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Chapter 10.2
Symmetric Encryption and Message Confidentiality

Book Reading: Computer Security Principles and Practice (3ed), 2015, p. 637-668

Symmetric Encryption and Message Confidentiality

- also known as: conventional encryption, secret-key, or single-key encryption
  - only alternative before public-key crypto in 70's
  - still most widely used alternative
  - has ingredients: plaintext, encryption algorithm, secret key, ciphertext, and decryption algorithm
- generically classified along dimensions of:
  1. type of operations used
  2. number of keys used
  3. way in which the plaintext is processed

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Cryptanalysis

- attacks:
  - ciphertext only - least info, hardest
  - known plaintext - some plain/cipher pairs
  - chosen plaintext - get own plain/cipher pairs
  - chosen ciphertext - rarer
  - chosen text - rarer
- only weak algs fail a ciphertext-only attack
- usually design algs to withstand a known-plaintext attack

Computationally Secure Algs

- encryption is computationally secure if:
  - cost of breaking cipher exceeds info value
  - time required to break cipher exceeds the useful lifetime of the info
- usually very difficult to estimate the amount of effort required to break
- can estimate time/cost of a brute-force attack (see Ch 2)

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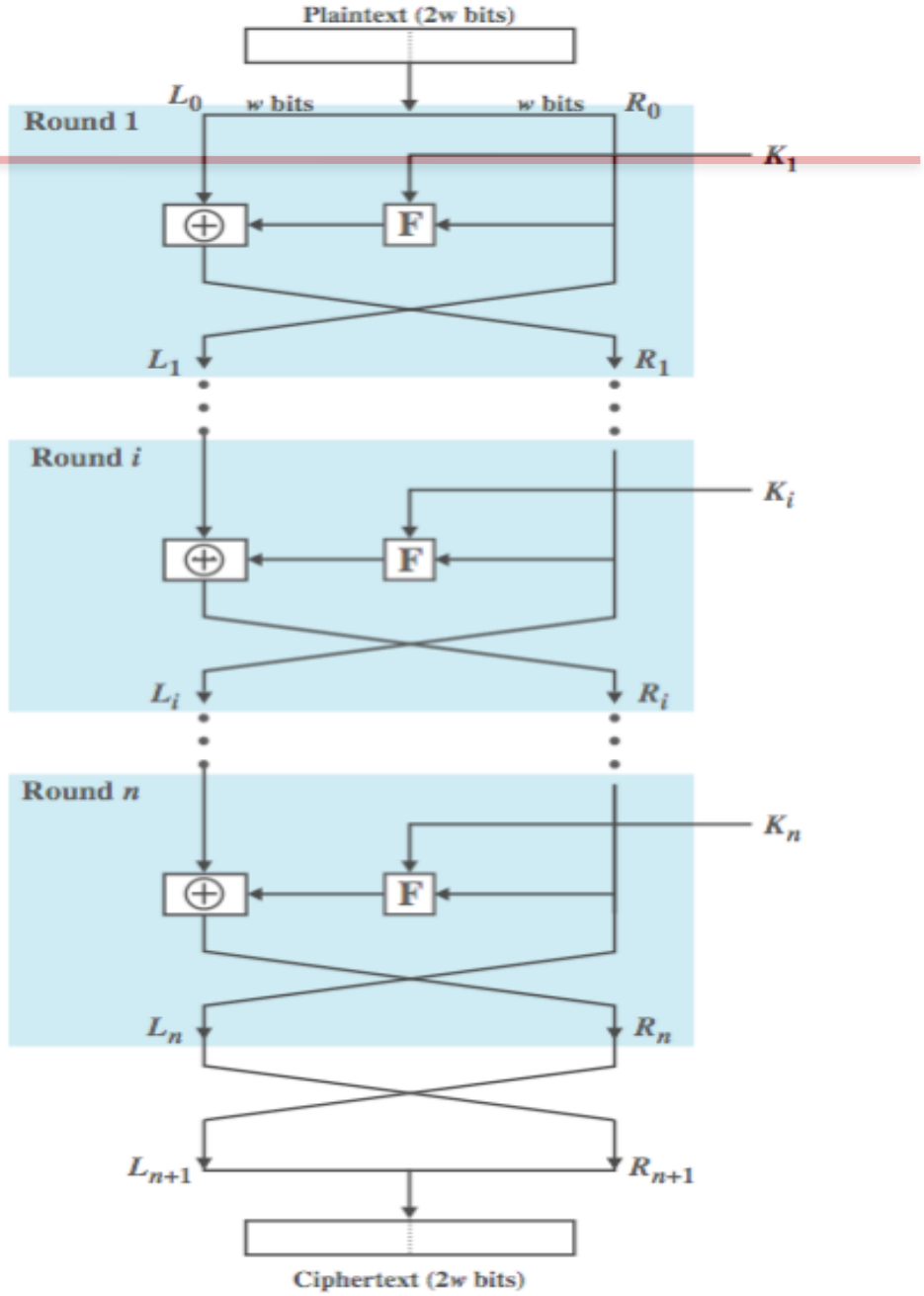
Feistel Cipher Structure



Block Cipher Structure

- have a general iterative block cipher structure
  - with a sequence of rounds
  - with substitutions / permutations controlled by key
- parameters and design features:
  - block size
  - key size
  - number of rounds
  - subkey generation algorithm
  - round function
  - also: fast software en/decrypt, ease of analysis

