Pattern Recognition with C++ OpenMP and CUDA

Giulio Bazzanti Niccolò Biondi

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Table of Contents

- Introduction
- Sequential
- OpenMP
- CUDA
- 6 Experiments
- 6 Conclusions

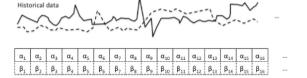
Introduction

SAD and Pattern Recognition

Pattern Recognition

Introduction

- Search some Queries inside Historical data
- The goal is to find the data subset which matches query
- Difficult to find the exact subset





γ1	γ ₂	γ3	γ ₄	γs	γ ₆
δ_1	δ2	δ_3	δ_4	δ_5	δ_6



Introduction

Parallel Pattern Recognition

- Sum of absolute differences (SAD)
 - Similarity measure between data subsets
 - Calculated by the absolute difference between each subset values
 - Find the nearest data subset
- SAD computation for Patter Recognition can be parallelized, since

$$R_q[i] = \sum_{j=0}^{len(q)} |D[i+j] - q[j]|$$



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Parallel solutions

CUDA

Introduction 000

OpenMP



Sequential

Sequential approach

Algorithm 1 Pattern Recognition

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- 1: Generate the Historical Data
- 2: Generate all the Queries
- 3: Create a vector R with length equal to Queries number
- 4: **for** query *q* in Queries **do**
- Create a vector R_q with length len(Data) len(q) + 1 5:
- Compute $R_q[i] = \sum_{i=0}^{len(q)} |D[i+j] q[j]|$
- $R[q] = arg \min R_q$
- 8: end for
- 9: return R



Sequential approach

Sequential 000

- Analyze all queries
- Scrolls across the Historical Data
- Has to find min SAD value for each guery
- Total computation time $\Rightarrow O(nrq) + O(nq)$



Sequential approach

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Problems

- very long Historical Data
- a lot of queries



${\sf OpenMP}$

OpenMP approach

- Use OpenMP library to leverage the CPU multi-threads computation
- Use OpenMP for loop, critical section, parallel section

Parallelism levels

- Parallelism on Queries
- Parallelism on Historical Data



Parallelism on queries

- Is the higher parallelism level in Pattern Recognition
- Each thread computes one query thought a OpenMP for loop
- Each thread analyzes its query like in the Sequential approach
- Each thread computes the min SAD value for the query



Parallelism on data

- Use parallel section to analyze Historical Data in parallel
- Each thread computes SAD values w.r.t. the same Query
- Threads wait the other threads termination through implicit barrier

Different approaches

- Privatization method
- Lock method



Parallelism on data

Privatization method

- Private variables for each thread
- Thread stores min SAD privately
- Each result vector (R_q) length depends on threads number

Lock method

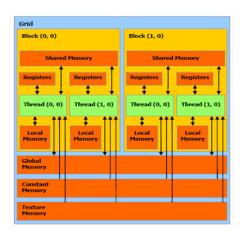
- No minimum SAD value search
- Shared variable stores query min SAD
- Critical section (with flush) to write SAD values



CUDA

CUDA approach

- Each block has a Shared memory
- Each tread has a personal and block id
- The total blocks number depends on Historical Data length and threads per block
- Four different implementations





Naive Implementation

- The Queries and Historical Data are stored in Global memory
- Each thread reads a Historical Data chuck and compare it with all queries
- The Result vector is in global memory



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A lot of reads/writes in/from Global memory



Private Implementation

- The Queries and Historical Data are stored in Global memory
- Each thread has a private variable
- The threads compute SAD values w.r.t. its data chunk
- Make local SAD computations and write once



CUDA

Private Implementation

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Constant Implementation

- Use the Constant memory to store the Queries
- In each block Shared memory is written a Historical Data chunk
- Fach block thread loads a value from Historical Data



Constant Implementation

- Use the Constant memory to store the Queries
- In each block Shared memory is written a Historical Data chunk
- Each block thread loads a value from Historical Data
- A thread can see only a fragment of data to compute SAD
 - Each thread in each block compute the remaining SAD values for a query

$$R[idx - r + q \cdot len(R)] + = Ds[threadId] - Qc[r + q \cdot len(Q)]$$



Constant Implementation

- Use the Constant memory to store the Queries
- In each block Shared memory is written a Historical Data chunk
- Fach block thread loads a value from Historical Data
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Possible race condition \Rightarrow atomic write



Tiling Implementation

- Both Historical Data and Queries are stored in Shared memory
- Each thread loads in Shared memory Historical Data and Queries chunks



