# 代码库

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### 1 数论

#### 1.1 中国剩余定理

```
typedef long long LL;
LL a[N], r[N], n;
void exgcd(LL a, LL b, LL &d, LL &x, LL &y) {
        if (!b) \{d = a; x = 1; y = 0;\}
        else {exgcd(b, a%b, d, y, x); y = (a / b) * x;}
}
LL ex_CRT(LL *m, LL *r, int n) {
        LL M = m[1], R = r[1], x, y, d;
        for(int i = 2; i <= n; i ++) {
                exgcd(M, m[i], d, x, y);
                if ((r[i] - R) % d) return -1;
                x = (r[i] - R) / d * x % (m[i] / d);
                R += x * M;
                M = M / d * m[i];
                R \% = M;
        }
        return R > 0 ? R : R + M;
}
1.2 FFT
struct comp {
        double r, i;
        comp(double r = 0.0, double i = 0.0) : r(r), i(i) {}
        comp operator + (const comp &b) const {return comp(r + b.r, i + b.i);}
        comp operator - (const comp &b) const {return comp(r - b.r, i - b.i);}
        comp operator * (const comp &b) const
                {return comp(r * b.r - i * b.i, r * b.i + i * b.r);}
}a[N], b[N], c[N];
void FFT(comp *a, int n, int type) {
        for(int i = 1; j = 0; i < n - 1; i ++) {
                for(int s = n; j ^= s >>= 1, ~j&s;);
                if (i < j) swap(a[i], a[j]);</pre>
        }
        for(int m = 1; m < n; m <<= 1) {
                double u = pi / m * type;
                comp wm(cos(u), sin(u));
                for(int i = 0; i < n; i += (m << 1)) {
                        comp w(1, 0);
```

```
for(int j = 0; j < m; j ++) {
                                comp \&A = a[i + j + m], \&B = a[i + j], t = w * A;
                                A = B - t; B = B + t; w = w * wm;
                        }
                }
        }
        if (type == -1) for(int i = 0; i < n; i ++) a[i].r /= n;
}
     组合数学(卢卡斯定理、线性筛逆元)
//Lucas
int C(int n, int m) {
        if (n < m) return 0;</pre>
        if (n < P \&\& m < P) return fac[n] * inv[m] % P * inv[n - m] % P;
        return C(n % P, m % P) * C(n / P, m / P);
}
//inv
for(int i = 1; i <= n; i ++)
        inv[i] = (LL)inv[P % i] * (P - P / i) % P;
1.4 辛普森自适应积分
double getf(double x) {
}
double calc(double 1, double f1, double fimd, double fr) {
        return (fl + fmid * 4 + fr) * 1 / 6;
}
double simpson(double 1, double mid, double r, double fl, double fmid, double fr, double s) {
        double m1 = (1 + mid) / 2, m2 = (mid + r) / 2;
        double f1 = getf(m1), f2 = getf(m2);
        double g1 = calc(mid - 1, f1, f1, fmid), g2 = calc(r - mid, fmid, f2, fr);
        if (fabs(g1 + g2 - s) < eps) return g1 + g2;
        return simpson(1, m1, mid, f1, f1, mid, g1) +
                   simpson(mid, m2, r, fmid, f2, fr, g2);
}
1.5
     线性筛
int tot, prime[N], mu[N], sum[N];
bool check[N];
void getmu() {
        mu[1] = 1;
        for(int i = 2; i < n; i ++) {
```

```
if (!check[i]) {
                         prime[++ tot] = i;
                         mu[i] = -1;
                 }
                 for(int j = 1; j <= tot; j ++) {</pre>
                         if (i * prime[j] > n) break;
                         check[i * prime[j]] = 1;
                         if (i % prime[j]) mu[i * prime[j]] = -mu[i];
                         else {
                                  mu[i * prime[j]] = 0;
                                  break;
                         }
                 }
        }
        for(int i = 1; i < n; i ++) sum[i] = sum[i - 1] + mu[i];</pre>
}
```

