

**LEMMA 1.**

$$\sum_{j \leq x} \Lambda(j) E(x/j) = \sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k \quad (x \geq 2).$$

*Proof.* By the Lemma of II.3 (sum of a convolution),

$$\begin{aligned} S := \sum_{j \leq x} \Lambda(j) E(x/j) &= \sum_{j \leq x} [\Lambda * (u * \nu)](j) \quad (\text{Lemma: } E \text{ sum-function of } u * \nu) \\ &= \sum_{j \leq x} (\ell * \nu)(j) \quad (\Lambda * \nu = \ell) \\ &= \sum_{j \leq x} \nu(j) \sum_{k \leq x/j} \log k \quad (\ell = \log; \text{Lemma again}) \\ &= \sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k \quad (x \geq 2). \quad // \end{aligned}$$

**LEMMA 2.**

$$\psi(2n) \geq \log \binom{2n}{n}.$$

*Proof.* Take  $x = 2n$  in Lemma 1. As each  $E(\cdot) \leq 1$ ,

$$S \leq \sum_{j \leq 2n} \Lambda(j) = \psi(2n).$$

But

$$\sum_{k \leq x} \log k - 2 \sum_{k \leq x/2} \log k = \sum_{k=n+1}^{2n} \log k - \sum_{k=1}^n \log k = \log \left( \frac{(n+1)(n+2) \dots (2n)}{1 \cdot 2 \dots n} \right) = \log \binom{2n}{n}. \quad //$$

**Th. 3 (Chebyshev's Lower Estimates).** For  $\epsilon > 0$  and  $x$  large,

- (i)  $\psi(x) \geq (\log 2 - \epsilon)x$ ;
- (ii)  $\theta(x) \geq (\log 2 - \epsilon)x$ ;
- (iii)  $\pi(x) \geq (\log 2 - \epsilon)li(x)$ .

*Proof.* (i) Let  $N := \binom{2n}{n}$  as above. This is the largest of the  $2n+1$  terms in the binomial expansion of  $(1+1)^{2n} = 2^{2n}$  (by Pascal's triangle), so  $2^{2n} \leq (2n+1)N$ . So by the Lemma above,

$$\psi(2n) \geq \log N \geq 2n \log 2 - \log(2n+1).$$

Given  $x$ , take  $n$  with  $2n \leq x < 2n + 2$ . Then (i) follows as

$$\psi(x) \geq \psi(2n) \geq (x - 2) \log 2 - \log(x + 1).$$

(ii) This follows from (i) as  $(\psi(x) - \theta(x))/x \rightarrow 0$  (Cor. above).

(iii) This follows from (ii) by the first Theorem of this section. //

**Cor. 5.**  $\pi(x) \geq (\log 2 - \epsilon)x / \log x$ .

*Proof.*  $\psi(x) \leq \pi(x) \log x$  (first Prop. of this section and (i)). //

**THEOREM (Chebyshev, 1849-51)** (Mastery Question, 2013).

$$\liminf \pi(x)/li(x) \leq 1 \leq \limsup \pi(x)/li(x).$$

In particular, if the limit exists, it is 1 (as in PNT).

*Proof.* For all  $\epsilon > 0$  there exists  $x_0$  such that for  $x \geq x_0$

$$\ell - \epsilon \leq \frac{\pi(x)}{x / \log x}, \quad \frac{\pi(x)}{x / \log x} \leq L + \epsilon.$$

For the lower bound, integration by parts gives, as  $0 < \pi(u) \leq u$ ,

$$\sum_{p \leq x} \frac{1}{p} \geq \sum_{x_0 < p \leq x} \frac{1}{p} = \int_{x_0}^x \frac{d\psi(u)}{u} = \frac{\pi(x)}{x} - \frac{\pi(x_0)}{x_0} + \int_{x_0}^x \frac{\pi(t)}{t^2} dt,$$

$$\geq -1 + \int_{x_0}^x \frac{\pi(t)}{t^2} dt \geq -1 + (\ell - \epsilon) \int_{x_0}^x \frac{dt}{t \log t} \geq (\ell - \epsilon) \log \log x + O_\epsilon(1)$$

( $\int^x dt/(t \log t) = \int^x d \log t / \log t = \log \log x$ ). But by Mertens' Second Th.,

$$\sum_{p \leq x} 1/p = \log \log x + c_1 + O(1/\log x).$$

Combining,

$$\log \log x + c_1 + O(1/\log x) \geq (\ell - \epsilon) \log \log x + O_\epsilon(1) : \quad 1 \geq \ell - \epsilon.$$

This holds for all  $\epsilon > 0$ . So  $\ell \leq 1$ . The upper bound is similar but simpler. //

In 1851, Chebyshev also proved *Bertrand's postulate* of 1845: for any  $n \geq 2$  there is a prime  $p$  between  $n$  and  $2n$ ; see 2013 Problems and Solutions 8 for Erdős' elementary proof of 1932.