}_{\vec{y}^T} \end{aligned}$$

$$= \vec{x}\vec{x}^T + \vec{y}\vec{y}^T$$

Recall:

For any real matrix  $A$ ,  $A^T A \succeq 0$ .

$\Rightarrow \vec{x}\vec{x}^T \succeq 0$  and  $\vec{y}\vec{y}^T \succeq 0$

$\Rightarrow \vec{x}\vec{x}^T + \vec{y}\vec{y}^T \succeq 0$

Claim: If  $A \succeq 0$  and  $B \succeq 0$ , then  $A+B \succeq 0$ .

Proof:

For any vector  $\vec{v}$ ,

$$\vec{v}^T (A+B) \vec{v} = \underbrace{\vec{v}^T A \vec{v}}_0 + \underbrace{\vec{v}^T B \vec{v}}_0 \quad (\text{since } A \succeq 0, B \succeq 0)$$

$\succeq 0$

$\Rightarrow A+B \succeq 0$

□

## SUMMARY:

Definition:

Symmetric matrix  $A$  is positive definite

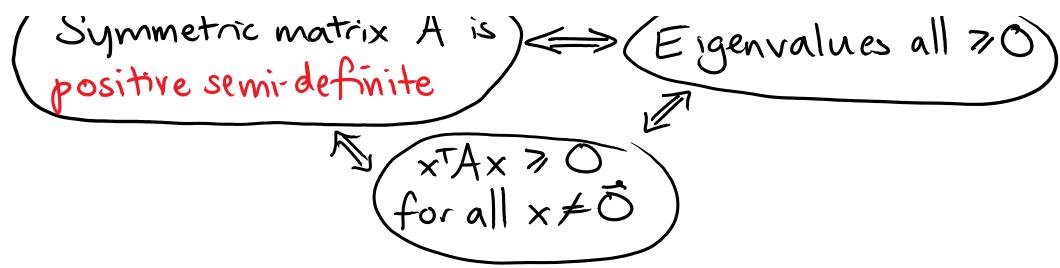
Eigenvalues all  $> 0$

$\vec{x}^T A \vec{x} > 0$   
for all  $\vec{x} \neq \vec{0}$

we proved this!

Symmetric matrix  $A$  is positive semi-definite

Eigenvalues all  $\geq 0$



Ways to show that  $A \geq 0$

- compute all e-values, check  $\geq 0$  (lame)
- write  $A = B + C$  with  $B \geq 0$  and  $C \geq 0$  (or  $B + C + D$ , etc)
- show that  $x^T A x \geq 0$  for all  $x$

Maybe also consider  $\text{rank}(A), \dots$

Ways to show  $A \not\geq 0$

- find an e-value  $< 0$
- find  $\vec{x} \neq \vec{0}$  so  $x^T A x < 0$
- show that a submatrix of  $A$  is not  $\geq 0$
- check  $\text{Det}(A) < 0$

$$\sum_{j=1}^n d_j$$

Exercise: Which of these matrices is p.s.d.?

$$A = \begin{pmatrix} 3 & 4 & 5 & 0 & 1 \\ 4 & 2 & 1 & 0 & 1 \\ 5 & 1 & 10 & 2 & 1 \\ 0 & 0 & 2 & 2 & 1 \\ 1 & 1 & 1 & 2 & 3 \end{pmatrix}$$

No!  $A^T \neq A$

$$B = \begin{pmatrix} 3 & 1 & 1 & 0 & 1 \\ 1 & 2 & 1 & 0 & 1 \\ 1 & 1 & -1 & 2 & 1 \\ 0 & 0 & 2 & 2 & 1 \\ 1 & 1 & 1 & 1 & 3 \end{pmatrix}$$

No!  $\vec{e}_3^T B \vec{e}_3 = -1 < 0$

$$C = \begin{pmatrix} 5 & 2 & 0 & 2 & 1 \\ 2 & 6 & 1 & 2 & 0 \\ 0 & 1 & 5 & 0 & 0 \\ 2 & 2 & 0 & 6 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{pmatrix}$$