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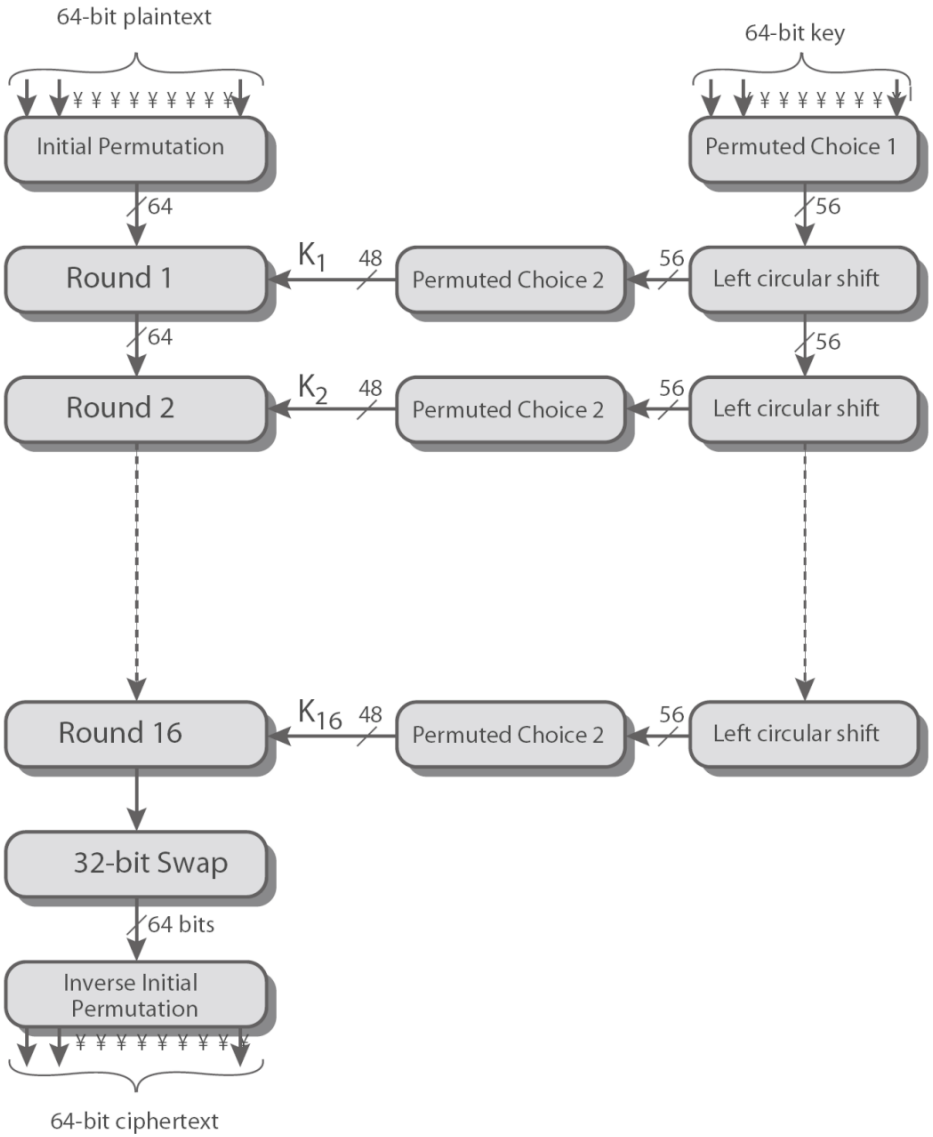
Data Encryption Standard (DES)



Triple DES (3DES)

- first used in financial applications
- in DES FIPS PUB 46-3 standard of 1999
- uses three keys & three DES executions:

$$C = \mathrm{E}(K_3, \mathrm{D}(K_2, \mathrm{E}(K_1, P)))$$

- decryption same with keys reversed
- use of decryption in second stage gives compatibility with original DES users
- effective 168-bit key length, slow, secure
- AES will eventually replace 3DES

