## Complexity Analysis: Fundamental Laws of Network Dynamics

**Theorem 1.** The optimization problem of finding the optimal state  $S^*$ , as defined in Law IV of the Fundamental Laws of Network Dynamics, belongs to complexity class  $\mathbf{P}$ .

*Proof.* Consider the objective function from Law IV:

$$F(S_t) = \alpha(S_t) O(S_t) + \beta(S_t) L(S_t),$$

where:

- $\alpha(S_t) = \frac{|\{(i,j) : R(i,j) \neq D\}|}{|\Omega|(|\Omega|-1)/2}$  is the fraction of non-disjunctive relations;
- $\beta(S_t) = 1 \alpha(S_t)$  is the fraction of disjunctive relations;
- $O(S_t) = \sum_{i < j} \Lambda(i, j)$  is the realized value from cohesive relations;
- $L(S_t) = \sum_{\substack{i < j \\ R(i,j) = D}} (V_i I_i + V_j I_j)$  is the latent value from disjunctive relations.

Since  $V_i \ge 0$  and  $I_i \ge 0$  for all nodes  $X_i$  (as they represent capacities/propensities), the weights

$$w_{ij} = V_i I_i + V_j I_j$$

are non-negative.

## Extreme cases:

1. All relations cohesive (C):  $\alpha = 1$ ,  $\beta = 0$ , so

$$F(S) = O(S) = \sum_{i < j} w_{ij}.$$

2. All relations disjunctive (D):  $\alpha = 0$ ,  $\beta = 1$ , so

$$F(S) = L(S) = \sum_{i < j} w_{ij}.$$

3. Mixed assignments: For a mix of C and D (ignoring N for now),

$$F(S) = \alpha \cdot \left(\sum_{C} w_{ij}\right) + (1 - \alpha) \cdot \left(\sum_{D} w_{ij}\right).$$

Since the total  $\sum w_{ij}$  is fixed and weights are non-negative, this weighted average cannot exceed the total sum.

Neutral relations (N) do not increase  $F(S_t)$ , since they contribute to neither  $O(S_t)$  nor  $L(S_t)$ , while reducing  $\alpha$  and  $\beta$  away from their extremes.

Thus, the maximum is:

$$\max F(S_t) = \sum_{i < j} w_{ij}.$$

## Polynomial-time algorithm:

1. Compute the total weight

$$W = \sum_{i < j} w_{ij},$$

which requires  $O(n^2)$  operations for n nodes.

2. Return either the all-C or all-D assignment (constant time decision).

Additionally, simulating the dynamics for a polynomial number of steps (per Laws V and VII) requires updating  $O(n^2)$  relations per step, which is also polynomial in n.

Hence, the problem is solvable in polynomial time, i.e. in class  $\mathbf{P}$ .