

# How Do We Really Find Eigenvalues and Why Should You Care?

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# Outline

## 1 Introduction

- Definition and Existence of Eigenvalues
- Equivalence of Roots and Eigenvalues

## 2 Invariant Subspaces

- The Power Method
- Simultaneous Iteration
- The QR Algorithm

## 3 Large Sparse Eigenvalue Problems

- Arnoldi Process
- Symmetric Lanczos Process
- Orthogonal Polynomials and Numerical Integration

# Definition of Eigenvalues and Eigenvectors

Let  $A \in \mathbb{C}^{n \times n}$ . We say that  $\lambda \in \mathbb{C}$  is an eigenvalue and  $x \in \mathbb{C}^n \setminus \{0\}$  is a corresponding eigenvector, if

$$Ax = \lambda x.$$

# Existence of Eigenvalues

- $\lambda \in \mathbb{C}$  is an eigenvalue of  $A \in \mathbb{C}^{n \times n}$  if and only if  $\lambda I - A$  is singular.
- if and only if  $\det(\lambda I - A) = 0$ .
- if and only if  $\lambda$  is a root of the characteristic polynomial  $p(\lambda) = \det(\lambda I - A)$
- $p(\lambda)$  is a polynomial of degree  $n$ . By the fundamental theorem of algebra  $p(\lambda)$  has  $n$  roots, counting multiplicities.

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# Example

Let  $A = \begin{bmatrix} 1 & 2 \\ 2 & 1 \end{bmatrix}$ .

- The characteristic polynomial is  $p(\lambda) = \lambda^2 - 2\lambda - 3$ .
- The spectrum of  $A$  is  $\sigma(A) = \{3, -1\}$



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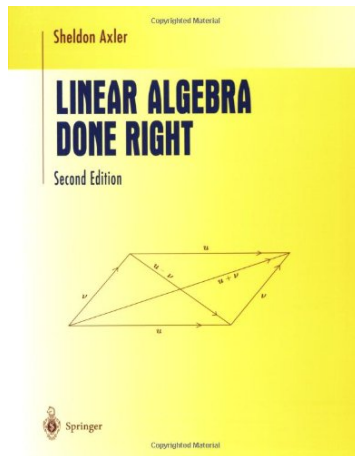
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# Linear Algebra Done Right



# No Determinants

- Linear algebra can be done better without determinants (Sheldon Axler, Down with Determinants!)
- I find it hard to conceive of a situation in which the numerical value of a determinant is needed (Henry Thacher, SIAM News, September 1988).

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# Existence of Eigenvalues

## Proof.

Let  $x \in \mathbb{C}^n \setminus \{0\}$  and consider the Krylov sequence

$$x, Ax, A^2x, \dots$$

There exists a minimal integer  $1 \leq k < n$  such that  $\{x, Ax, \dots, A^k x\}$  is a set of linearly dependent vectors. There exists constants  $a_0, a_1, \dots, a_k$  such that

$$a_0 x + a_1 Ax + \dots + a_k A^k x = 0.$$

Define  $p(z) = a_0 + a_1 z + \dots + a_k z^k = c(z - r_1) \dots (z - r_m)$ , then

$$c(A - r_1 I) \dots (A - r_m I)x = 0.$$



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# Companion Matrix

Let  $p(\lambda) = \lambda^n + a_{n-1}\lambda^{n-1} + \cdots + a_1\lambda + a_0$  be a monic scalar polynomial. We define the companion matrix of  $p$  to be the  $n \times n$  matrix

$$A_p = \begin{bmatrix} -a_{n-1} & \cdots & -a_1 & -a_0 \\ 1 & & & \\ & \ddots & & \\ & & 1 & \end{bmatrix}.$$



## Eigenvalues of Companion Matrix

## Theorem

The eigenvalues of  $A_p$  are the roots of the polynomial  $p(\lambda)$ .

Proof.

The equation  $A_p v = \lambda v$  is equivalent to

$$\begin{bmatrix} -a_{n-1}v_1 - \cdots - a_1v_{n-1} - a_0v_n \\ v_1 \\ \vdots \\ v_{n-1} \end{bmatrix} = \begin{bmatrix} \lambda v_1 \\ \lambda v_2 \\ \vdots \\ \lambda v_n \end{bmatrix}.$$

Therefore,  $v_1 = \lambda^{n-1} v_n, v_2 = \lambda^{n-2} v_n, \dots, v_{n-1} = \lambda v_n$ .

# Summary

- The eigenvalues of a matrix  $A$  are the roots of the characteristic polynomial.
- The roots of a polynomial are the eigenvalues of a companion matrix.
- We cannot expect to solve the eigenvalue problem in a finite number of steps.