No! Observe the submatrix  
 $C' = \begin{pmatrix} 5 & 1 \\ 1 & 0 \end{pmatrix}$   $\text{Det} C' = -1 < 0$

$$\text{or: } (-x-1)C' \begin{pmatrix} -x \\ 1 \end{pmatrix} = 5x^2 - 2x = x(5x-2) < 0 \text{ for } 0 < x < \frac{2}{5}$$

$$D = \begin{pmatrix} 8 & -2 & 3 & 1/2 & 0 \\ -2 & 6 & 1 & 1 & -1 \\ 3 & 1 & 10 & 2 & 1 \\ 1/2 & 1 & 2 & 5 & 1 \\ 0 & -1 & 1 & 2 & 7 \end{pmatrix}$$

Yes!  
 $D = D^T \checkmark$

Also  $D$  is diagonally dominant:  $\forall i, D_{ii} > \sum_{j \neq i} |D_{ij}|$ .  
 Assume for contradiction  $D\vec{v} = \lambda\vec{v}$  with  $\lambda < 0, \vec{v} \neq \vec{0}$ .

Let  $i = \arg\max |v_j|$ .

We may assume that  $v_i > 0$ . (If  $v_i < 0$ , multiply by -1)

$$\begin{aligned} (\vec{D}\vec{v})_i &= D_{ii}v_i + \sum_{j \neq i} D_{ij}v_j \\ &\geq D_{ii}v_i - \sum_{j \neq i} |D_{ij}| \cdot |v_j| \\ &\geq v_i \cdot \left( D_{ii} - \sum_{j \neq i} |D_{ij}| \right) \\ &> 0. \quad \times \text{ contradiction} \quad \square \end{aligned}$$

WHY CARE? one reason: optimization

Recall: A function  $T: V \rightarrow W$  is **linear** if

- $T(x+y) = T(x) + T(y)$  for all  $x, y \in V$ ,
- $T(\alpha x) = \alpha T(x)$  for all scalars  $\alpha$ .

Linear functions  $\longleftrightarrow$  Matrices  
w/ fixed bases  
for  $V$  &  $W$

Not all functions are linear... but calculus lets us give **affine approximations** to differentiable functions.

Calculus  $\rightarrow$  linearity

$$f(x) - f(x_0) \approx f'(x_0) \cdot (x - x_0)$$

for  $f: \mathbb{R} \rightarrow \mathbb{R}$  differentiable

Or for  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$  differentiable,

$$\vec{f}(\vec{y}) - \vec{f}(\vec{y}_0) \approx \frac{\partial \vec{f}}{\partial x_1}(\vec{y}_0) \cdot (\vec{y} - \vec{y}_0)_1 + \dots + \frac{\partial \vec{f}}{\partial x_n}(\vec{y}_0) \cdot (\vec{y} - \vec{y}_0)_n$$

$$= \begin{pmatrix} \frac{\partial \vec{f}}{\partial x_1} & \frac{\partial \vec{f}}{\partial x_2} & \dots & \frac{\partial \vec{f}}{\partial x_n} \end{pmatrix} \begin{pmatrix} \vec{y} - \vec{y}_0 \end{pmatrix}$$

$$\sum_{j=1}^n \frac{\partial \vec{f}}{\partial x_j} \cdot \vec{e}_j^\top$$

But often, an affine approximation isn't enough!

