

How to compose point-free functions

utashih

Warm-up: eta-reduction

- The following two functions are equivalent under eta-conversion:

$\lambda x \rightarrow f\ x$

and

f

Warm-up: eta-reduction

- Haskell's *partially applying* functionality enables us to remove the last (rightmost) parameter of *curried* functions

```
sum xs = foldr (+) 0 xs
```

By applying eta-reduction

```
sum = foldr (+) 0
```

Point-free style

- also called *tacit programming*
- (*from Wikipedia*) Function definitions do not identify the arguments (points) on which they operate; they merely compose other functions
- considered unnecessarily obscure
- ...and we're going into this obscurity!

Examples from H-99

```
last = foldr1 (const id)
reverse = foldl (flip (:)) []
```

```
elementAt :: [a] -> Int -> a
elementAt = (last .) . flip take
```

```
isPalindrome :: (Eq a) => [a] -> Bool
isPalindrome = (==) <*> reverse
```

Composition of functions

- $(.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$
- $(f \ . \ g) \ x = f \ (g \ x)$
- will be used ubiquitously

Example from Codewars

- WeIrD StRiNg CaSe

```
toWeirdCase :: String -> String
```

```
toWeirdCase =
```

```
    unwords . map (zipWith ($) weird) . words
```

```
    where weird = cycle [toUpper, toLower]
```

Digression: functor

- types that can be mapped over (with restrictions)

`class Functor f where`

`fmap :: (a -> b) ~> f a -> f b`

Hom functor

- For each type r , the (covariant) hom functor $Hom(r, -)$ maps
 - each type a to type $r \rightarrow a$
 - each function of type $a \rightarrow b$ to a function of type $(r \rightarrow a) \rightarrow (r \rightarrow b)$

`fmap :: (a -> b) ~> (r -> a) -> (r -> b)`

Hom functor

- a more familiar name would be the Reader functor

```
instance Functor ((->) r) where  
    fmap = (.)
```

Example: concatMap

```
concatMap :: (a -> [b]) -> [a] -> [b]  
concatMap = (concat .) . map
```

- Or equivalently

```
concatMap f xs = concat (map f xs)
```

Example: concatMap

- To validate the equivalence

`concatMap f xs = concat (map f xs)`

`concatMap f xs = concat $ map f xs`

`concatMap f xs = concat . map f $ xs`

`concatMap f = concat . map f`

`concatMap f = (.) concat (map f)`

`concatMap f = (.) concat $ map f`

`concatMap f = (.) concat . map $ f`

`concatMap = (.) concat . map`

`concatMap = (concat .) . map`

Example: concatMap

- The key idea is to find the rightmost parameter by changing (nested) function applications into single composed function applied to single argument

```
concatMap f xs = concat (map f xs)
```

```
concatMap f xs = concat $ map f xs
```

```
concatMap f xs = concat . map f $ xs
```

```
concatMap f      = concat . map f
```

Example: concatMap

- Note that it is incorrect to take the second step like

`concatMap f xs = concat $ map f xs`

`concatMap f xs = concat . map $ f xs`

- The latter formula corresponds to

`concat (map (f xs))`

- Function application associates to the left while (\$) associates to the right; thus the outermost application is `(map f) xs` instead of `map (f xs)`

Example: elementAt

```
elementAt :: [a] -> Int -> a  
elementAt = (last .) . flip take
```

```
(last .) :: (r -> [a]) -> (r -> a)  
flip take :: [a] -> (Int -> [a])
```

...so we have more dots!

- Composing function with 2 arguments

$$(. :) :: (c \rightarrow d) \rightarrow (a \rightarrow b \rightarrow c) \rightarrow a \rightarrow b \rightarrow d$$
$$(. :) = (.) \ . \ (.)$$
$$f \ .: g = \lambda x \ y \rightarrow f \ (g \ x \ y)$$

- Thus we have

$$\text{concatMap} = \text{concat} \ .: \text{map}$$

- This operator is defined in Data.Composition in the *composition* package

Back to the hom functor

- Of course, it is an applicative functor/monad as well
- Some more definitions:

```
instance Applicative ((->) r) where
```

```
    pure = const
```

```
    f <*> g = \r -> f r (g r)
```

```
instance Monad ((->) r) where
```

```
    f >>= g = \r -> g (f r) r
```

Example from H-99

```
isPalindrome :: (Eq a) => [a] -> Bool  
isPalindrome = (==) <*> reverse
```

```
isPalindrome = \x -> x == reverse x  
              = \x -> (==) x (reverse x)  
              = \x -> (==) <*> reverse $ x  
              = (==) <*> reverse
```

Example from Codewars

- Equal Sides of An Array

```
findEvenIndex :: [Int] -> Int
findEvenIndex = fromMaybe (-1) . elmIndex True .
    (zipWith (==) <$> scanl1 (+) <*> scanr1 (+))
```

- Note that `fmap`, `(<$>)` and `(.)` serve the same purpose

Example from Codewars

- Sort the Odd

```
sortArray :: [Int] -> [Int]
```

```
sortArray = replaceOdd <$> id <*> sort .  
    filter odd
```

```
replaceOdd :: [Int] -> [Int] -> [Int]
```

Example from Codewars

- Are they the “same”

```
comp :: [Integer] -> [Integer] -> Bool
```

```
comp = (. sort) . (==) . sort . map (^2)
```

```
comp xs ys = sort (map (^2) xs) == sort ys
```

Example: map filter

```
mp :: (b -> Bool) -> (a -> b) -> [a] -> [b]
mp p f xs = filter p (map f xs)
mp p f     = filter p . map f
mp         = (. map) . (.) . filter
```

Partial applying to $(.)$

- Putting it off

$f :: (a \rightarrow b)$

$(f \ .) :: (r \rightarrow a) \rightarrow (r \rightarrow b)$

- Bringing it forward

$g :: (a \rightarrow b)$

$(. \ g) :: (b \rightarrow s) \rightarrow (a \rightarrow s)$

In all: understanding (.)

- Composition of functions
- Functor map of hom functor
- Continuation-like

“Boilerplate” pattern

```
func x y = f (g x y)
func      = (f .) . g
```

```
func x = f (g x) (h x)
func    = f <$> g <*> h
        = f . g <*> h
```

```
func x y = f (g x) (h y)
func      = (. h) . f . g
```

Try it out

```
agreeLen :: (Eq a) => [a] -> [a] -> Int
agreeLen x y = length $
    takeWhile (\(a, b) -> a == b) (zip x y)
```

- Check your answer at <http://pointfree.io/>

Typical misuse

```
mapWithIndex :: (Int -> a -> b) -> [a] -> [b]
mapWithIndex =
    (snd .) . ('mapAccumL' 0) . ((.) . (,) . (+ 1) <*>)
```

- This should be

```
mapWithIndex f xs =
    snd $ mapAccumL (\i x -> (i+1, f i x)) 0 xs
```

```
mapAccumL :: (a -> b -> (a, c)) -> a -> [b] -> (a, [c])
```

Q&A