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# Definition of dominant eigenvalue

We say that the eigenvalue  $\lambda$  is dominant, if there exists a positive number  $r < |\lambda|$  such that

$$\sigma(A) \setminus \{\lambda\} \subseteq \{z \in \mathbb{C} : |z| \leq r\}.$$

- If  $\lambda$  is dominant, then every eigenvector  $v$  associated with  $\lambda$  is called a dominant eigenvector.

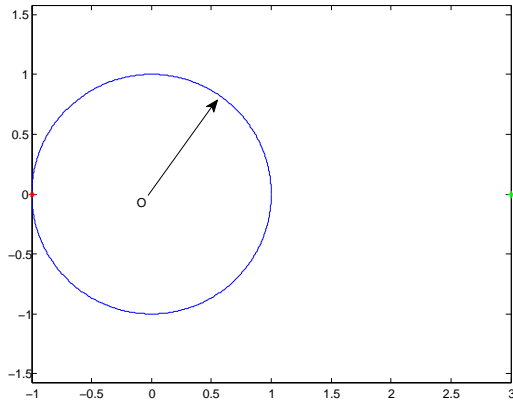
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# The Power Method

Given a nonzero vector  $x \in \mathbb{C}^n$ , the power method forms the Krylov sequence of vectors

$$x, Ax, A^2x, \dots$$

- If  $x$  is not an unlucky choice and if  $A$  has a dominant eigenvector, then the Krylov sequence will converge to a dominant eigenvector of  $A$ .

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Let  $A = \begin{bmatrix} 1 & 2 \\ 2 & 1 \end{bmatrix}$  and  $x_0 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$ .

- Then we update via

$$x_{n+1} = \frac{Ax_n}{\|Ax_n\|}$$

where  $\|\cdot\|$  is the vector 2 norm.

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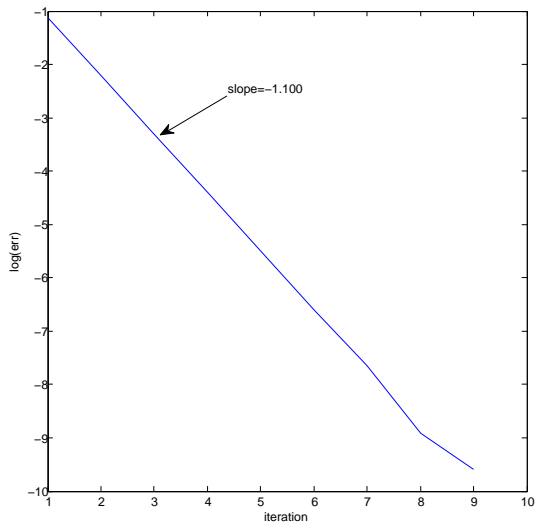
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# Example

$n$	$x_n$	$n$	$x_n$
1	$\begin{bmatrix} 0.4472 \\ 0.8944 \end{bmatrix}$	6	$\begin{bmatrix} 0.7081 \\ 0.7061 \end{bmatrix}$
2	$\begin{bmatrix} 0.7809 \\ 0.6247 \end{bmatrix}$	7	$\begin{bmatrix} 0.7068 \\ 0.7074 \end{bmatrix}$
3	$\begin{bmatrix} 0.6805 \\ 0.7328 \end{bmatrix}$	8	$\begin{bmatrix} 0.7072 \\ 0.7070 \end{bmatrix}$
4	$\begin{bmatrix} 0.7158 \\ 0.6983 \end{bmatrix}$	9	$\begin{bmatrix} 0.7071 \\ 0.7071 \end{bmatrix}$
5	$\begin{bmatrix} 0.7042 \\ 0.7100 \end{bmatrix}$	10	$\begin{bmatrix} 0.7071 \\ 0.7071 \end{bmatrix}$

# Example



# Summary

- The power method generally converges to a dominant eigenvector of  $A$ .
- The rate of convergence is linear and has contraction factor  $\frac{r}{|\lambda|}$ .

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# Subspace Iteration

- When computing eigenvectors, our real object of interest are eigenspaces.
- If the eigenspace is one-dimensional, then a single eigenvector  $v$  is a basis for the space.
- We can view the power method as

$$S, AS, A^2S, \dots$$

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# Definition of Invariant Subspace

Let  $A \in \mathbb{F}^{n \times n}$ . A subspace  $S$  of  $\mathbb{F}^n$  is invariant under  $A$  if  $Ax \in S$  whenever  $x \in S$ . That is,

$$AS \subseteq S.$$

# Eigenvectors form Invariant Subspaces

Let  $v_1, \dots, v_k$  be eigenvectors of  $A$  corresponding to the eigenvalues  $\lambda_1, \dots, \lambda_k$ .