## 2 图论

#### 2.1 Tarjan

```
//SCC
int dfn[N], low[N], dfs_clock, belong[N], SCC, size[N];
int st[N], top;
bool in[N];
void tarjan(int x) {
        dfn[x] = low[x] = ++dfs_clock;
        st[top++] = x;
        in[x] = 1;
        for(int i = head[x]; i; i = nxt[i])
                if (!dfn[to[i]]) {
                        tarjan(to[i]);
                        low[x] = min(low[x], low[to[i]]);
                } else if (in[to[i]]) low[x] = min(low[x], dfn[to[i]]);
        if (low[x] == dfn[x]) {
                SCC ++; size[SCC] = 0;
                for(int y = 0; y != x;) {
                        y = st[--top];
                        in[y] = 0;
                        belong[y] = num;
                        size[num] ++;
                }
        }
```

```
}
//BCC
int dfn[N], low[N], dfs_clock, bccno[N], have[N], size[N], bcc, top;
bool iscut[N];
void tarjan(int x, int fa) {
        low[x] = dfn[x] = ++dfs_clock;
        int child = 0;
        for(int i = head[x]; i; i = nxt[i]) {
                if (!dfn[to[i]]) {
                        child ++;
                         tarjan(to[i], x);
                         low[x] = min(low[to[i]], low[x]);
                         if (low[to[i]] >= dfn[x]) iscut[x] = 1;
                } else low[x] = min(low[x], dfn[to[i]]);
        }
        if (fa < 0 && child == 1) iscut[x] = 0;</pre>
}
bool vis[N];
void dfs(int x) {
        vis[x] = 1; size[bcc] ++;
        for(int i = head[x]; i; i = nxt[i])
                if (!vis[to[i]]) {
                         if (!iscut[to[i]]) dfs(to[i]);
                        else if (bccnoo[to[i]] != bcc)
                                 bccno[to[i]] = bcc, have[bcc] ++;
                }
}
2.2
    最大流
const int N = 1010, M = 100010;
const int INF = 1e9;
struct edge{
        int to, v;
}E[M << 1];
int head[N], nxt[M << 1], cnt = 1;</pre>
void add(int x, int y, int z) {
        E[++ cnt] = (edge)\{y, z\}; nxt[cnt] = head[x]; head[x] = cnt;
        E[++ cnt] = (edge)\{x, 0\}; nxt[cnt] = head[y]; head[y] = cnt;
}
```

```
int S, T, d[N], cur[N];
bool mklevel() {
        memset(d, -1, sizeof d);
        Q.push(S);
        d[S] = 0;
        while(!Q.empty()) {
                int x = Q.front(); Q.pop();
                for(int i = head[x]; i; i = nxt[i])
                         if (d[E[i].to] == -1 \&\& E[i].v) {
                                 d[E[i].to] = d[x] + 1;
                                 Q.push(E[i].to);
                         }
        }
        return d[T] != -1;
}
int dfs(int x, int a) {
        if (x == T \mid \mid a == 0) return a;
        int flow = 0;
        for(int &i = cur[x]; i; i = nxt[i])
                if (d[E[i].to] == d[x] + 1 && E[i].v) {
                         int f = dfs(E[i].to, min(E[i].v, a - flow));
                         E[i].v -= f;
                         E[i ^1].v += f;
                         flow += f;
                         if (f == a) break;
                }
        if (!flow) d[x] = -1;
        return flow;
}
int Dinic() {
        int ans = 0;
        while(mklevel()) {
                for(int i = 0; i <= T; i ++) cur[i] = head[i];</pre>
                ans += dfs(S, INF);
        }
        return ans;
}
```

#### 2.3 最小费用流

```
int cost, flow;
struct edge {
        int from, to, v, c;
}E[M];
void ins(int x, int y, int z, int c) {
        E[++ cnt] = (edge)\{x, y, z, c\};
        nxt[cnt] = heda[x]; head[x] = cnt;
}
void add(int x, int y, int z, int c) {
        ins(x, y, z, c); ins(y, x, 0, -c);
}
int S, T, d[N], from[N], Q[M];
bool inq[N];
bool spfa() {
        int 1 = 0, r = -1;
        for(int i = 0; i <= T; i ++)</pre>
                if (d[E[i].to] > d[x] + E[i].c && E[i].v) {
                         d[E[i].to] = d[x] + E[i].c;
                        from[E[i].to] = i;
                        if (!inq[E[i].to]) {
                                 Q[++ r] = E[i].to;
                                 inq[E[i].to] = 1;
                        }
                }
        return d[T] != INF;
}
void mcf() {
        int x = INF;
        for(int i = from[T]; i; i = from[E[i].from])
                x = min(x, E[i].v);
        for(int i = from[T]; i; i = from[E[i].from]) {
                E[i].v -= x;
                E[i ^1].v += x;
        cost += x * d[T];
        flow += x;
}
```