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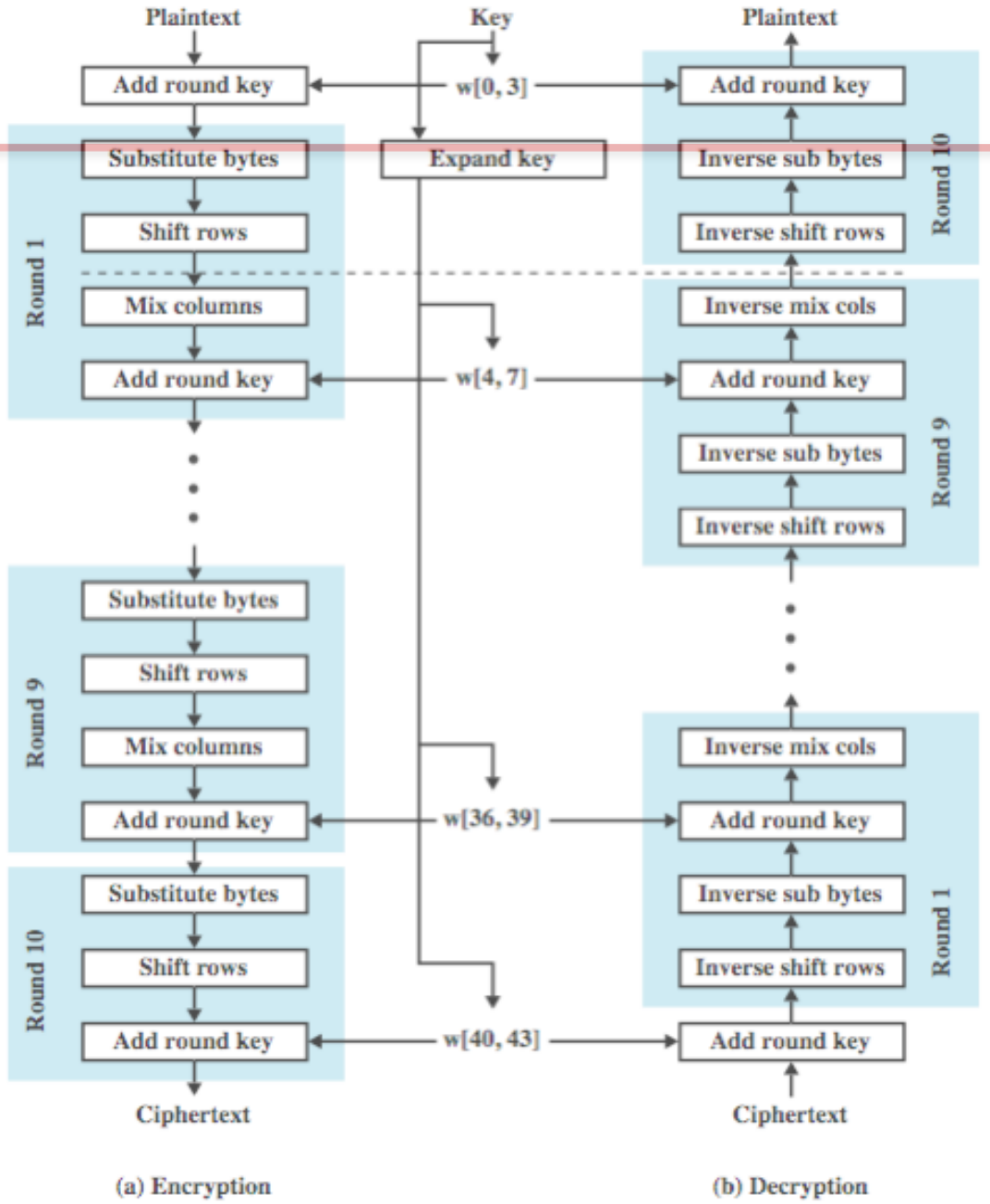
Advanced Encryption Standard (AES)



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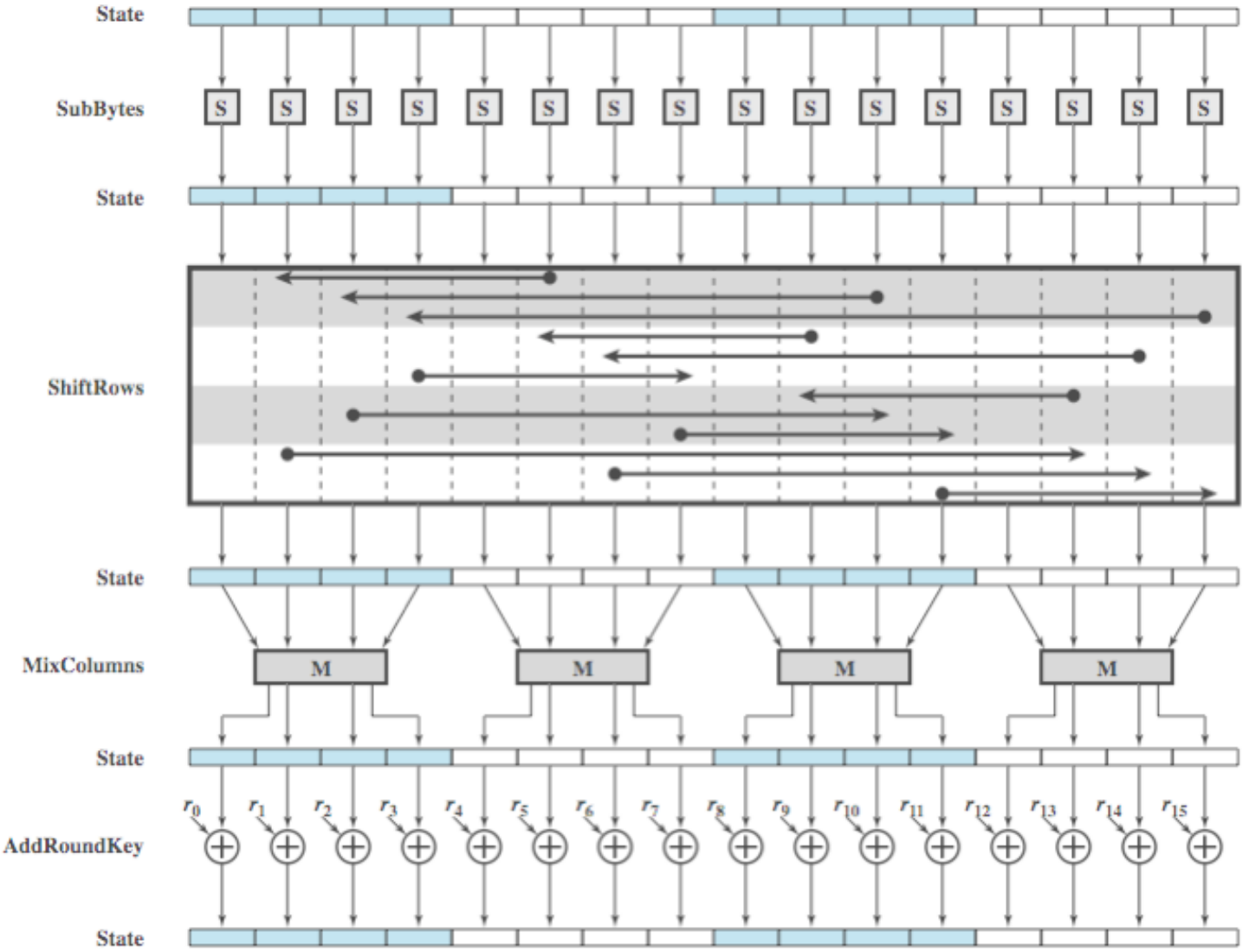
AES Round Structure