Eg., studying local optima requires a quadratic approx:

$$f(x) = f(x_0) + f'(x_0) \cdot (x - x_0) + \frac{1}{2} f''(x_0) \cdot (x - x_0)^2 + O((x - x_0)^3)$$

### Single-variable functions

local minimum  $\leftrightarrow f'(x_0) = 0, f''(x_0) > 0$

local maximum  $\leftrightarrow f'(x_0) = 0, f''(x_0) < 0$   
(if  $f''(x_0) \neq 0$ )

### Multi-variable functions $f: \mathbb{R}^n \rightarrow \mathbb{R}$

$$\begin{aligned} f(y) = f(y_0) + \sum_j \frac{\partial f}{\partial y_j}(y_0) \cdot (y - y_0)_j \\ + \frac{1}{2} \sum_{j,k} \frac{\partial^2 f}{\partial y_j \partial y_k}(y_0) (y - y_0)_j (y - y_0)_k \\ + O(\|y - y_0\|^3) \end{aligned}$$

For  $f: \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $y_0$  is a local minimum if

$$\frac{\partial f}{\partial y_j}(y_0) = 0 \quad \forall j = 1, \dots, n$$

and  $\underbrace{\frac{1}{2} \sum_{j,k} \frac{\partial^2 f}{\partial y_j \partial y_k}(y_0) (y - y_0)_j (y - y_0)_k}_{> 0}$

This is a **quadratic form** in the variables  
 $(y - y_0)_1, \dots, (y - y_0)_n$

**Definition:** A **quadratic form** is a degree-two **homogeneous polynomial**.

$$x^2 + 2xy + y^2 \quad \checkmark$$

$$x^2 + 2x \quad \times$$

$$x^2 - 2xz + 2y^2 - 2yz + z^2 \quad \checkmark$$

$$x^2 + 2xy + y \quad \times$$

**Observe:** Quadratic forms

↑  
Symmetric matrices

$$x^2 + 2xy + y^2 = (x \ y) \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$$

It also equals  $(x \ y) \begin{pmatrix} 1 & 2 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$  and  $(x \ y) \begin{pmatrix} 1 & 3/2 \\ 1/2 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$ ,  
but the symmetric form is nicer.

**Corollary:** Symmetric matrices let us study **homogeneous quadratic functions**.

### Examples:

- $$f(x,y) = x^2 + (2b)xy + y^2$$

$$= (x \ y) \underbrace{\begin{pmatrix} 1 & b \\ b & 1 \end{pmatrix}}_A \begin{pmatrix} x \\ y \end{pmatrix}$$

$$\lambda_1 + \lambda_2 = \text{Tr}(A) = 2$$

$$\lambda_1 \cdot \lambda_2 = \text{Det}(A) = 1 - b^2$$

$\Rightarrow A \succ 0$  if  $b < 1$

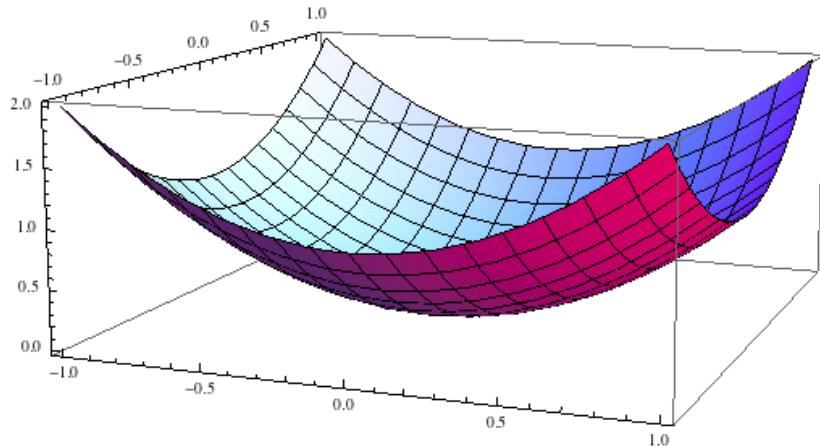
$A \succcurlyeq 0$  if  $b \leq 1$

$A \not\succ 0$  if  $b > 1$

( $A$  is indefinite)

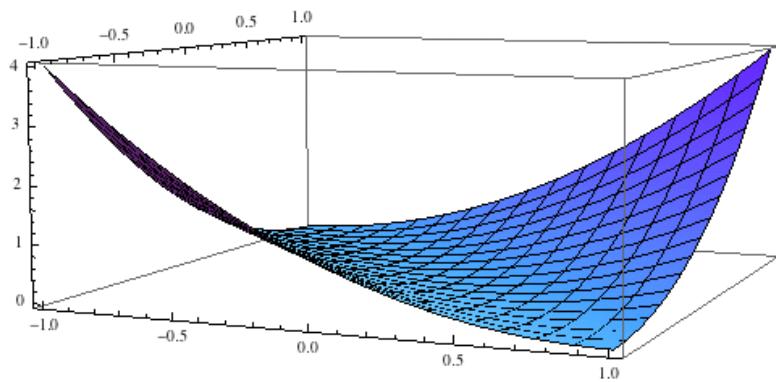
$b = 0;$

`Plot3D[x^2 + 2 b x y + y^2, {x, -1, 1}, {y, -1, 1}]`

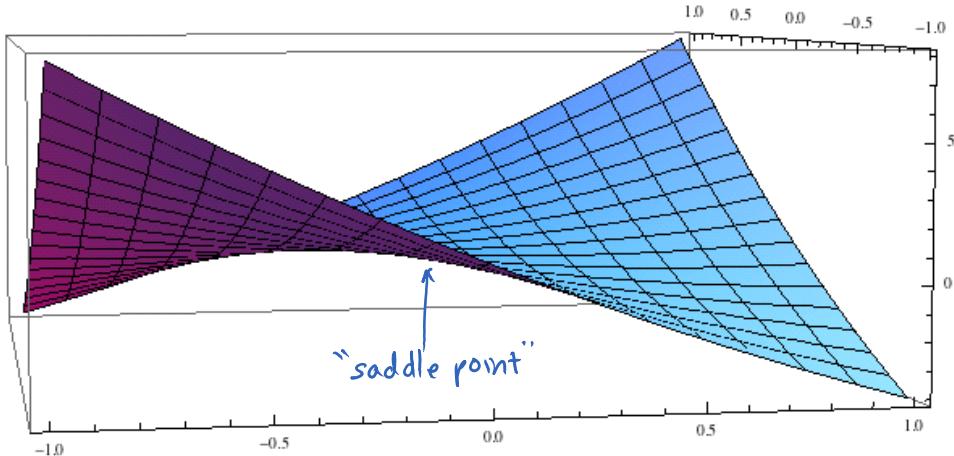


$b = 1;$

`Plot3D[x^2 + 2 b x y + y^2, {x, -1, 1}, {y, -1, 1}]`



```
b = 3;
Plot3D[x^2 + 2 b x y + y^2, {x, -1, 1}, {y, -1, 1}]
```



$$\begin{pmatrix} 1 & b \\ b & 1 \end{pmatrix} \quad \begin{matrix} \text{Eigenvalues} \\ 1+b \\ 1-b \end{matrix} \quad \begin{matrix} \text{Eigenvectors} \\ (1, 1) \\ (1, -1) \end{matrix}$$

$$\begin{pmatrix} 0 & b \\ b & 0 \end{pmatrix} + I \quad \begin{matrix} \\ \parallel \\ \end{matrix}$$

$$\Rightarrow \begin{pmatrix} 1 & b \\ b & 1 \end{pmatrix} = (1+b) \cdot \frac{1}{2}(1)(1, 1) + (1-b) \cdot \frac{1}{2}(-1)(1, -1)$$

$$\begin{aligned} \Rightarrow f(x, y) &= (x, y) \begin{pmatrix} 1 & b \\ b & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = x^2 + y^2 + 2bx \\ &= (1+b) \cdot \left( \frac{1}{2}(1, 1) \cdot (x, y) \right)^2 \\ &\quad + (1-b) \cdot \left( \frac{1}{2}(1, -1) \cdot (x, y) \right)^2 \\ &= \frac{1+b}{2} (x+y)^2 + \frac{1-b}{2} (x-y)^2 \end{aligned}$$

Observe: Any quadratic form can be expressed as a sum of squares like this, and it is  $\geq 0$  iff all coefficients are  $\geq 0$ .