### 3 数据结构

#### 3.1 K-D 树

```
const int N = 100010;
int D, p[N], tot, root;
struct node {
        int d[2], mn[2], mx[2], 1, r, D, size, sum, v;
        int& operator [] (int x) {return d[x];}
}t[N], now;
inline bool cmp(int x, int y) {return t[x][D] < t[y][D];}</pre>
#define L t[o].l
#define R t[o].r
\#define\ mid\ ((l\ +\ r)\ >>\ 1)
inline void Push_up(int o) {
        for(int i = 0; i < 2; i ++) {
                t[o].mn[i] = min(t[o].mn[i], min(t[L].mn[i], t[R].mn[i]));
                t[o].mx[i] = max(t[o].mx[i], max(t[L].mx[i], t[R].mx[i]));
        }
        t[o].sum = t[L].sum + t[R].sum + t[o].v;
        t[o].size = t[L].size + t[R].size + 1;
}
inline int build(int 1, int r, int dir) {
        D = dir;
        nth_element(p + 1, p + mid, p + r + 1, cmp);
        int o = p[mid];
        for(int i = 0; i < 2; i ++) t[o].mn[i] = t[o].mx[i] = t[o][i];
        t[o].sum = t[o].v;
        L = 1 < mid ? build(1, mid - 1, dir ^ 1) : 0;
        R = mid < r ? build(mid + 1, r, dir ^ 1) : 0;
        Push up(o);
        return o;
}
inline void dfs(int o){
        if (!o) return;
        dfs(L);
        p[++cnt] = o;
        dfs(R);
}
inline void rebuild(int &o) {
        cnt = 0;
        dfs(o);
```

```
o = bulid(1, cnt, t[o].D);
}
inline void Insert(int &o, int dir) {
        if (!o) {
                o = ++tot; t[o] = now;
                for(int i = 0; i < 2; i ++) t[o].mn[i] = t[o].mx[i] = t[o][i];
                t[o].D = dir;
                t[o].size = 1;
                t[o].sum = t[o].v;
                return;
        if (now[dir] < t[o][dir]) {</pre>
                Insert(L, dir ^ 1);
                Push_up(o);
                if ((double)t[L].size > (double)t[o].size * 0.7) rebuild(o);
        } else {
                Insert(R, dir ^ 1);
                Push_up(o);
                if ((double)t[R].size > (double)t[o].size * 0.7) rebuild(o);
        }
}
inline double getans(int o, int k) {
        if (!o) return INF;
}
inline double calc(int o, double k) {
        if (!o) return INF;
        double ans = ....;
}
inline void query(int o, double k) {
        if (!o) return 0;
        double dl = calc(L, k), dr = calc(R, k), d0 = getans(o, k);
        ans = max(ans, d0);
        if (dl < dr) {
                if (dr > ans \&\& R) query(R, k);
                if (dl > ans && L) query(L, k);
        } else {
                if (dl > ans && L) query(L, k);
                if (dr > ans && R) query(R, k);
        }
}
```

#### 3.2 可持久化 Trie

```
int t[M][2], rt[N], id[M], tot;
inline void Ins(int pre, int x, int k) {
        int now = rt[k] = ++tot; id[tot] = k;
        for(int i = 30; i >= 0; i --) {
                int j = (x >> i) & 1;
                t[now][j ^ 1] = t[pre][j ^ 1];
                t[now][j] = ++tot; id[tot] = k;
                now = tot;
                pre = t[pre][j];
        }
}
inline int query(int 1, int r, int x) {
        int ans = 0, now = rt[r];
        for(int i = 30; i >= 0; i --) {
                int j = ((x >> i) & 1) ^ 1;
                if (id[t[now][j]] >= 1) ans |= 1 << i; else j = 1;
                now = t[now][j];
        }
        return ans;
}
3.3 Link-Cut-Tree
struct LCT {
        int c[N][2], fa[N], v[N], mx[N];
        bool rev[N];
        int st[N], top;
#define L c[x][0]
#define R c[x][1]
        void Push_up(int x) {
                mx[x] = x;
                if (v[mx[L]] > v[mx[x]]) mx[x] = mx[L];
                if (v[mx[R]] > v[mx[x]]) mx[x] = mx[R];
        }
        void Push_down(int x) {
                if (rev[x]) rev[x] = 0, rev[L] ^= 1, rev[R] ^= 1, swap(L, R);
        }
        bool not_root(int x) {
                return c[fa[x]][0] == x || c[fa[x]][1] == x;
        }
        void Rotate(int x) {
```

```
if (not\_root(y)) c[z][c[z][1] == y] = x;
                fa[x] = z; fa[y] = x; fa[c[x][r]] = y;
                c[y][1] = c[x][r]; c[x][r] = y;
                Push_up(y);
        }
        void preview(int x) {
                top = 0; st[++top] = x;
                for(;not_root(x); x = fa[x]) st[++top] = fa[x];
                for(int i = top; i; i --) Push_down(st[i]);
        void splay(int x, int y = 0) {
                for(preview(x); not_root(x); Rotate(x))
                        if (not_root(y = fa[x]))
                                Rotate(c[y][1] == x \hat{c}[fa[y][1]] == y ? x : y);
                Push_up(x);
        }
        void access(int x, int y = 0) {
                for(;x;splay(x), c[x][1] = y, y = x, x = fa[x]);
        }
        void makeroot(int x) {
                access(x); splay(x); rev[x] ^= 1;
        void link(int x, int y) {
                makeroot(x); fa[x] = y;
        void cut(int x, int y) {
                makeroot(x); access(y); splay(y);
                if (c[y][0] == x) c[y][0] = fa[x] = 0;
        int query(int x, int y) {
                makeroot(x); access(y); splay(y);
                return mx[y];
        }
};
3.4
    可持久化线段树
struct LCT {
        int c[N][2], fa[N], v[N], mx[N];
        bool rev[N];
        int st[N], top;
#define L c[x][0]
```

int  $y = fa[x], z = fa[y], l = (c[y][1] == x), r = l^1;$ 

```
#define R c[x][1]
        void Push_up(int x) {
                mx[x] = x;
                if (v[mx[L]] > v[mx[x]]) mx[x] = mx[L];
                if (v[mx[R]] > v[mx[x]]) mx[x] = mx[R];
        }
        void Push_down(int x) {
                if (rev[x]) rev[x] = 0, rev[L] ^= 1, rev[R] ^= 1, swap(L, R);
        }
        bool not_root(int x) {
                return c[fa[x]][0] == x || c[fa[x]][1] == x;
        void Rotate(int x) {
                int y = fa[x], z = fa[y], l = (c[y][1] == x), r = l^1;
                if (not\_root(y)) c[z][c[z][1] == y] = x;
                fa[x] = z; fa[y] = x; fa[c[x][r]] = y;
                c[y][1] = c[x][r]; c[x][r] = y;
                Push_up(y);
        void preview(int x) {
                top = 0; st[++top] = x;
                for(;not_root(x); x = fa[x]) st[++top] = fa[x];
                for(int i = top; i; i --) Push_down(st[i]);
        }
        void splay(int x, int y = 0) {
                for(preview(x); not_root(x); Rotate(x))
                        if (not_root(y = fa[x]))
                                Rotate(c[y][1] == x \hat{c}[fa[y][1]] == y ? x : y);
                Push_up(x);
        void access(int x, int y = 0) {
                for(;x;splay(x), c[x][1] = y, y = x, x = fa[x]);
        void makeroot(int x) {
                access(x); splay(x); rev[x] ^= 1;
        }
        void link(int x, int y) {
                makeroot(x); fa[x] = y;
        void cut(int x, int y) {
                makeroot(x); access(y); splay(y);
                if (c[y][0] == x) c[y][0] = fa[x] = 0;
```

