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Asn1 Electronic Typed Answers

September 25, 2021

1.

We must find constants $c > 0$ and $n \geq 1$ such that

$$1/n^2 \leq c \cdot 1/n$$

$$0 \leq c \cdot 1/n - 1/n^2 \text{ for all } n \geq n_0$$

$$0 \leq 1/n (c - 1/n) \text{ for all } n \geq n_0$$

$$1/n \leq c \text{ for all } n \geq n_0$$

Therefore when c is 1 and n_0 is 1 this equality holds true for all $n \geq n_0$

2.

We must find constants $c > 0$ and $n \geq 1$ such that

$f(n)$ is $O(g(n))$

This means that there exist constants $c > 0$ and $n \geq 1$ such that

$$f(n) \leq c g(n) \text{ for all } n \geq n_0$$

both are positive, so we can multiply both sides by $1/f(n) \cdot g^*(n)$

$$f(n)/f(n) \cdot g(n) \leq g(n) / f(n) \cdot g(n) \text{ for all } n \geq n_0$$

$$1/g(n) \leq c \cdot 1/f(n) \text{ for all } n \geq n_0$$

Therefore, when $c = 1$ and $n_0 = 1$ the inequality holds true for all $n \geq n_0$

3.

We must find constants $c > 0$ and $n \geq 1$ such that $n - 1$ is $O(1)$

To prove by contradiction we will assume that $n - 1$ is $O(1)$ is true and derive the contradiction

If $n - 1$ is $O(1)$ then there are constants $c > 0$ and $n_0 \geq 1$ such that:

$n - 1 \leq c(1)$ for all $n \geq n_0$

$n \leq c - 1$ for all $n \geq n_0$

Therefore, this cannot be true because c is a constant whereas n can grow without bounds, so $n-1$ is not $O(n)$.

4.

Done → MostTimes.java submission

5.

The algorithm does not terminate. The example I will show is in the case we have an array L with $n \% 2 = 0$ with $x = L[-1]$. The problem with the algorithm is the second if statement. The reason we increment i by 2 in the else statement is to check through all the even indexes in the array for x , then in the second if statement we set i to 1 when we are “done with all the even indexes” to then start checking the odd ones in hopes to find x . The problem with this is that if the element is at the last slot in the array if n is even then i will be set back to 1 because it is equal to $n - 1$ and we compare $L[i]$ which is $L[1]$ with x on the current loop/value of i . In the case where n is odd, if it is the last element in the array the same issue occurs, i gets set to 1 and then we compare $L[i]$ which is $L[1]$ with x .

Lets say x that we are looking for is the last element in the array of an even number of elements, take the array: {4, 5, 6, 7, 8, 9}, look at the table below if we are looking for $x = 9$:

Iteration	Statement	i	$L[i]$ in this iteration
1	if $L[i] = x$	0	4
	else do $i += 2$	2	4
	if $i \geq n - 1$	Not Yet	4
2	if $L[i] = x$	2	6
	else do $i += 2$	4	6
	if $i \geq n - 1$	Not Yet	6
3	if $L[i] = x$	4	8
	else do $i += 2$	6	8

	if $i \geq n - 1$	True	8
4	if $L[i] = x$	1	5
	else do $i += 2$	3	5
	if $i \geq n - 1$	Not Yet	5
5	if $L[i] = x$	3	7
	else do $i += 2$	5	7
	if $i \geq n - 1$	Yes and x is currently $L[i]$ after the incrementation in the else statement but is now set to 1 because i is $n - 1$ currently, one more loop just to make sure	7
6	if $L[i] = x$	1	5
	else do $i += 2$	3	5
	if $i \geq n - 1$		5
7	if $L[i] = x$	3	7
	else do $i += 2$	5	7
	if $i \geq n - 1$	Yes and x is currently in index but this will repeat infinitely	7

What I am trying to explain is that i is never allowed to get to the last element in the array because every time it does, the loop is started again with $i = 1$. To fix this if the second if statement was $i > n$. Also in the scenario that x is not in the array, this will cause an infinite loop because i is always being incremented and reset to 1 until x is found.

I coded this out in the IDE and have the solution with my suggested correction, happy to provide it as context. However, my response is that the algorithm does not ALWAYS terminate in the case where x is at the last element in the array L .

6.

This algorithm always terminates because we use variable j to be incremented and reset to try to find x , however, each time we loop through the array i is incremented by 1, and so when we loop through the array $n - 1$ times, the loop will terminate and return -1 because x was not found.

Additionally, this algorithm does not produce the correct output for any L with x as its last element. Like the algorithm in question 5, we are still checking the even indexes for x first then we are checking the odd ones. However, when j gets incremented to the last element in the

array ($j = n - 1$) where x lies after the else statement, the last if statement does not allow us to check that position because it sets j to 1 starting the loop again from one, incrementing i and therefore this will repeat until $i = n$ and the loop terminates returning -1 because we never checked the last element so we could not find x .

These are proved using the example I gave in question 5. The loop terminates when i is incremented to 6 after the 5th loop → always terminates.

Also when x is the last element in the array, when j goes from $0 \rightarrow 2 \rightarrow 4 \rightarrow 1 \rightarrow 3 \rightarrow 5$ we do not start a loop with $j = 5$ because 5 is $n - 1$ so j becomes 1 and this repeats until $i \geq n$.