

Counting Peg Solitaire Solutions

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1 The Algorithm

To count the number of ways of solving a peg solitaire problem, I used a depth first search approach. This means I would make as many moves as possible in a row before backtracking when the game has been completed or no new moves can be made.

I used this approach instead of a breadth first search as the stack used in DFS stays relatively short as opposed to the queue used in BFS, which for a grid with 5 rows ran out of memory almost immediately. This is because BFS adds every possible move to the queue at a given position which means it gets very large very quickly. If the shortest path is required, then BFS would be better as the first result it finds would be the shortest.

Another important change I have made is I have ignored multi-move turns. This is because I don't think that doing 2 moves in one turn counts as a separate solution, as it is just combining the moves in another solution. It also speeds up the program by a lot and keeps the stack empty.

The code was written in Scala, where I created a class for the board that contains the current position of the pieces. This was then used in a DFS algorithm. For convenience, the numbers have been altered so that the counting starts from 0, so the top position is marked by the number 0 and the bottom right is marked by $\frac{n(n-1)}{2} - 1$ when there are n rows.

2 The Results

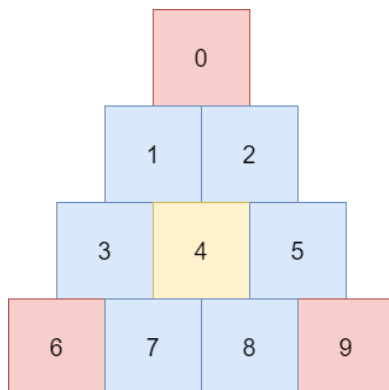
2.1 4 Rows

When there are 4 rows on the board, the starting empty squares can be split into three equivalence classes, where squares in the same class are related by symmetry:

$$E_1 = \{0, 6, 9\}$$

$$E_2 = \{1, 2, 3, 5, 7, 8\}$$

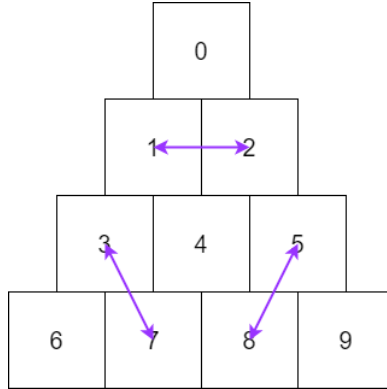
$$E_3 = \{4\}$$



The classes E_1 and E_3 have no solutions. However, the 6 squares in E_2 are able to be solved, and each starting square has a unique empty square. The following table shows the number of ways we can get from a starting empty square to a ending filled square:

Start	End									
	0	1	2	3	4	5	6	7	8	9
0	-	-	-	-	-	-	-	-	-	-
1	-	-	14	-	-	-	-	-	-	-
2	-	14	-	-	-	-	-	-	-	-
3	-	-	-	-	-	-	-	14	-	-
4	-	-	-	-	-	-	-	-	-	-
5	-	-	-	-	-	-	-	-	14	-
6	-	-	-	-	-	-	-	-	-	-
7	-	-	14	-	-	-	-	-	-	-
8	-	-	-	-	14	-	-	-	-	-
9	-	-	-	-	-	-	-	-	-	-

This can also be shown on a diagram:



I also calculated the shortest sequence of moves, where more than one jump can be made per move, using a BFS algorithm. There were two sequences of shortest length starting from square 1 being empty:

$$6 - 1 \ 0 - 3 \ 8 - 6 - 1 \ 5 - 0 - 3 - 5 \ 9 - 2$$

$$6 - 1 \ 0 - 3 \ 8 - 6 - 1 \ 5 - 3 - 0 - 5 \ 9 - 2$$

2.2 5 Rows

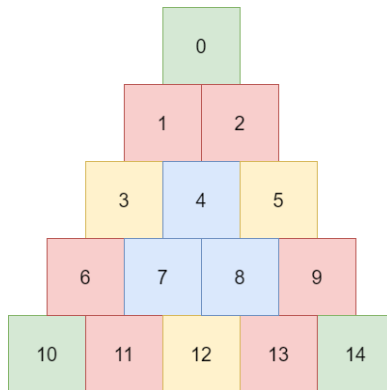
The results are much more interesting when there are 5 rows. In this case it is always possible to solve the problem no matter which square was taken out first. There are 4 equivalence classes in this case:

$$E_1 = \{0, 10, 14\}$$

$$E_2 = \{1, 2, 6, 9, 11, 13\}$$

$$E_3 = \{3, 5, 12\}$$

$$E_4 = \{4, 7, 8\}$$



There are thousands of solutions for each starting square, and all of the squares apart from those in E_4 can reach more than one ending square:

Start	End														
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	6816	-	-	-	-	-	3408	-	-	3408	-	-	16128	-	-
1	-	720	-	-	-	8064	-	-	-	-	3408	-	-	2688	-
2	-	-	720	8064	-	-	-	-	-	-	-	2688	-	-	3408
3	-	-	8064	51452	-	-	-	-	1550	-	-	8064	-	-	16128
4	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
5	-	-	-	-	-	-	-	-	-	-	-	-	1550	-	-
6	-	8064	-	-	-	15452	-	1550	-	-	16128	-	-	8064	-
7	-	-	-	-	-	1550	-	-	-	-	-	-	-	-	-
8	-	-	-	1550	-	-	-	-	-	-	-	-	-	-	-
9	3048	-	-	-	-	-	2688	-	-	720	-	-	8064	-	-
10	-	3048	-	-	-	16128	-	-	-	-	6816	-	-	3408	-
11	-	-	2688	8064	-	-	-	-	-	-	-	720	-	-	3408
12	-	8064	-	-	-	-	-	1550	-	-	16128	-	51452	8064	-
13	-	2688	-	-	-	8064	-	-	-	-	3408	-	-	720	-
14	-	-	3048	16128	-	-	-	-	-	-	-	3048	-	-	6816

We can also put these on 4 diagrams:

