# COMP 2280 - Introduction to Computer Systems

Module 4- Branching and Control Statements



- In this module we will learn
  - The LC-3 branching instructions,
  - The LC-3 condition codes,
  - How to implement basic control structures.

#### **Condition Codes**

- CPUs have bits that represent conditions that have occurred during the last arithmetic or logic instruction executed.
- These bits may indicate that the result was zero, negative, positive, overflow, carry, etc.
- Each of these bits is known as a condition code.
- LC-3 has three condition codes
  - ∘ N negative
  - ∘ P − positive
  - $\circ$  Z zero
- Condition codes can be used by branching/control instructions to change the order in which instructions are executed.

#### Condition Codes

- The instructions in LC-3 that affect the condition codes are: add, and, Id (direct),
   Idi (indirect), Idr (relative/base-offset), and not.
- Each time one of these instructions is executed, the three condition codes are set/reset based on the result.
- eg)
  - AND R0,R0,#0; N=0,P=0,Z=I
  - ADD R0,R0,#-I; N=I,P=0,Z=0

#### **Condition Codes**

- Condition codes may be examined by some LC-3 instructions to determine what to do next.
- This is the idea of branching.
- We use this idea in High-Level-Languages all the time.

#### Example:

```
x = y + 10;
if (x > 0) //a branch instr
```

- In order to write useful assembly language programs, we must be able to implement control structures like those we have in HLL.
- Let's look at the LC-3 branching instructions that will allow us to do this.

#### Branching/Control Statements

- These statements are used to alter the flow of control within a program.
- This is done by changing the program counter (PC) register.
- Three types
  - conditional branch based on result of last instruction
  - unconditional branch always (ever heard of "goto"?)
  - trap changes PC to execute an "OS service routine".
     Once trap finishes, returns to execute next instruction.

- Can be used to implement
  - if/else statement
  - while loop
  - for loop
  - do-while loop
  - and other looping structures
- The key when using conditional branches is the condition codes
  - n negative
  - o z zero
  - p positive

The conditional branch instruction has the form

brX label

where X = n, z, p, nz, np, zp, or nzp

eg) This example loops 10 times.

AND R0,R0,#0

ADD R0,R0,#10

loop

ADD R0,R0,#-I BRP loop

- In order to use conditional branching, you usually do some operation that affects the condition codes and use the appropriate branch instruction
  - brn branch if z = 0, n = I, p = 0
  - brz branch if z = I, n = 0, p = 0
  - brp branch if z = 0, n = 0, p = 1
  - brnz branch if (z = I or n = I) and (p = 0)
  - brnp branch if (z = 0) and (n = 1 or p = 1)
  - brzp branch if (z = I or p = I) and (n = 0)
  - brnzp (same as br) branch if n=1 or z=1 or p=1 (this is an unconditional branch)

An Example of an If/Else statement (see if-else.asm).

```
;This example implements a simple if/else statement
          .orig x3000
          and R2,R2,#0
                                       :R2 <- 0
          and RI,RI,#0
          add RI,RI,#10
                                       ;RI < -10
; if RI > 0 R2++, RI-=2 else R2=RI, RI--
          brnz else
                                       ;RI <= 0 do else clause
if
                                       ; if (RI > 0) do if clause
         add R2,R2,#1
          add RI,RI,#-2
          br endif
else
                                       ; else (RI \leq 0)
          add R2,R1,#0
          add RI,RI,#-I
endif
                                       ; end of if/else
         add R3,R1,R2
          halt
          .end
```

- An example of a while loop (see while.asm).
- Read a char and print it out 12 times.

;This example implements a simple while loop

```
.orig x3000
                             ; read char into r0
          trap x20
         and RI,RI,#0
          add RI,RI,#12
loop
          brz done
                             ; while (R1 != 0)
                             ; display char in r0
          trap x21
          add RI,RI,#-I
          br loop
done
                             ; end of while loop
          halt
          .end
```

An example of a do-while loop (see do-while.asm)

;This example implements a simple do-while loop

.orig x3000

trap x20 ;read char into r0

and RI,RI,#0

add R1,R1,#5 ;loop 5 times

;notice no test is done before entering the loop

loop

trap x21 ;print character in r0

add RI,RI,#-I

brp loop

done ;end of while loop

halt

.end

#### Converting an if to a branch

Given this if statement, which would be the equivalent branch?

```
ADD R0, R0, R0
BRz if ;if true, goto the if code
BRnp else ; else, go to else code
if ADD R0, R0, #1 ;if code
BR endif ;skip over the else
else AND R0, R0, #0 ; else code
BR endif ;jump to end of if
endif ADD R0, R0, #2
```

ADD R0, R0, R0

BRnp else ;if false, go to else

ADD R0, R0, #1 ;if code

BR endif ;skip over the else

else AND R0, R0, #0 ; else code

endif ADD R0, R0, #2

Much better!

Too much branching!

#### Converting an if to a branch

- When choosing an if condition it's better to choose one that results in the least number of "else" branches.
- The reason for this is exactly what we saw in the last slide.
  - Branch if the condition is **false**, otherwise keep running the code as usual.
- Here is an example of a multi-condition if (easily written as a nested if):

```
ADD R3, R1,#-10 ;check if R1>10
BRNZ endif ;not >10, we skip
ADD R3, R2,#2 ;check if R2==-2
BRNP endif ;was not -2, we skip
;do stuff here
;do stuff here
endif
;last thing here
```

```
if(x>10){ //check if R1>10
    if(y==-2){ //check if R2==-2
        //do stuff here
        //do stuff here
        //do stuff here
     }
}
//last thing here
```

#### Note:

- branching is PC-relative, ie. target address is made by adding signed offset (IR[8:0]) to current PC. (More meaningful when we see addressing modes)
- PC has already been incremented by FETCH stage.
- Target must be within -255 and 256 words of BR instruction.
- This is because we only have 9 bits to hold this offset.



- When a branch is not taken, the next sequential instruction is executed.
- Some instructions do not affect the CC.
   In that case, the CC is unchanged.

Assembler Inst.

```
JMP Rx
```

- Jump is an unconditional branch -- <u>always</u> taken
  - Target address is the contents of a register
  - Copies the register contents to the PC
  - Allows any target address
- E.g.:

```
LD R4,EXITPTR

JMP R4; target can be anywhere

EXITPTR .fill EXIT
```

#### Or simply use:

```
BR EXIT ; target must be within range
```



- Why have both a JMP and a BR?
  - JMP can go beyond the -255 to 256 distance, whereas BR can't.
  - JMP can jump to anywhere in memory.



- A character can be read from the keyboard using trap x20 (or GETC)
- The character is read into the low 8 bits of register R0.
- The top 8 bits of R0 is set to 0...0.

### Printing to Console

- A trap can be used to print a string to the console.
- The trap # is x22. You can also use PUTS.
- The address of the string needs to be stored into R0 (use LEA instruction)

```
Eg)
```

```
.orig x3000
lea R0,strHi ;get addr of string
    trap x22 ;print string
...
    halt
strHi .stringz "Hello\n";string
.end
```