Substitute Bytes

- a simple table lookup in S-box
  - a 16×16 matrix of byte values
  - mapping old byte to a new value
    - e.g. {95} maps to {2A}
  - a permutation of all possible 256 8-bit values
- constructed using finite field properties
  - designed to be resistant to known cryptanalytic attacks
- decrypt uses inverse of S-box

Shift Rows

- on encrypt left rotate each row of State by 0,1,2,3 bytes respectively
- decrypt does reverse
- to move individual bytes from one column to another and spread bytes over columns

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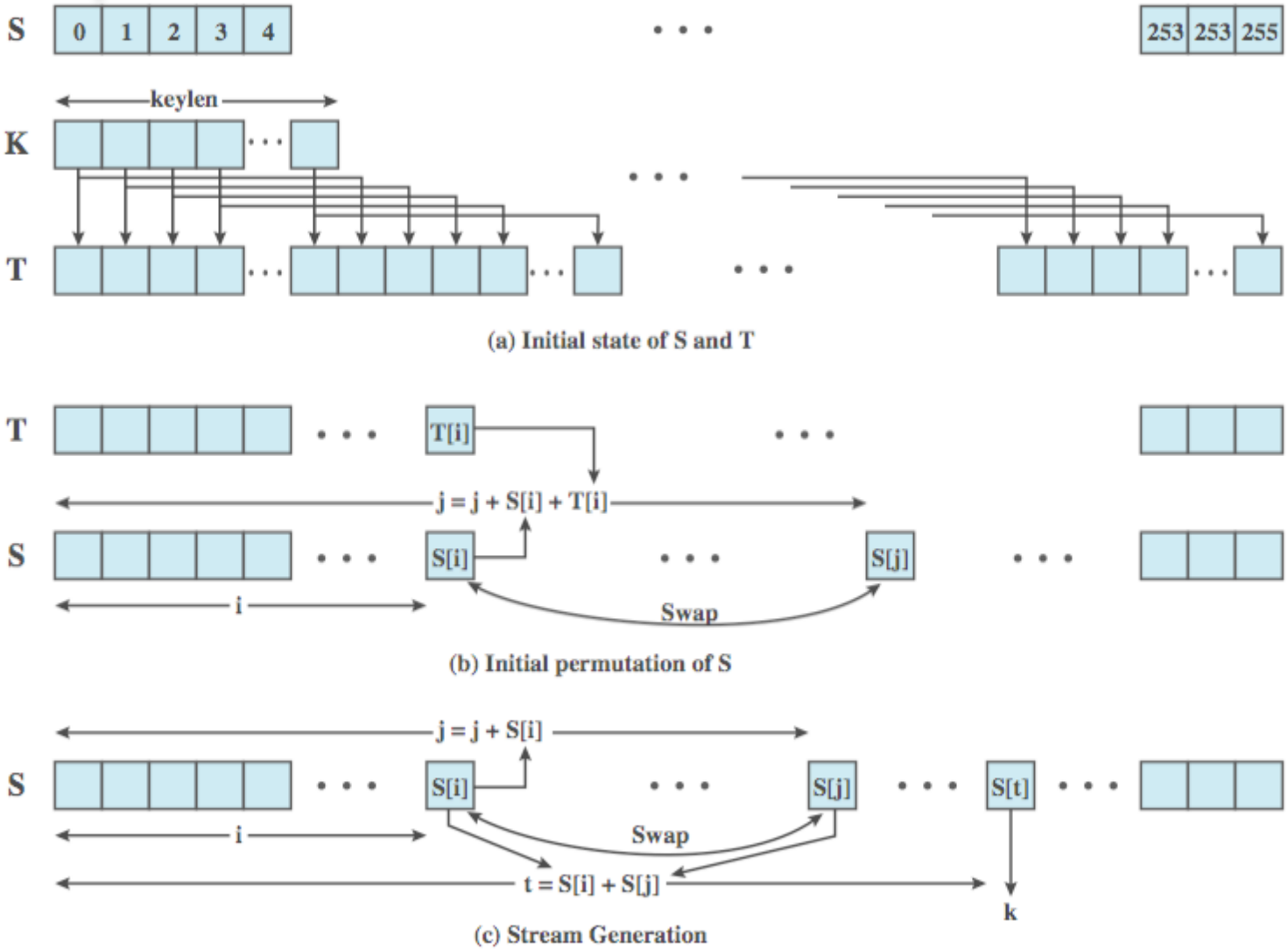
Mix Columns & Add Key

