

Domain Name System

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CS 3103: Compute Networks and Protocols

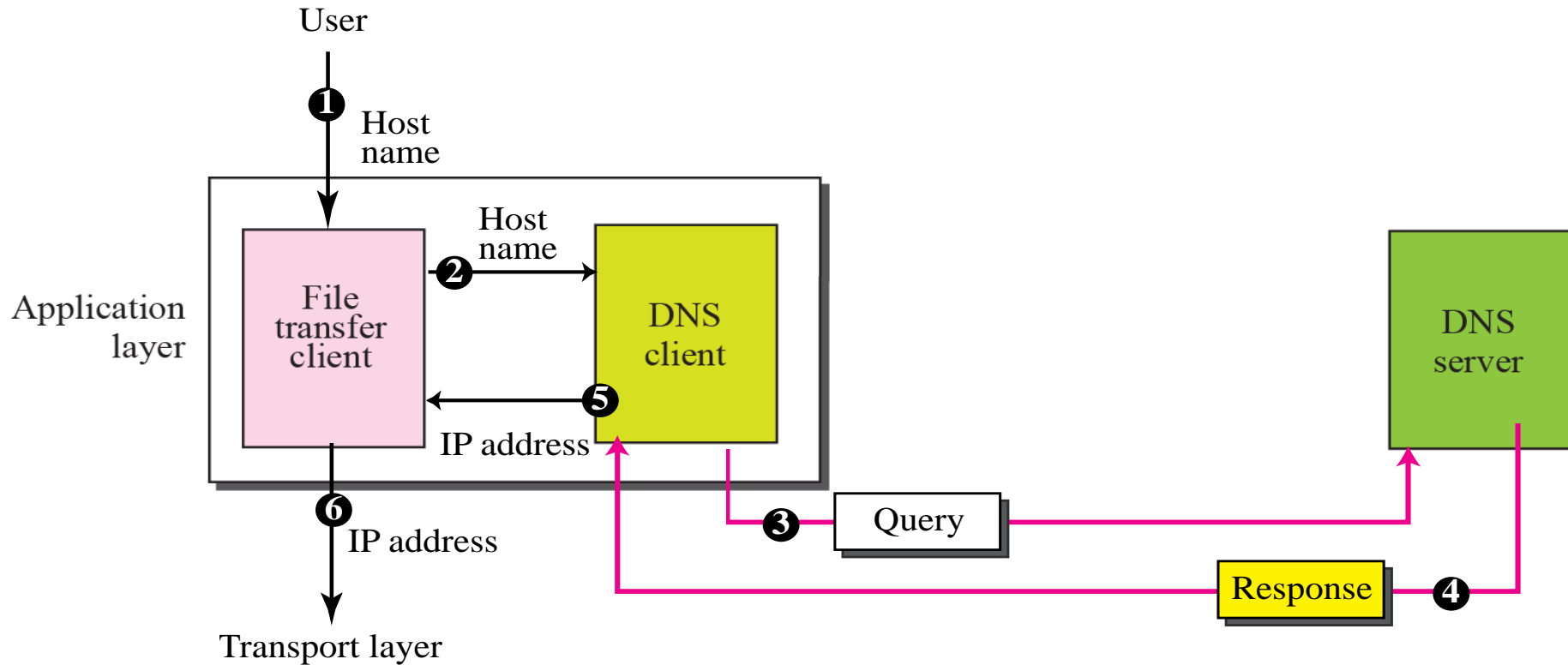
Names Vs. Addresses

- ❑ Names are easier for human to remember
 - ❖ www.comp.nus.edu.sg vs. 137.132.90.2
- ❑ Addresses can be changed without changing names
 - ❖ move www.comp.nus.edu.sg to 203.126.100.199
- ❑ Name could map to multiple addresses
 - ❖ google.com maps to multiple replicas of the Web site and to different “nearby” addresses in different geographies to reduce latency
- ❑ Multiple names could map to the same address

Domain Name System (DNS)

