

It's now time to give our formal language a semantics.

The meaning of not:

p	$\neg p$
T	F
F	T

The meaning of or:

p	q	$(p \vee q)$
T	T	T
T	F	T
F	T	T
F	F	F

The meaning of and:

p	q	$(p \wedge q)$
T	T	T
T	F	F
F	T	F
F	F	F

The meaning of iff:

p	q	$(p \equiv q)$
T	T	T
T	F	F
F	T	F
F	F	T

How do we account for the meaning of if then?

Like so

p	q	$(p \supset q)$
T	T	T
T	F	F
F	T	T
F	F	T

Three explanations.

$$(p \supset q) = \neg p \vee q.$$

$(p \supset q)$ = it's false that p is true but q is false; true otherwise = $\neg(p \wedge \neg q)$

And also this:

We want to $p \supset q$ to license the following two inferences.

Modus Ponens

1. $P \supset Q$
2. P
3. So Q

Modus Tollens

1. $P \supset Q$
2. $\neg Q$
3. So $\neg P$

But we don't want our analysis of *if then* to license the following inference.

Affirming the consequent (No good)

1. $P \supset Q$
2. Q
3. So P

How do we set up the truth tables for $P \supset Q$ such that we get MP, MT, but not Affirming the consequent?

For MP we need a truth table that looks like this.

p	q	($p \supset q$)
T	T	T
T	F	F
F	T	?
F	F	?

Whenever ($p \supset q$) is true and p is true, we want q to also be true.

So we set ($p \supset q$) to be true when p is true and q is true, and false when p is true but q is false.

For MT we need a truth table that looks like this.

p	q	($p \supset q$)
T	T	?
T	F	F
F	T	?
F	F	T

Whenever ($p \supset q$) is true and q is false, we want p to be false.

So for p is true and q is false, we want ($p \supset q$) to be false. And for p is false and q is false, we want ($p \supset q$) to be true.

So any two-place operator that licenses MP and MT will have the following truth table.

p	q	($p \supset q$)
T	T	T
T	F	F
F	T	?
F	F	T

So the question becomes: what do we say when p is false but q is true?

If we let ($p \supset q$) be false when both p and q are false, then the only time ($p \supset q$) and q are both true, p is true. But that's the inference *affirming the consequent*, which is no good.

Notice also that defining ($p \supset q$) in this way will make ($p \supset q$) equivalent to ($p \equiv q$).

p	q	($p \supset q$) (not correct)
T	T	T
T	F	F
F	T	F
F	F	T

So to avoid having ($p \supset q$) license ($p \supset q$), q, therefore p, we have to make it so that when p is false and q is true, ($p \supset q$) is true.

p	q	($p \supset q$)
T	T	T
T	F	F
F	T	T
F	F	T

On this second way of doing things, when ($p \supset q$) is true and q is true, it's not forced that p is also true.

p	q	($p \supset q$)
T	T	T
T	F	F
F	T	T
F	F	T

Hence the above is the correct semantics for \supset .