- Mix Columns
  - operates on each column individually
  - mapping each byte to a new value that is a function of all four bytes in the column
  - use of equations over finite fields
  - to provide good mixing of bytes in column
- Add Round Key
  - simply XOR State with bits of expanded key
  - security from complexity of round key expansion and other stages of AES

Stream Ciphers

- processes input elements continuously
- key input to a pseudorandom bit generator
  - produces stream of random like numbers
  - unpredictable without knowing input key
  - XOR keystream output with plaintext bytes
- are faster and use far less code
- design considerations:
  - encryption sequence should have a large period
  - keystream approximates random number properties
  - uses a sufficiently long key

RC4


S
0
1
2
3
4
253
253
255
keylen
K
T
(a) Initial state of S and T
T
T[i]
j = j + S[i] + T[i]
S
S[i]
S[j]
i
Swap
(b) Initial permutation of S
j = j + S[i]
S
S[i]
S[j]
S[t]
i
Swap
t = S[i] + S[j]
k
(c) Stream Generation

Modes of Operation

- block ciphers process data in blocks
  - e.g. 64-bits (DES, 3DES) or 128-bits (AES)
- for longer messages must break up
  - and possibly pad end to blocksize multiple
- have 5 five **modes of operation** for this
  - defined in NIST SP 800-38A
  - modes are: ECB, CBC, CFB, OFB, CTR

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Electronic Codebook (ECB)

- simplest mode
- split plaintext into blocks
- encrypt each block using the same key
- "codebook" because have unique ciphertext value for each plaintext block
  - not secure for long messages since repeated plaintext is seen in repeated ciphertext

