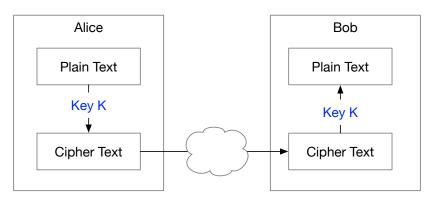
# Introduction to blockchain With bitcoin as an example

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# Symmetric cryptography

Consider symmetric cryptoraphy, where



The problem: how do I securely transfer key K over the network?

# Asymmetric cryptography, RSA algorithm

Find 3 very large positive integers *e*, *d*, *n* s.t.

$$(m^e)^d \equiv m \pmod{n}, \ \forall m, \ 0 \leq m \leq n$$

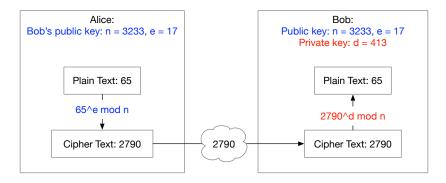
Knowing e, n or even m it's extremely hard to find d.

- ▶ Public key: n, e
- ▶ Private key: n, d
- ▶ Message: *m*

### Operations:

- Encryption(m):  $c = m^e \pmod{n}$
- ▶ Decryption(c):  $m' = c^d \pmod{n} = (m^e)^d \pmod{n}$ ,  $m' = m \pmod{n}$

## RSA algorithm, example



Anyone who wants to talk to Bob can retrieve Bob's public key, use it to encrypt the message, and know that only holder of the corresponding private key can decrypt.

## RSA algorithm in practice

In practice, key pairs are much longer.

\$ ssh-keygen -t rsa -b 4096

Each key pair corresponds with an **identity**.

#### The two operations:

- Encryption/Decryption: Alice uses Bob's public key to encrypt, Bob uses his private key to decrypt
- Signing/Validation: Alice uses her own private key to sign, others use Alice's public key to verify

Vs symmetric encryption: more computation, but solves the problem of key transfer (and many others)

Related concepts: digital signature, certificate, public key infrastructure



## To design a cryptocurrency

## Imagine designing a cryptocurrency

- An account is a public/private key pair!
- ▶ If I pay someone (a public key identity), I sign with my private key: others can verify I made the payment, and I cannot refute later on

But how does anyone know I have enough money left to make the payment? *Double spending* 

- The centralized way: we all agree on (and preinstall) having one party to trust, who keeps track of everyone's account balances
- ► The decentralized way: is it possible to have all of us keep track of account balances together?

## To design a cryptocurrency - cont

#### Considerations:

- Distributed
- Trustless, no a-prori trust relationship established
- ► Consensus, everyone agrees on account balances

#### Why is this problem unique:

- ▶ I don't trust anyone I talk to! (think Raft, dynamo, etc)
- But we can still agree on something

#### Distributed trustless consensus - intuition

Distributed consensus: leadership election, and  ${\bf state}$ -machine  ${\bf replication}^1$ 

If we can agree on the entire history of transactions (state machine), we know if someone is trying to double spend. <sup>2</sup>

To agree on the history,

- what if we assume: there are more peers in the network who are honest, than those who are not
- we can have everyone tell everyone else their view of the entire history, and hope they converge...

<sup>&</sup>lt;sup>1</sup>Think Raft consensus algorithm

<sup>&</sup>lt;sup>2</sup>Think log structured merge tree (e.g. writes in Cassandra) / journaling file system / git 

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# Blockchain, introduction<sup>3</sup>

Blockchain makes a different assumption

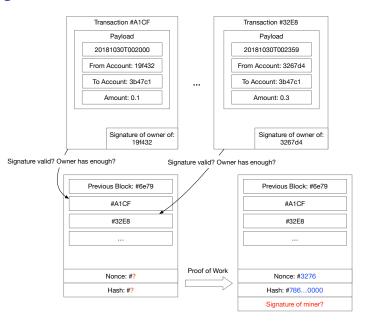
honest peers own more computation power than those who are dishonest

Consequently, if it takes computation power to grow the history, then honest peers can grow the history faster than dishonest ones