- ❑ The Domain Name System is
 - ❖ a distributed database implemented in a hierarchy of **DNS servers**, and
 - ❖ an application-layer protocol that allows hosts to query the distributed database.
- ❑ The DNS servers are often UNIX machines
 - ❖ running the Berkeley Internet Name Domain (BIND) software.
- ❑ The DNS protocol runs over UDP on port 53.
- ❑ Used by other application-layer protocols
 - ❖ E.g., HTTP, SMTP, and FTP

A Typical Scenario



DNS: Other Use Cases

□ Host aliasing

- ❖ Canonical hostname → alias hostname
- ❖ E.g., webserver1.abc.com → www.abc.com

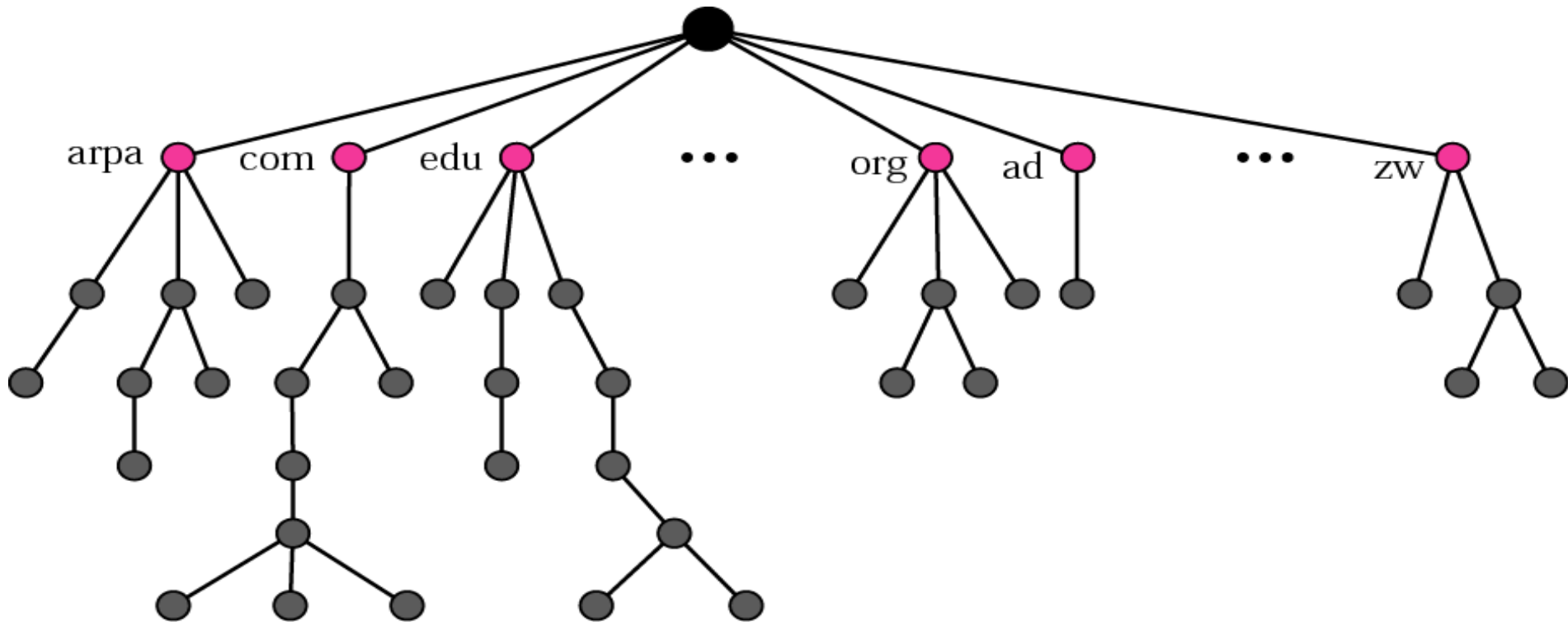
□ Mail server aliasing

- ❖ Bob's account with Hotmail → bob@hotmail.com
- ❖ Web and mail servers share the same

□ Load distribution

- ❖ Load balancing among replicated servers
- ❖ Rotate among the set of IP addresses in replies

Hierarchical Name Space



- ❑ Maximum 128 levels (level 0 to level 127)
- ❑ The label for each level is a string with maximum of 63 characters
- ❑ The root label is a null string

Hierarchical Name Space

❑ Name allocation decentralized to domains

host.sub-subdomain.. . .subdomain.domain[.ROOT]

host: machine name, can be an alias

sub-subdomain: department (comp, eng, ceg, math)

subdomain: institution, company, geography, provider

domain: most significant segment (edu, com, org, net, gov)

❑ Domain: a sub-tree of the domain name space

❑ Domain name

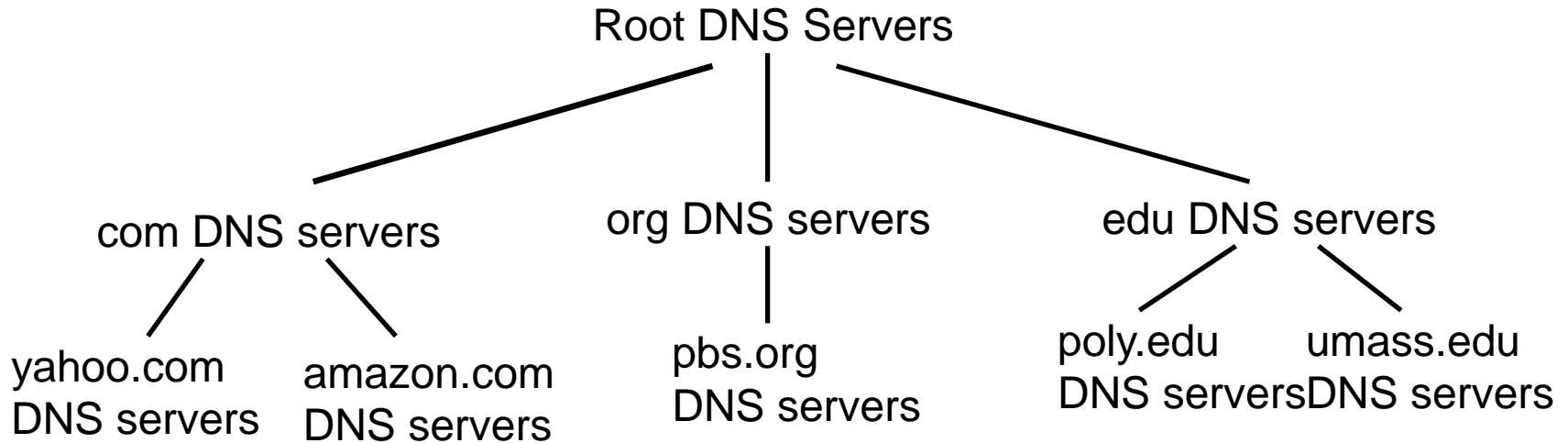
- ❖ Fully Qualified Domain Name (FQDN)

- ❖ Partially Qualified Domain Name (PQDN)

Implementation

- ❑ *Distributed database* implemented in hierarchy of many *name servers*
 - ❖ A distributed database storing resource records (RR)
 - ❖ Client-server query-reply
- ❑ Host, routers and name servers communicate to *resolve* domain names
 - ❖ core Internet function, implemented as application-layer protocol
 - ❖ complexity at network's "edge"

Distributed, Hierarchical Database



client wants IP for www.amazon.com; 1st approx:

- ❑ client queries a root server to find com DNS server
- ❑ client queries com DNS server to get amazon.com DNS server
- ❑ client queries amazon.com DNS server to get IP address for www.amazon.com

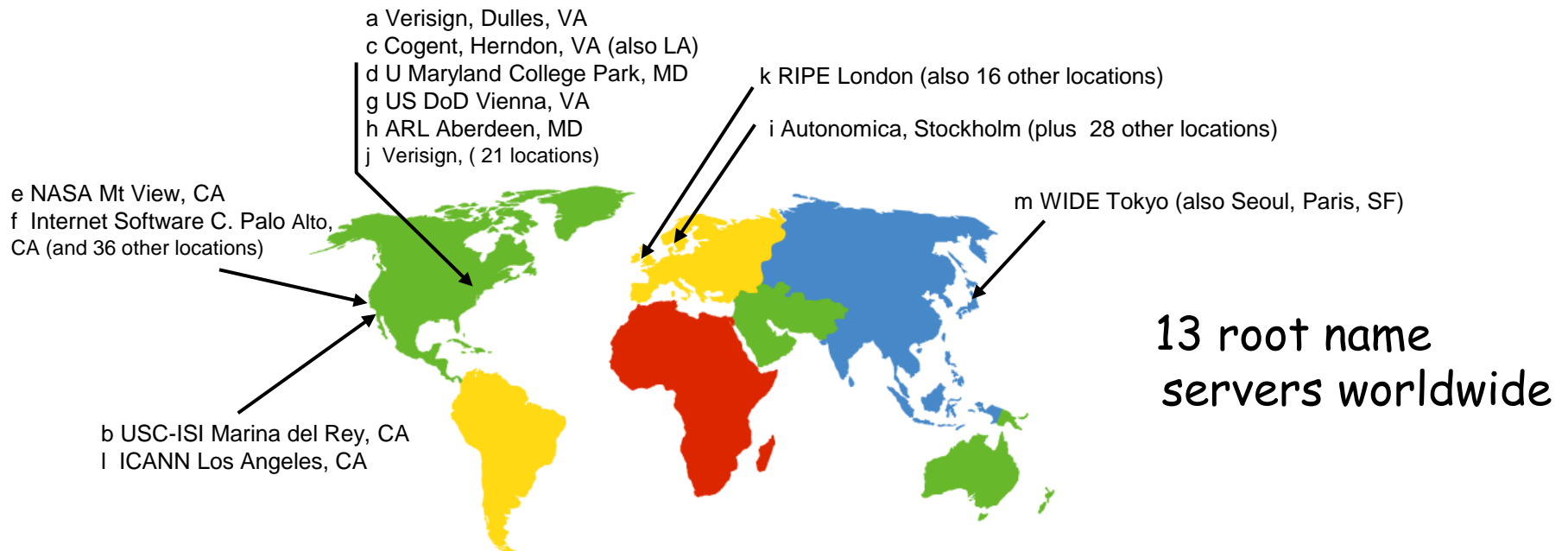
Discussion: Design Choices

- ❑ Why not flat name space?
- ❑ Why not centralize DNS?
 - Single point of failure
 - Traffic volume
 - Distant centralized database
 - Maintenance

Scalability!!

DNS: Root name servers

- ❑ contacted by local name server that can not resolve name
- ❑ root name server:
 - ❖ contacts authoritative name server if name mapping not known
 - ❖ gets mapping
 - ❖ returns mapping to local name server



TLD and Authoritative Servers

Top-level domain (TLD) servers:

- ❖ responsible for com, org, net, edu, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, ca, jp
- ❖ Network Solutions maintains servers for com TLD
- ❖ Educause for edu TLD

Authoritative DNS servers:

- ❖ organization's DNS servers, providing authoritative hostname to IP mappings for organization's servers (e.g., Web, mail).
- ❖ can be maintained by organization or service provider

Local Name Server

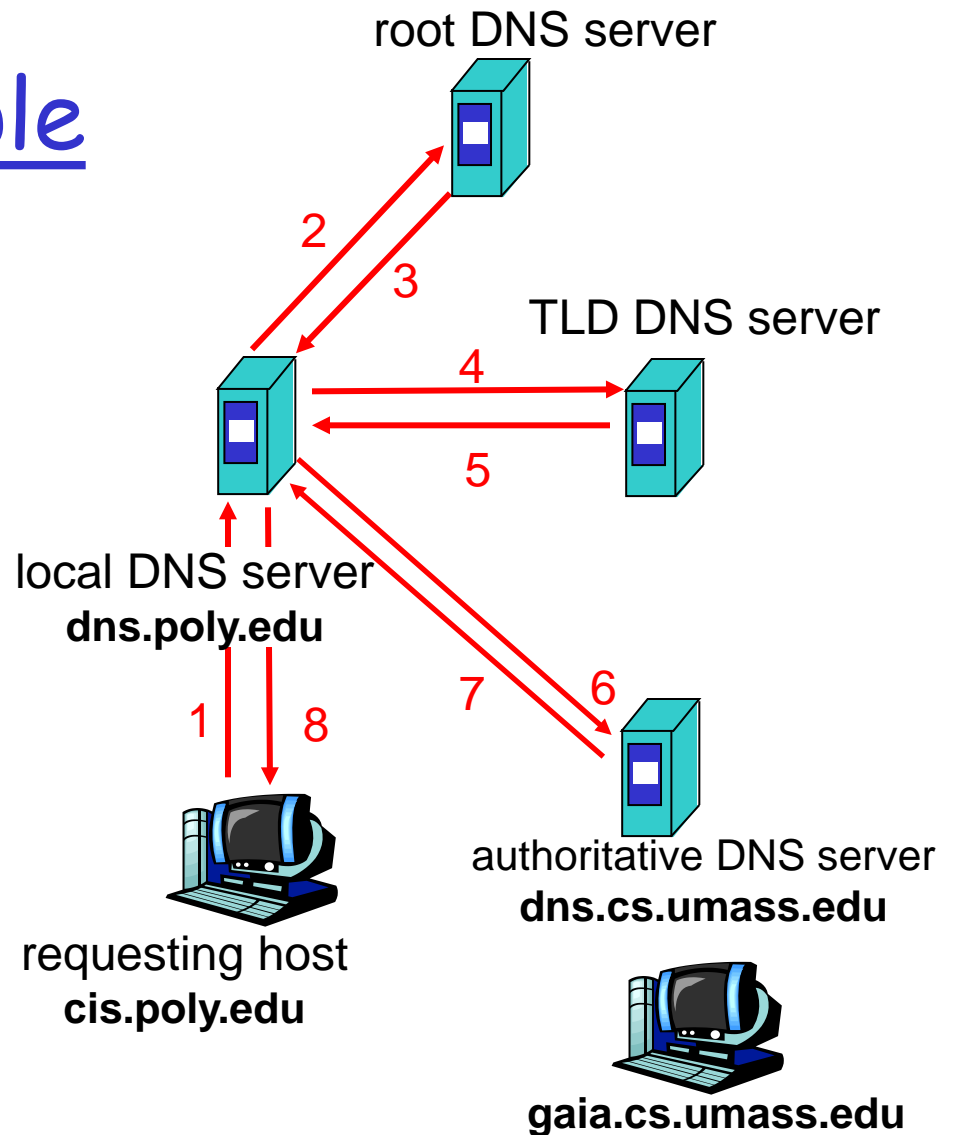
- ❑ Does not strictly belong to hierarchy
 - ❖ Like a "client" (the hierarchy is the "server")
- ❑ Each ISP (residential ISP, company, university) has one
 - ❖ also called "default name server"
- ❑ When host makes DNS query, query is sent to its local DNS server
 - ❖ acts as proxy, forwards query into hierarchy

DNS name resolution example

- ❑ Host at cis.poly.edu wants IP address for gaia.cs.umass.edu

iterated query:

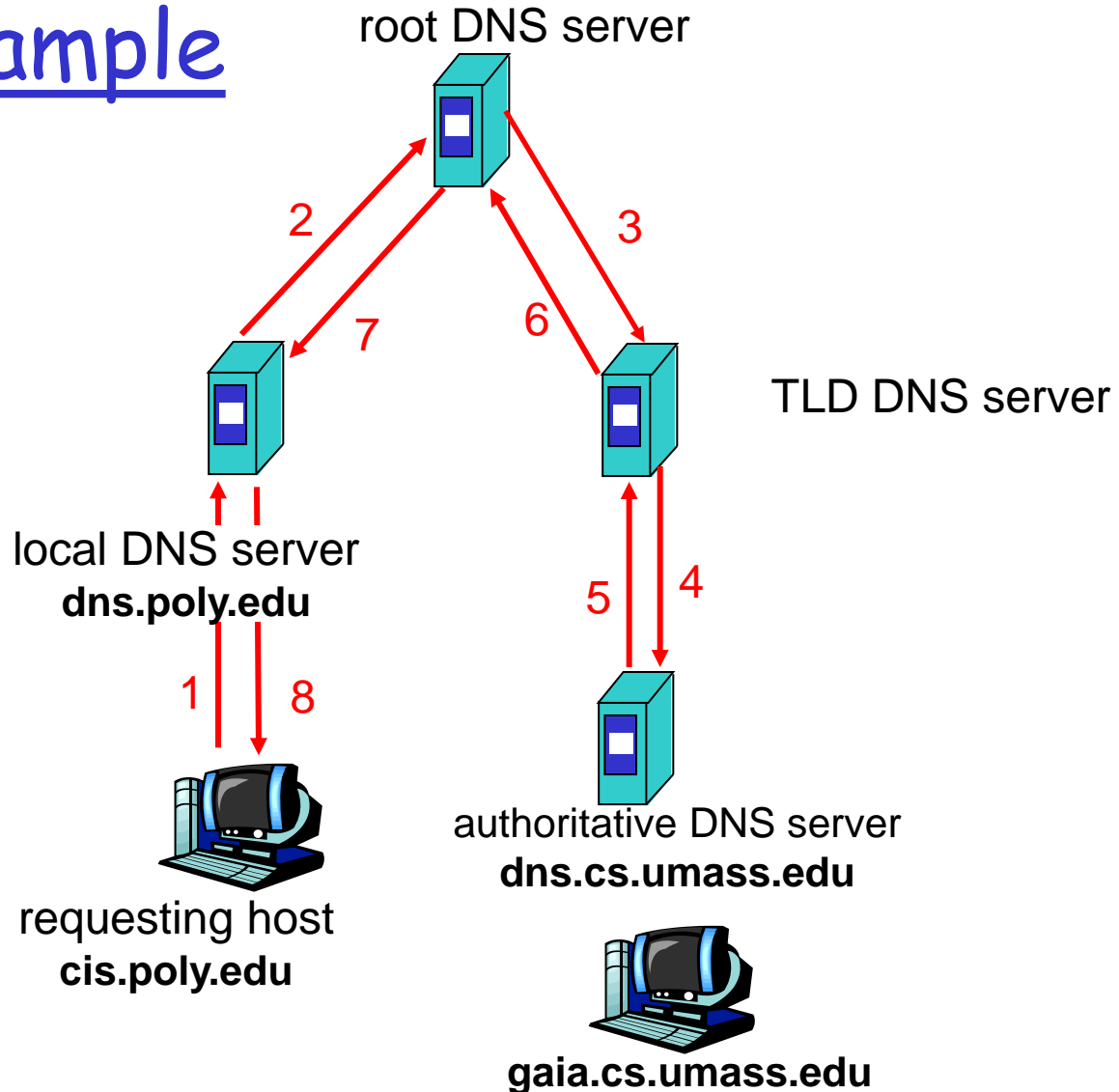
- ❖ contacted server replies with name of server to contact
- ❖ "I don't know this name, but ask this server"



DNS name resolution example

recursive query:

- ❖ puts burden of name resolution on contacted name server
- ❖ heavy load?