Tiling Implementation

- Both Historical Data and Queries are stored in Shared memory
- Each thread loads in Shared memory Historical Data and Queries chunks
- For each Historical Data chuck threads compute SAD w.r.t. all the Queries
 - Need multiple Queries loads in Shared memory (at least 1 for each Query)
 - Different Query portions loaded in the same block Shared memory (t < q)

$$R[idx - r + p \cdot T + q \cdot len(R)] + = |Ds[threadId] - Qs[r]|$$



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Experiments

One Historical Data vector and multiple Queries

- Pattern Recognition tested on:
 - Variable Historical Data length
 - Variable Query length
 - Variable Queries number
 - Variable CPU and GPU threads

Hyper-parameter	Default Value
Historical Data	100000
Query	1000
Queries Number	10
Threads Number	12
Kernels per Block	128



Testing Hypothesis

One Historical Data vector and multiple Queries

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Stabilize results

For each test configuration are performed 10 run for stable results

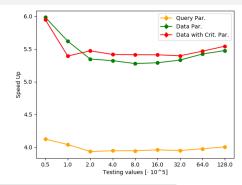


 Sequential
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Historical Data length

- Historical Data with size in [50000, 1280000]
- Table compares the OpenMP and CUDA implementations w.r.t. the Sequential one (with mean and std)

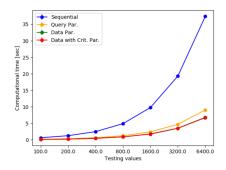


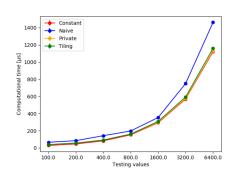
	len(Data)	Sequential [sec]	OpenMP [sec]	CUDA [µsec]
Î	50000	$3.20 \pm 0.37\%$	$0.53 \pm 0.02\%$	$178.6 \pm 5.81\%$
	200000	$12\pm1.03\%$	$2.24\pm0.52\%$	$181.3 \pm 5.51\%$
	800000	$47.8\pm0.02\%$	$9.05\pm1.02\%$	$171.2 \pm 8.61\%$
	320000	$191\pm0.34\%$	$36\pm1.21\%$	$182.8 \pm 2.91\%$
	1280000	$780 \pm 6.58\%$	$142\pm2.57\%$	$181.3 \pm 3.72\%$



Query length

- Query length from 100 to 6400
- ullet CPP computational times reported in sec; GPU ones in μsec
- No parallelism on queries ⇒ **proportional** grown

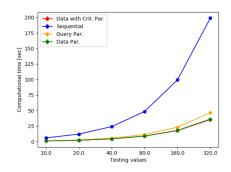


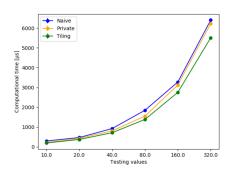




Queries number

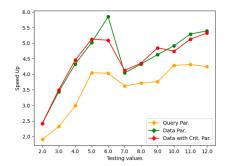
- Queries number in [10, 320]
- CUDA Constant implementation can not be tested

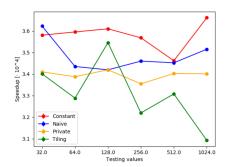






- Threads number from 1 (for Relative Speedup) to 12
- GPU kernels per block from 32 to 1024 (device maximum)







	OpenMP				CUDA			
Metric	Query	Privatization	Lock	Naive	Private	Tiling	Constant	
S_p	4.03	5.87	5.09	$2.2 \cdot 10^4$	$2.8 \cdot 10^4$	$3.0 \cdot 10^4$	$3.4 \cdot 10^4$	
RS_p	4.37	4.48	4.15	(-)	1.24	1.34	1.51	
E	0.67	0.97	0.85	0.22	0.28	0.30	0.34	

- Reported performances for the parallel Pattern Recognition solutions. All the values are computed with a 100000 Historical Data vector, 1000 Query length, 10 total Queries, 6 CPU threads number and 128 CUDA kernels for each block.
- Relative Speedup is computed w.r.t. the baseline methods (Sequential and Naive)
- The values are the mean obtained in 10 different runs



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Conclusions

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- Experimental results underline the gap between CPU and GPU parallel implementations
- OpenMP Privatization method reports almost linear Speedup and Efficiency close to 1
- CUDA show its parallelism power in all the implementations, in particular thanks to the Shared and Constant memories



Thanks for the attention