- Peers trust the chain with the most computation power invested (i.e. the longest), and grow based on that
- ► In order to grow the history, peers perform a **proof of work**, which is computationally non-trivial

 $<sup>^3</sup>$ As in the classic bitcoin whitepaper, since the field has seen lots of divergents

## Building a block



## Proof of work

After assembling X transactions in a block<sup>4</sup>, find a nonce value to attach to the block, s.t. the hash<sup>5</sup> of the entire block has an agreed number of 0s at its end

- A block is only considered valid and complete, if one such nonce is attached
- No known efficient algorithm exists to find one such value, (mostly) trials
- Alternatives exist

Proof-of-work "verifies" a block (in turn, grows a chain)

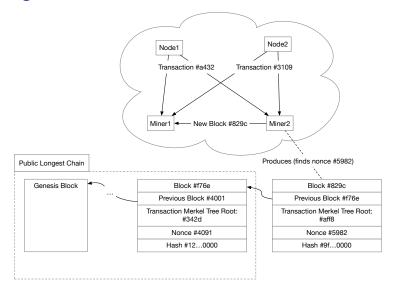
Peers who carry out proof-of-work are called **miners**.



<sup>&</sup>lt;sup>4</sup>X depends on the size of transactions, and is usually above 1000 in bitcoin

<sup>&</sup>lt;sup>5</sup>E.g. sha256

## Growing a chain



Genesis block: the start of the chain (e.g. allocating X coins to the creator)

# Mining and incentive

Why should anyone verify someone else's transaction (mine)?

- Miners are awarded a fixed number of bitcoins for each block mined
- Transactions can attach a transaction fee
- Winner takes all

How are new coins introduced to the system?

- Genesis block
- Mining 'adds' new coins to the system (analogy: miners spend work to find gold nuggets to add to world gold circulation)

## Operation

What if two blocks are created at the same time?

A temporary fork. Can introduce an ordering mechanism, or miners may work on different forks, until one gets longer than the other and propagates.

As the system progresses,

- ► Number of 0s to satisfy proof-of-work increases <sup>6</sup>
- Mining reward halves every X blocks<sup>7</sup>
- ► A hard fork could be made for backward incompatible upgrade, or confirmed compromise

Max number of coins is fixed (inflation-free!)

$$\lim_{n\to\infty}\sum_{i=1}^n\frac{1}{2^i}=1$$

<sup>&</sup>lt;sup>6</sup>The system aims to create 1 new block every 10 minutes, and it does so by adjusting the difficulty of proof-of-work

<sup>&</sup>lt;sup>7</sup>Current reward: 12.5 btc, X = 210000. How to incentivize afterwards: transaction fee

## Tamper resistance and privacy

#### Temper resistance

- Assume dishonest peers work together to produce proof-of-work for a block in which a transaction source account double spends
- ▶ Dishonest peers have to grow the chain faster than the honest ones, so that the honest ones use their chain
- Subverting a previous block in the chain gets exponentially more difficult (grow everything afterwards, and faster)
- ▶ If you have this much computation power, it'd be more in your interest to mine, rather than to subvert

## Privacy

- ▶ No tie between account 736f with a physical identity
- The network does not store anything related with physical identity



### The trade-offs

#### Distributed and trustless come at a cost

- ▶ Global scale broadcast messages: vs unicast to a bank
- ► Proof-of-work: vs bank keeping records and trusting what they keep
- Transaction fee: pay to verifier, or pay to bank

#### But also with advantages

- Less a-priori
- Anonymous
- Mathematically hard to tamper / counterfeit

# Analogy to a currency

- Durable
- Portable
- Divisible
- ► Fungible (like symmetric in an equivalence relation)
- ► Intrinsic value?

#### Extensions - smart contract

With bitcoin, the *actions* we record on the chain are transactions. They can be anything else.

- Casting a vote for a presidential candidate
- Publishing a blog post
- **.**..

How do we express *anything*? The ethereum framework and Solidity programming language.

## Summary

- Public key cryptography
- Distributed trustless consensus
- Computation power assumption
- ► Proof-of-work

## References