DNS: caching and updating records

- Once (any) name server learns mapping, it *caches* the mapping
 - ❖ cache entries timeout (disappear) after some time
 - ❖ TLD servers typically cached in local name servers
 - Thus root name servers not often visited
- Update/notify mechanisms proposed IETF standard
 - ❖ RFC 2136

DNS records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

Type=A (Address)

- name is hostname
- value is IP address

Type=NS (Name Server)

- ❖ name is domain (e.g., foo.com)
- ❖ value is hostname of authoritative name server for this domain

Type=CNAME (Canonical NAME)

- name is alias name for some "canonical" (the real) name
- www.ibm.com is really servereast.backup2.ibm.com
- value is canonical name

Type=MX (Mail eXchange)

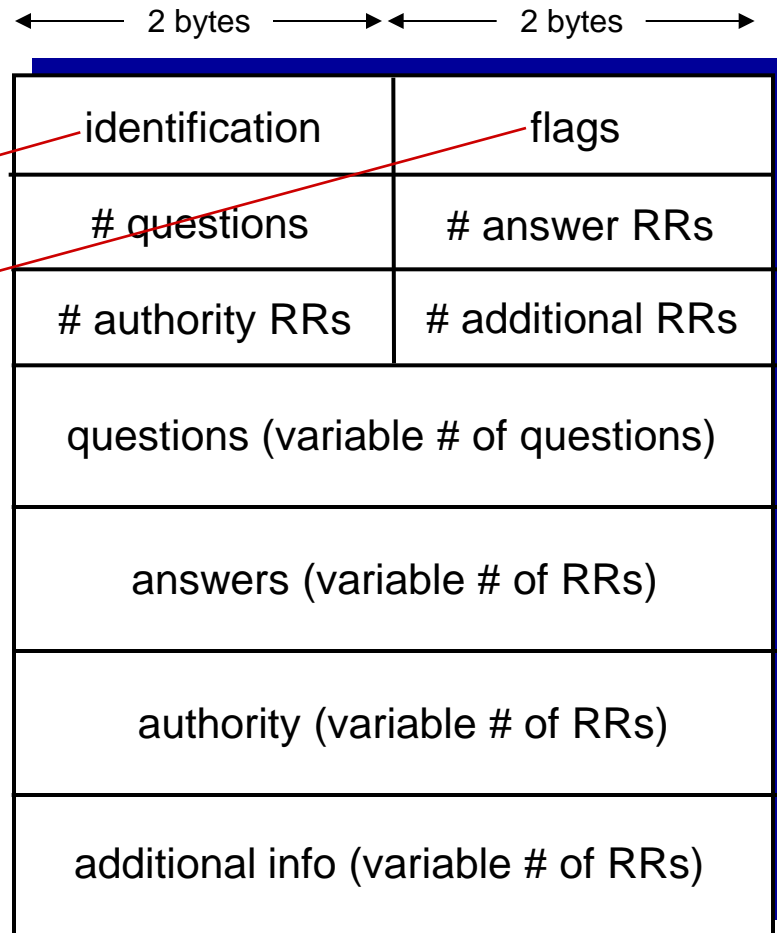
- value is name of mailserver associated with name