By writing  $A = \lambda_1 v_1 v_1^T + \dots + \lambda_n v_n v_n^T$ ,

$$x^T A x = \sum_j \lambda_j (v_j \cdot x)^2 \quad \checkmark$$

Claim:  $A \succeq 0 \iff A = M^T M$  for some matrix  $M$ .

Proof:

$$\begin{aligned} \text{If } A = M^T M, \quad x^T A x &= \|Mx\|^2 \geq 0 \\ \Rightarrow A \succeq 0. \quad \checkmark \end{aligned}$$

$$\text{If } A \succeq 0, \quad A = \sum_j \lambda_j v_j v_j^T.$$

$$\text{Let } M = \left( \frac{\sqrt{\lambda_j} v_j^T}{\lambda_j} \right) \rightarrow M^T M = A \quad \checkmark$$

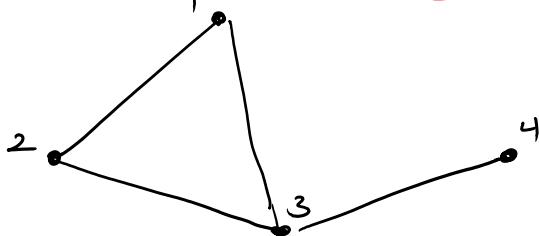
If  $A \succcurlyeq 0$ ,  $A = \sum_j \lambda_j v_j v_j^T$ .

$$\text{Let } M = \begin{pmatrix} \sqrt{\lambda_1} v_1^T \\ \sqrt{\lambda_2} v_2^T \\ \vdots \end{pmatrix} \Rightarrow M^T M = A \quad \square$$

(Remark: This also proves that  $A \succcurlyeq 0$  if and only if  $A$  is the Gram matrix for a set of  $n$  vectors.)

Example:

### The Laplacian of a graph



$$L = \begin{pmatrix} & 1 & 2 & 3 & 4 \\ 1 & 2 & -1 & -1 & 0 \\ 2 & -1 & 2 & -1 & 0 \\ 3 & -1 & -1 & 3 & -1 \\ 4 & 0 & 0 & -1 & 1 \end{pmatrix}$$

degrees along  
 edges  
 diagonals  
 edges

$L = E E^T$  for the incidence matrix  $E$

$$E = \sum_{\substack{\text{edges} \\ (u,v)}} (e_u - e_v) e_{(u,v)}^T$$

$$\Rightarrow L \succcurlyeq 0$$

$$x^T L x = \sum_{\substack{\text{edges} \\ (u,v)}} (x_u - x_v)^2$$

Question: Is  $L$  positive definite? (not just  $\succcurlyeq 0$ )

Answer: No!  $L(\mathbf{1}) = 0$ .

Observe: For  $f: \mathbb{R}^n \rightarrow \mathbb{R}$ , a "critical point" (where all derivatives  $\frac{\partial f}{\partial x_i} = 0$ ) is a local minimum if

$$\begin{pmatrix} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} & \dots \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2^2} & \ddots \\ \vdots & \ddots & \ddots \end{pmatrix}$$

is positive definite.

Geometrically:



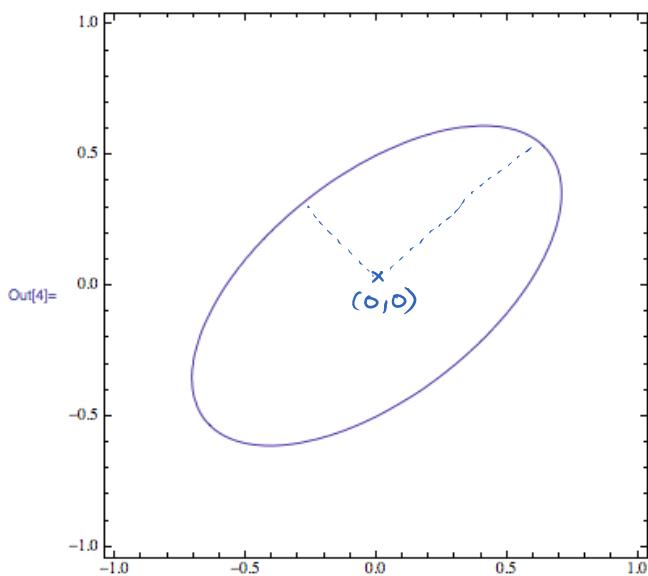
## Geometrically:

Quadratic forms  
assoc. to  
positive definite matrices

Ellipses  
and  
ellipsoids

```
In[1]:= A = {{3, -2}, {-2, 4}};
Eigenvalues[A // N]
Eigensystem[A] // N // Transpose // MatrixForm
ContourPlot[{x, y}.A.{x, y} == 1, {x, -1, 1}, {y, -1, 1}]
Out[2]= {5.56155, 1.43845}
```

```
Out[3]//MatrixForm=
{{5.56155, {-0.780776, 1.}}, {1.43845, {1.28078, 1.}}}
```

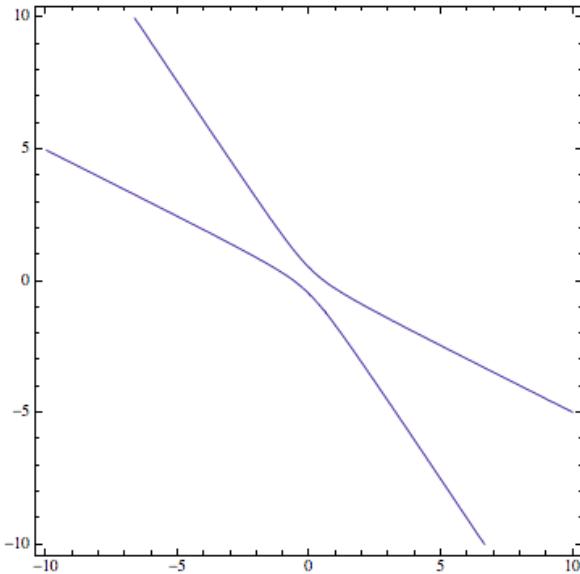


ratio  $\frac{\text{longest axis}}{\text{shortest axis}}$   
 $\longleftrightarrow$  condition number

```

A = {{3, 4}, {4, 4}};
Eigenvalues[A // N]
ContourPlot[{x, y}.A.{x, y} == 1, {x, -10, 10}, {y, -10, 10}]
{7.53113, -0.531129}

```



Ellipsoids also correspond to inner products/norms where some directions are weighted differently than others.

### Example: GRADIENT DESCENT

Goal: Solve  $A\vec{x} = \vec{b}$ .

Recall: If  $A$  is  $n \times n$ , Gaussian elimination solves for  $\vec{x}$  using  $O(n^3)$  basic operations.

Can we approximately solve for  $\vec{x}$  faster?

Define the cost function

$$\begin{aligned}
 C(\vec{x}) &= \|A\vec{x} - \vec{b}\|^2 \\
 &= \|A(\vec{x} - \vec{x}^*)\|^2 \quad \text{where } A\vec{x}^* = \vec{b} \\
 &= (\vec{x} - \vec{x}^*)^T A^T A (\vec{x} - \vec{x}^*)
 \end{aligned}$$

positive semi-definite

Algorithm: Repeat until convergence:

$$\vec{x}^t = \vec{x}^{t-1} - \alpha \left( \frac{\partial C}{\partial x_1}(\vec{x}^{t-1}), \dots, \frac{\partial C}{\partial x_n}(\vec{x}^{t-1}) \right)$$

i.e., step in the direction opposite the largest increase in  $C$ .