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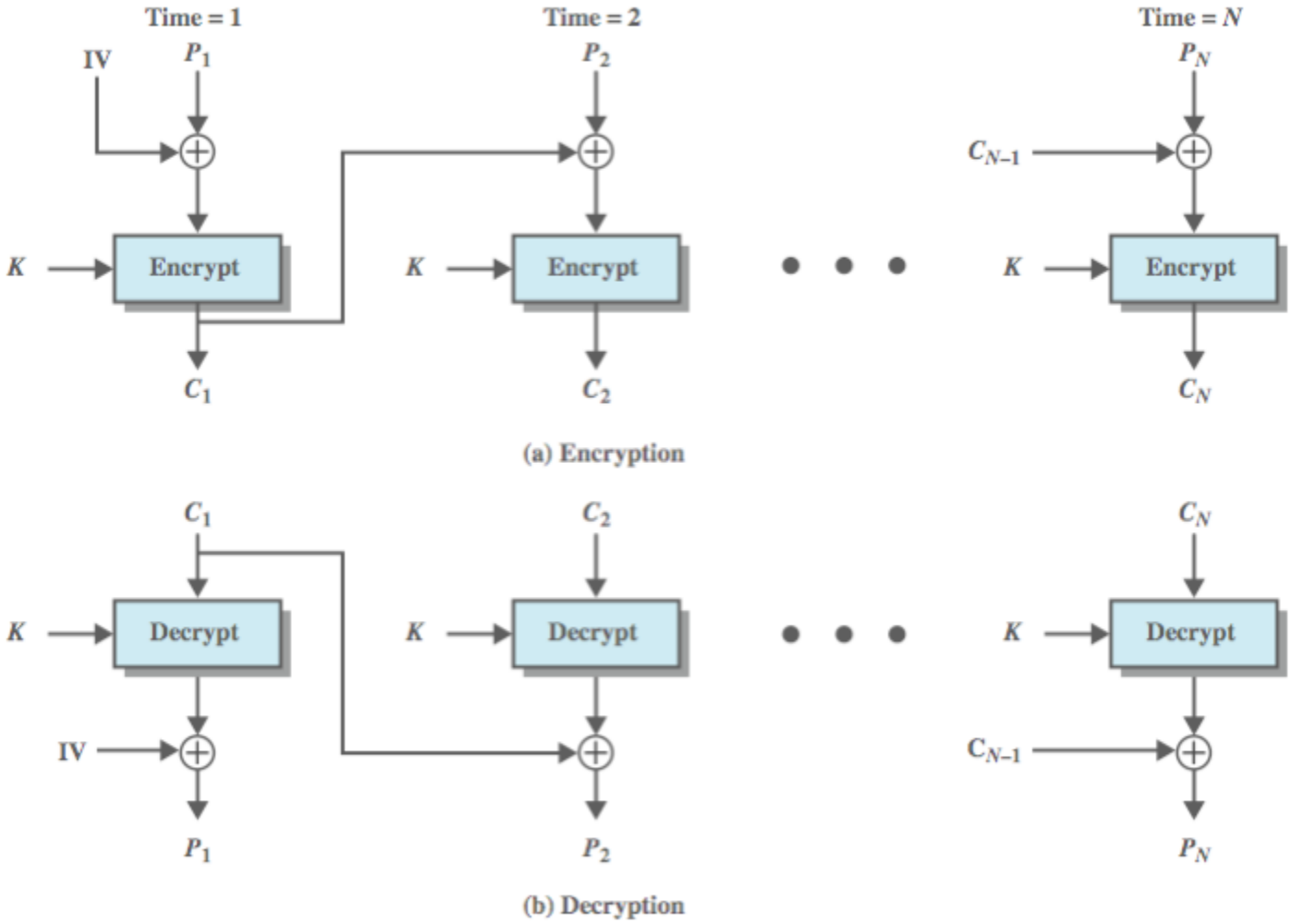
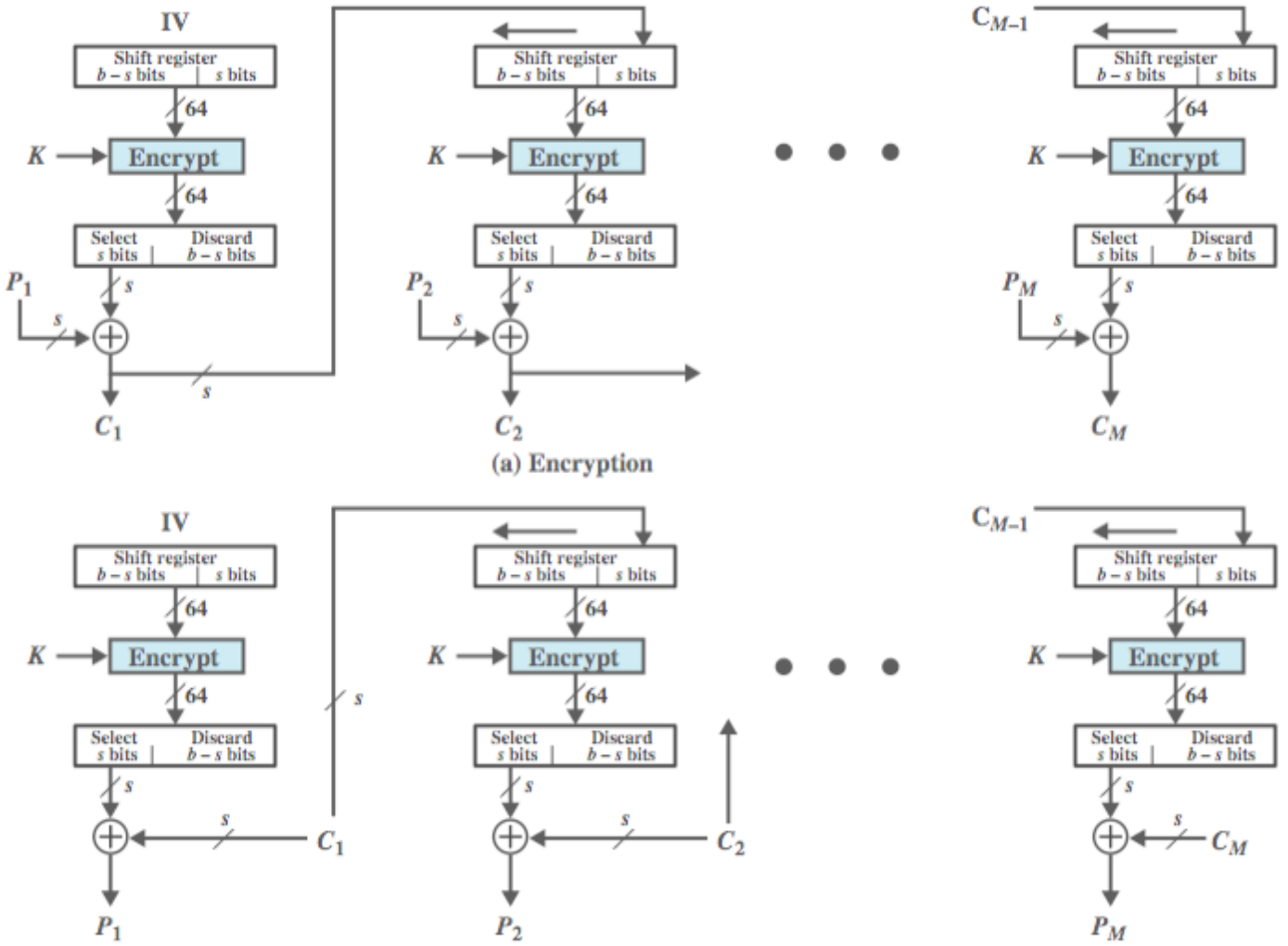
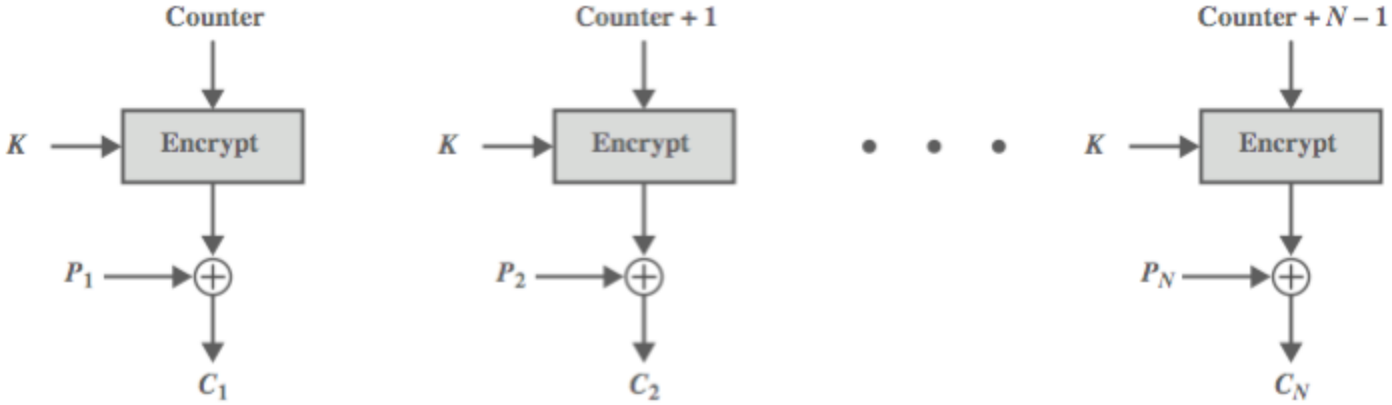
Cipher Block Chaining (CBC)



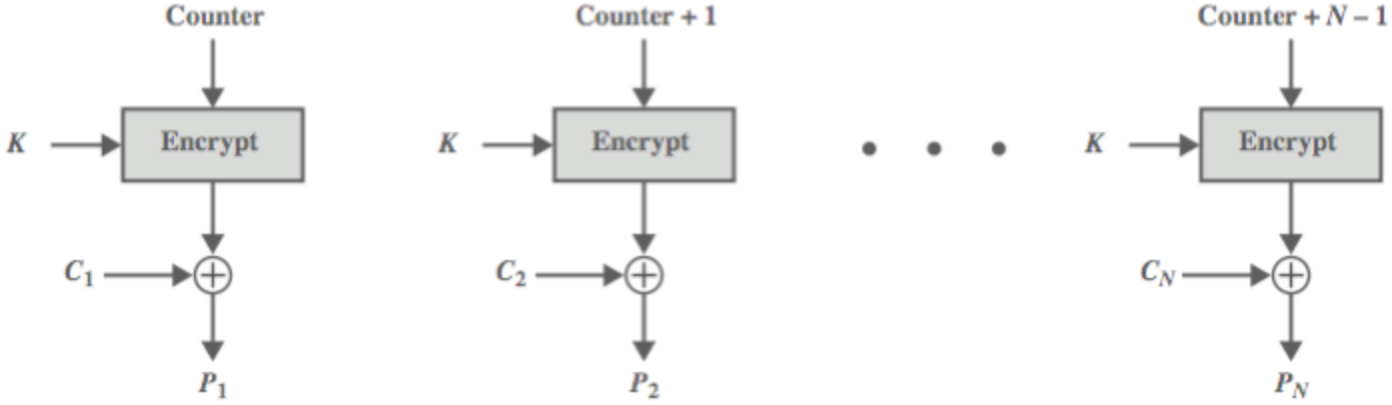
Cipher Feedback (CFB)



Counter (CTR)