## 4 字符串

#### 4.1 AC 自动机

```
int cnt = 1;
struct Trie {
        int ch[26], cnt, fail;
        bool sign;
}T[N];
inline int id(char c) {return c - 'A';}
void Ins(char *s) {
        int x = 1, y, l = strlen(s);
        for(int i = 0; i < 1; i ++) {</pre>
                y = id(s[i]);
                if (!T[x].ch[y]) T[x].ch[y] = ++cnt;
                x = T[x].ch[y];
        }
        T[x].sign = 1;
}
int Q[N];
void make_fail() {
        int 1 = 0, r = -1;
        Q[++r] = 1;
        while(1 <= r) {
                int x = Q[1++], y, j;
                for(int i = 0; i < 26; i ++) {</pre>
                         j = T[x].fail;
                         T[x].sign |= T[j].sign;
                         while(j && !T[x].ch[i]) j = T[j].fail;
                         if (T[x].ch[i]) {
                                 y = T[x].ch[i];
                                 T[y].fail = j ? T[j].ch[i] : 1;
                                 Q[++r] = y;
                         } else T[x].ch[i] = j ? T[j].ch[i] : 1;
                }
        }
```

```
}
```

#### 4.2 后缀数组

```
int n, m, sa[N], c[N], wa[N], wb[N], rank[N], height[N];
inline bool cmp(int *r, int a, int b, int l) {
        return r[a] == r[b] \&\& r[a + 1] == r[b + 1];
}
void DA(char *s, int *sa, int n, int m) {
        int *x = wa, *y = wb;
        for(int i = 0; i < m; i ++) c[i] = 0;
        for(int i = 0; i < n; i ++) c[x[i] = s[i]] ++;
        for(int i = 1; i < m; i ++) c[i] += c[i - 1];
        for(int i = n - 1; i \ge 0; i --) sa[-- c[x[i]]] = i;
        for(int j = 1; p = 0; p < n; j <<= 1, m = p) {
                for(int i = n - j; i < n; i ++) y[p ++] = i;
                for(int i = 0; i < n; i ++) if (sa[i] >= j) y[p ++] = sa[i] - j;
                for(int i = 0; i < m; i ++) c[i] = 0;
                for(int i = 0; i < n; i ++) c[x[y[i]]] ++;</pre>
                for(int i = 1; i < m; i ++) c[i] += c[i - 1];
                for(int i = n - 1; i \ge 0; i --) sa[-- c[x[y[i]]]] = y[i];
                swap(x, y); p = 1; x[sa[0]] = 0;
                for(int i = 1; i < n; i ++) x[sa[i]] = cmp(y, sa[i - 1], sa[i], j) ? p - 1 : p
        }
}
void calcheight(char *s, int *sa, int n) {
        int k = 0;
        for(int i = 1; i <= n; i ++) rank[sa[i]] = i;
        for(int i = 0; i < n; i ++) {
                if (k) k --;
                int j = sa[rank[i] - 1];
                while(s[i + k] == s[j + k]) k ++;
                height[rank[i]] = k;
        }
}
int main() {
        n = strlen(s); m = size_of_character_set;
        DA(s, sa, n + 1, 39);
        calcheight(s, sa, n);
```

```
}
```

```
4.3 回文自动机
```

```
int ch[N][26], fail[N], len[N], tot, cnt[N];
void ready() {
        len[0] = 0; len[1] = -1;
        fail[0] = 1; fail[1] = -1;
}
void Insert(char *s, int *cnt) {
        int now = 1, l = strlen(s), x, y, tmp;
        for(int i = 0; i < 1; ++ i) {
                x = s[i] - 'a';
                while(s[i] != s[i - len[now] - 1]) now = fail[now];
                if (!ch[now][x]) {
                        ch[now][x] = ++ tot;
                        len[tot] = len[now] + 2;
                y = ch[now][x];
                tmp = fail[now];
                if (tmp == -1) fail[y] = 0;
                else {
                        while(s[i] != s[i - len[tmp] - 1]) tmp = fail[tmp];
                        fail[y] = ch[tmp][x];
                }
                now = y;
                cnt[now] ++;
        }
}
```

### 4.4 Manacher 算法

```
//求串 S 中最长的回文子串的长度
char s[N];
int a[N];
int manacher(char *s) {
    memset(p, 0, sizeof p);
    int n = strlen(s);
    for(int i = 1; i <= n; i ++) a[i << 1] = b[i - 1];
    n = n << 1 | 1;
    int id = 0, mx = 0, ans = 0;
    for(int i = 1; i <= n; i ++) {
        if (mx > i) p[i] = min(p[2 * id - i], mx - i);
        while(i - p[i] - 1 > 0 && i + p[i] + 1 <= n
```