$$\frac{\partial C}{\partial x_j} = \frac{\partial}{\partial x_j} [(A\vec{x} - \vec{b})^T (A\vec{x} - \vec{b})]$$

$$\begin{aligned}
&= \frac{\partial}{\partial x_j} \sum_i (Ax - b)_i^2 \\
&= 2 \sum_i (Ax - b)_i \cdot \frac{\partial}{\partial x_j} [(Ax - b)_i] \\
&= 2 \sum_i (Ax - b)_i \cdot a_{ij} \\
\Rightarrow \left( \frac{\partial C}{\partial x_1}, \dots, \frac{\partial C}{\partial x_n} \right) &= 2 \sum_i \underbrace{(Ax - b)_i}_{(Ax - b)^T e_i} \cdot \underbrace{(a_{i1}, \dots, a_{in})}_{e_i^T A} \\
&= 2(Ax - b)^T (\sum_i e_i e_i^T) A \\
&= 2(Ax - b)^T A \\
&\quad (\text{As you might guess!})
\end{aligned}$$

Example:

```

A = {{2, -1.5}, {-1.5, 4}};
b = A.{3, 3};
Eigenvalues[A]

α = .01;
x = {0, -4};
list = {x};
For[k = 1, k ≤ 1000, k++,
  x = x - 2 α (A.x - b).A;
  AppendTo[list, x];
];
"The solution is:"
x
{4.80278, 1.19722}

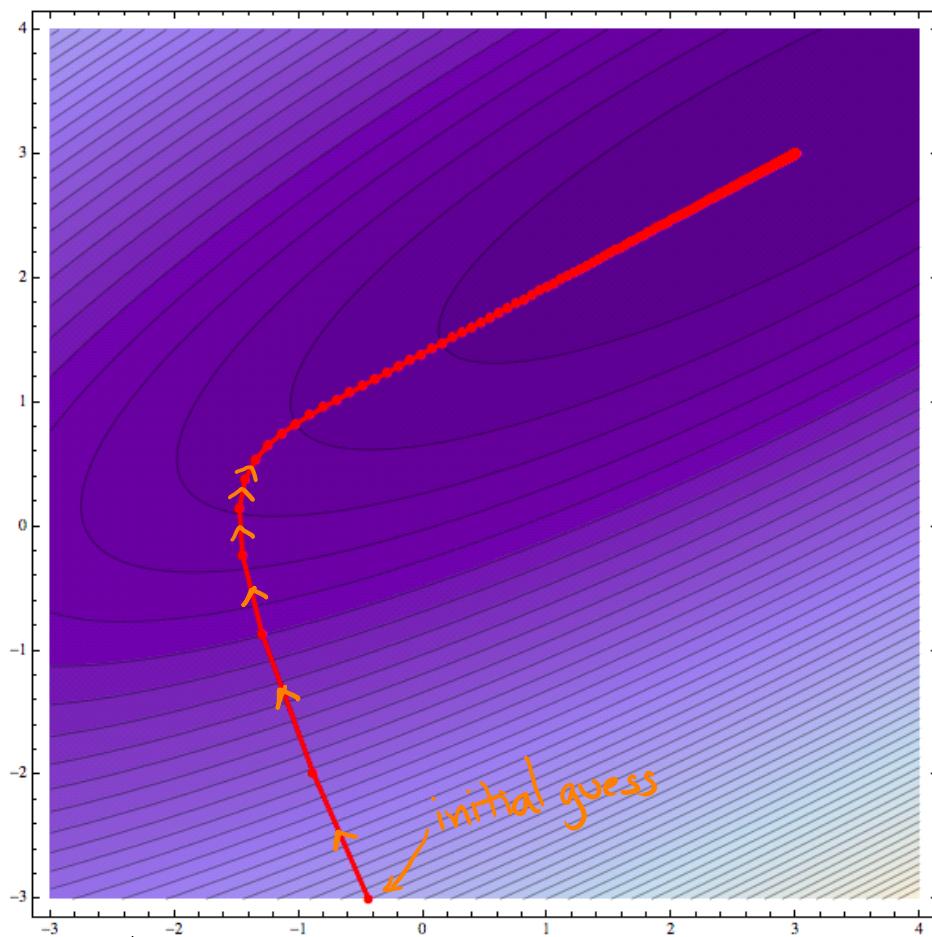
The solution is:
{3., 3.}

```

```

"Display the results";
iter = ListPlot[list, PlotRange -> 3 {{-1, 1}, {-1, 1}}, Joined -> True,
  PlotStyle -> {Red, Thickness[.005]}, PlotMarkers -> Automatic];
contour = ContourPlot[Norm[A.{x, y} - b]^2, {x, -3, 4}, {y, -3, 4}, Contours -> 50];
Show[contour, iter]

```



For small enough  $\alpha$ , this has to converge, since there is a unique local and global minimum.