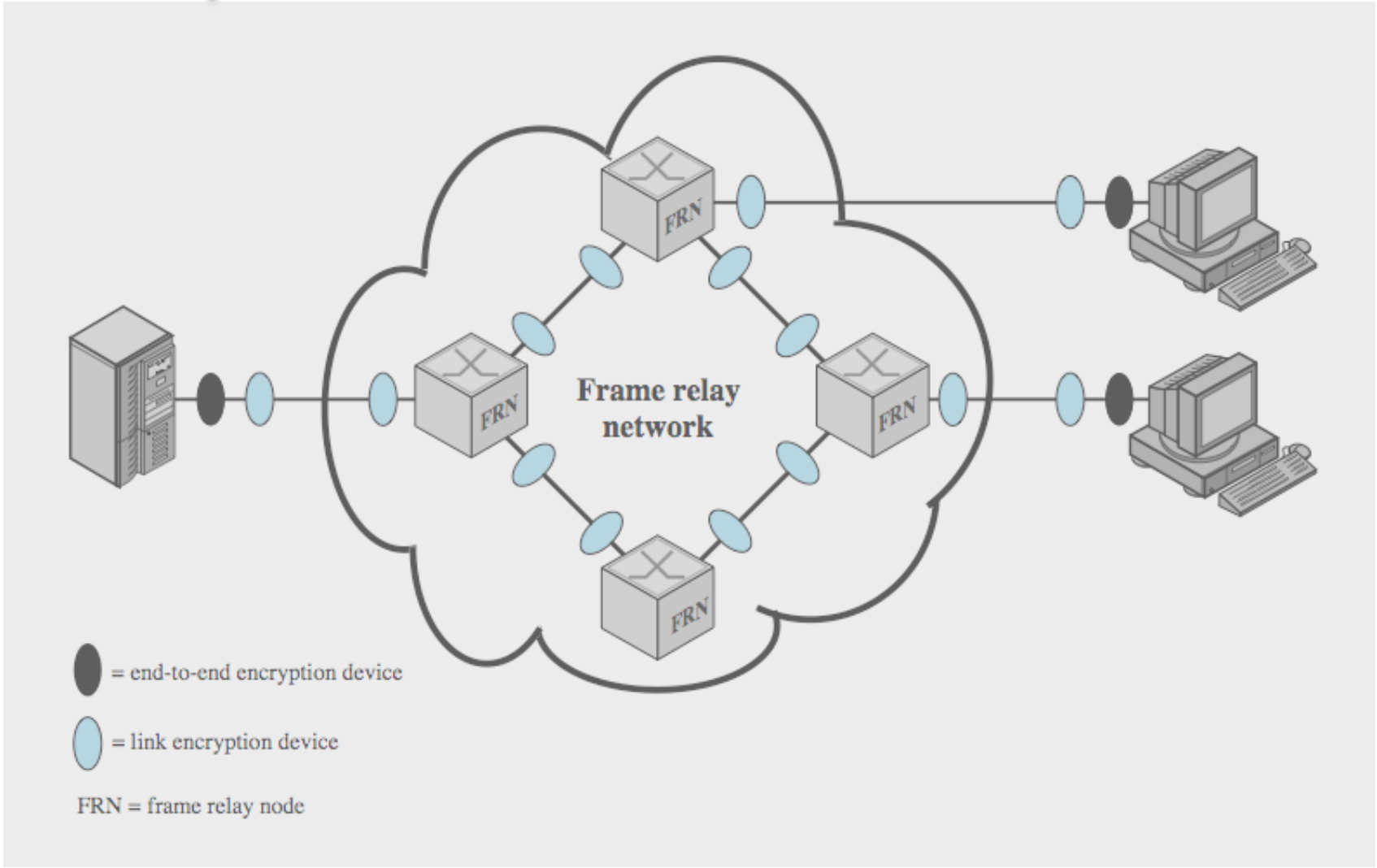


(a) Encryption



(b) Decryption

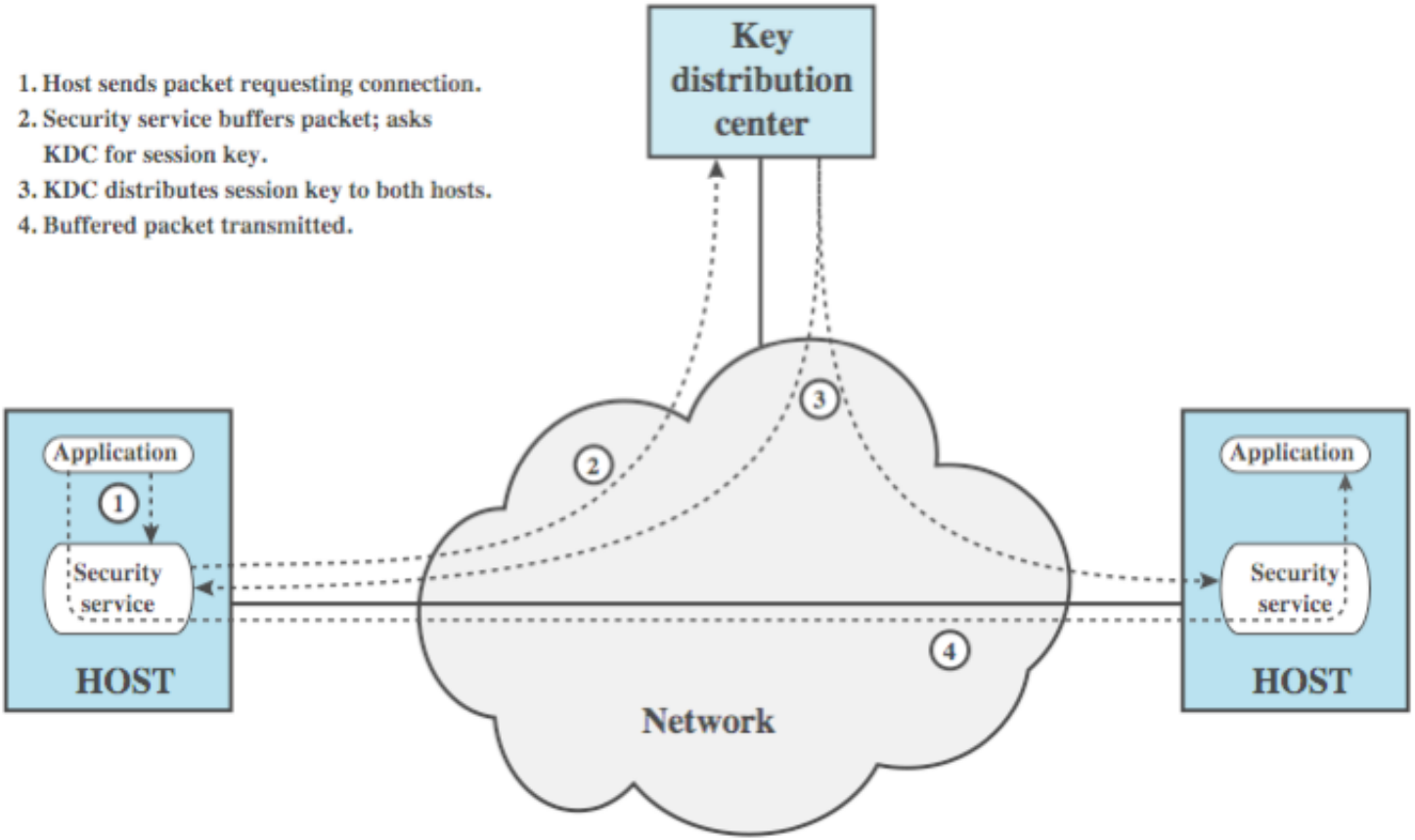
Location of Encryption



Key Distribution

- symmetric crypto needs a shared key:
- two parties A & B can achieve this by:
  - A selects key, physically delivers to B
  - 3rd party select keys, physically delivers to A, B
    - reasonable for link crypto, bad for large no's users
  - A selects new key, sends encrypted using previous old key to B
    - good for either, but security fails if any key discovered
  - 3rd party C selects key, sends encrypted to each of A & B using existing key with each
    - best for end-to-end encryption

Key Distribution



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Summary

- introduced symmetric encryption basics
- DES, 3DES and AES
- stream ciphers and RC4
- modes of operation
- location of encryption
- key distribution