DNS protocol, messages

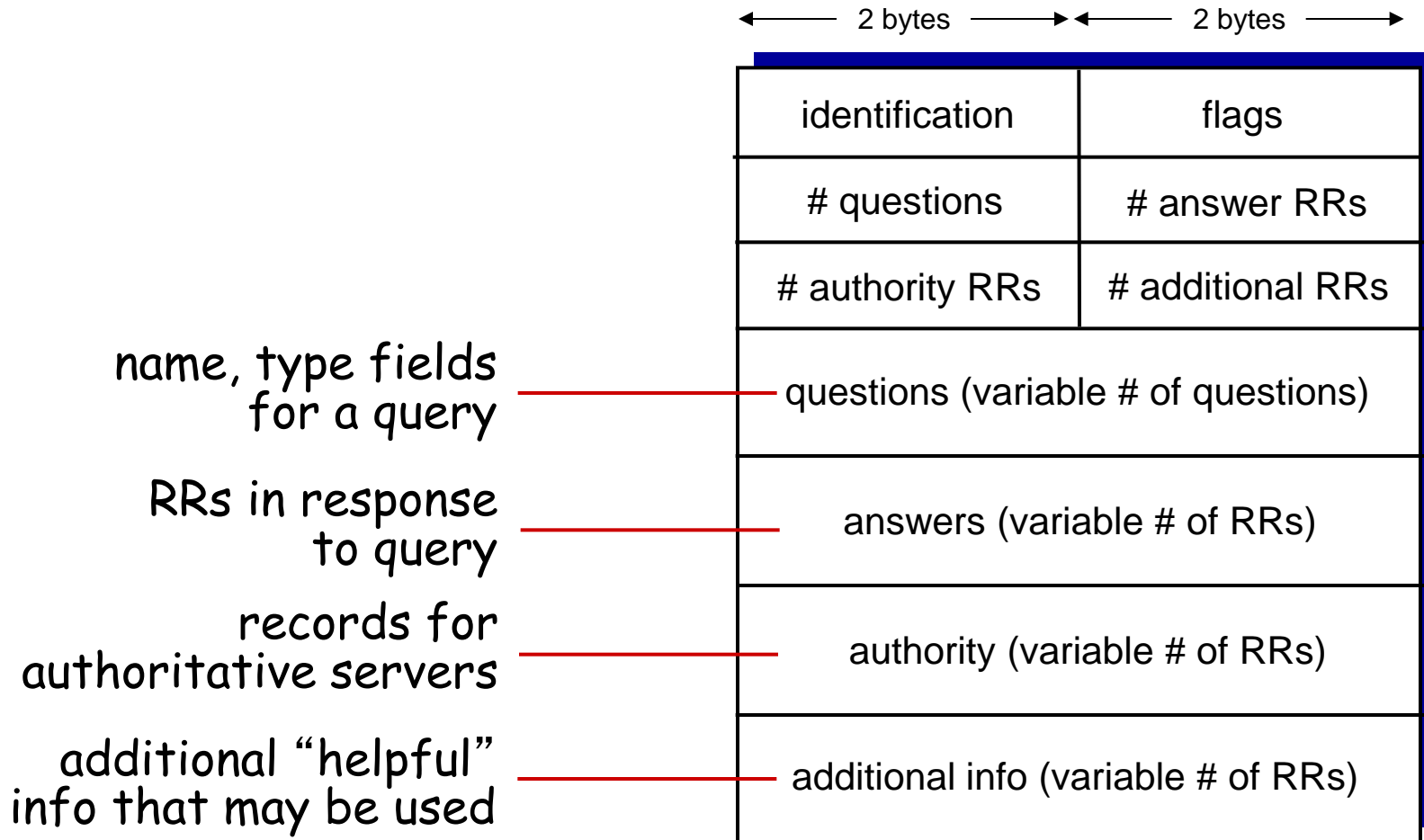
□ *query* and *reply* messages, both with same *message format*

msg header

- ❖ **identification**: 16 bit #
for query, reply to query
uses same #
- ❖ **flags**:
 - query or reply
 - recursion desired
 - recursion available
 - reply is authoritative



DNS protocol, messages



Inserting records into DNS

- ❑ example: new startup "Network Utopia"
- ❑ register name networkutopia.com at *DNS registrar* (e.g., Network Solutions)
 - ❖ provide names, IP addresses of authoritative name server (primary and secondary)
 - ❖ registrar inserts two RRs into com TLD server:

```
(networkutopia.com, dns1.networkutopia.com, NS)
(dns1.networkutopia.com, 212.212.212.1, A)
```

- ❑ create authoritative server Type A record for `www.networkutopia.com`; Type MX record for `networkutopia.com`
- ❑ *How do people get IP address of your Web site?*