Observe: Smaller condition number  
 $\Rightarrow$  Rounder ellipse  $\Rightarrow$  Faster convergence

Remark: Although not generally the fastest way of solving a set of linear equations, the gradient descent approach is very robust, and many variants are used in many applications.

Convergence analysis:

Let  $x^*$  solve  $Ax^* = b$ .

$$x_t = x_{t-1} - 2\alpha A^T(Ax_{t-1} - b)$$

$$\Rightarrow x_t - x^* = (x_{t-1} - x^*) - 2\alpha A^T A(x_{t-1} - x^*)$$

$$= (\mathbb{I} - 2\alpha A^T A)(x_{t-1} - x^*)$$

$\Rightarrow$  error magnitude

$$\|\vec{x}_t - \vec{x}^*\| \leq \|\mathbb{I} - 2\alpha A^T A\| \cdot \|x_{t-1} - x^*\|$$

If the e-values of  $A^T A$  are

$$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n \geq 0$$

$\uparrow$  these are (singular values of  $A$ )<sup>2</sup>

$\mathbb{I} - 2\alpha A^T A$  has e-values

$$1 - 2\alpha \lambda_1 \leq \dots \leq 1 - 2\alpha \lambda_n$$

$$\Rightarrow \|\mathbb{I} - 2\alpha A^T A\| = \max_i |1 - 2\alpha \lambda_i|$$

$$\max \{ |1 - 2\alpha \lambda_1|, |1 - 2\alpha \lambda_n| \}$$

The best choice of  $\alpha$  will make these equal in magnitude, centered on 0:

$$\Rightarrow -(1 - 2\alpha \lambda_1) = + (1 - 2\alpha \lambda_n)$$

$$\Rightarrow \alpha = \frac{1}{\lambda_1 + \lambda_n}$$

With this choice for  $\alpha$ ,

$$\begin{aligned} \|\mathbb{I} - 2\alpha A^T A\| &= 1 - \frac{2\lambda_n}{\lambda_1 + \lambda_n} \leq 1 - \frac{\lambda_n}{\lambda_1} \\ &= 1 - \frac{1}{\text{condition\# of } A^T A} \\ &= 1 - \frac{1}{(\text{condition\# of } A)^2} \end{aligned}$$

$$\Rightarrow \|x^t - x^*\| \leq (1 - \frac{1}{\kappa^2})^t \cdot \|x^0 - x^*\|$$

$$\Rightarrow t = \kappa^2 \left( \log \frac{1}{\epsilon} + \log \frac{1}{\|x^0 - x^*\|} \right) \quad \text{since } (1 - \frac{1}{\kappa^2})^{\kappa^2} \leq e^{-1}$$

ensures  $\|x^t - x^*\| \leq \epsilon$ .

Remark: The conjugate gradient algorithm has  $O(\sqrt{\kappa})$  dependence instead of  $O(\kappa^2)$  dependence, and

$$x^t = x^{t-1} - \alpha (Ax^{t-1} - b)$$

has  $O(K)$  dependence if  $A \geq \tilde{\sigma}^{\text{error}}$

One more motivation for positive semi-definite matrices:  
Semi-definite programming

This was touched on in the homework ...

SDPs  $\max \text{Trace}(AB)$  are everywhere!

$$\text{s.t. } A \geq 0$$

They can be solved efficiently and have a notion of duality  
roughly because the set of p.s.d. matrices is convex.