- For  $x \in \text{span}\{v_1, \dots, v_k\}$  we have  $x = c_1 v_1 + \dots + c_k v_k$ .
- $Ax = c_1 Av_1 + \dots + c_k Av_k = c_1 \lambda_1 v_1 + \dots + c_k \lambda_k v_k \in \text{span}\{v_1, \dots, v_k\}$ .
- If  $A$  is semisimple ( $A$  has  $n$  linearly independent eigenvectors), then every invariant subspace under  $A$  has this form.

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# Example

Let  $A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 2 & 1 \\ 0 & 0 & 2 \end{bmatrix}$  and  $S = \text{span} \left\{ \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \right\}$ .

- Note that  $S$  is not spanned by eigenvectors of  $A$ .

- For  $x \in S$ , we have  $x = \begin{bmatrix} 0 \\ a \\ b \end{bmatrix}$ .

- $Ax = \begin{bmatrix} 0 \\ 2a+b \\ 2b \end{bmatrix} \in S$ .

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# Definition of Dominant Invariant Subspace

A  $k$ -dimensional invariant subspace  $S$  is dominant if the eigenvalues  $\lambda_1, \dots, \lambda_k$  of the restricted operator  $A|_S$  dominate the spectrum of  $A$ .

- That is, there exists an  $r$  such that  $r < |\lambda_j|$  for  $j = 1, \dots, k$  and the rest of the spectrum lies within the disc  $\{z \in \mathbb{C} : |z| \leq r\}$ .

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# Subspace Iteration

Let  $S$  be a  $k$ -dimensional space, and not an unlucky choice. Then the sequence

$$S, AS, A^2S, \dots$$

will converge to a  $k$ -dimensional dominant invariant subspace under  $A$ , if  $A$  has one.



# Deflating the Problem

## Theorem

Let  $S = \text{span}\{x_1, \dots, x_k\}$  be invariant under  $A \in \mathbb{F}^{n \times n}$  and

$$\{x_1, \dots, x_k, x_{k+1}, \dots, x_n\}$$

be a basis for  $\mathbb{F}^n$ . Define  $X = [x_1, \dots, x_n]$  and  $B = X^{-1}AX$ , then

$$B = \begin{bmatrix} B_{11} & B_{12} \\ 0 & B_{22} \end{bmatrix}.$$

# Deflating the Problem

Proof.

The equation  $B = X^{-1}AX$  is equivalent to  $AX = XB$ . The  $j^{\text{th}}$  column of this equation is equivalent to

$$Ax_j = \sum_{i=1}^n x_i b_{ij}.$$

If  $S$  is invariant under  $A$ , then  $Ax_j \in \text{span}\{x_1, \dots, x_k\}$  for  $j = 1, \dots, k$ . Therefore,

$$b_{ij} = 0$$

for  $1 \leq j \leq k$  and  $k+1 \leq i \leq n$ . □

# Simultaneous Iteration

For  $i = 0, 1, 2, \dots$ , let  $S^{(i)} = A^i S$ .

- If  $\{q_1^{(i)}, \dots, q_k^{(i)}\}$  is a basis for  $S^{(i)}$ , then  $\{Aq_1^{(i)}, \dots, Aq_k^{(i)}\}$  is a basis for  $S^{(i+1)}$ .
- In practice we specify that  $\{q_1^{(i)}, \dots, q_k^{(i)}\}$  be orthonormal.
- We can obtain an orthonormal basis  $\{q_1^{(i+1)}, \dots, q_k^{(i+1)}\}$  for  $S^{(i+1)}$  by applying the Gram-Schmidt process to  $\{Aq_1^{(i)}, \dots, Aq_k^{(i)}\}$ .

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# Important Benefits

- The Gram-Schmidt process preserves all lower dimensional subspaces

$$\text{span} \left\{ q_1^{(i+1)}, \dots, q_j^{(i+1)} \right\} = \text{span} \left\{ Aq_1^{(i)}, \dots, Aq_j^{(i)} \right\}$$

for  $j = 1, \dots, k$ .

- We are simultaneously performing subspace iteration on subspaces of dimension  $1, \dots, k$ .
- We are always looking for the opportunity to break off a smaller problem of size  $j \times j$ , for  $j = 1, \dots, k$ .

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# Change of Coordinate System

Consider a step of simultaneous iteration with the standard basis vectors  $e_1, e_2, \dots, e_n$ .

- We get  $Ae_1, Ae_2, \dots, Ae_n$ , and then orthonormalize to get  $q_1, q_2, \dots, q_n$ .
- We perform the similarity transformation  $\hat{A} = Q^*AQ$ , where  $Q = [q_1, q_2, \dots, q_n]$ .
- The vectors  $q_1, q_2, \dots, q_n$  in the old coordinate system are  $e_1, e_2, \dots, e_n$  in the new coordinate system.

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- The vectors  $q_1, q_2, \dots, q_n$  in the old coordinate system are  $e_1, e_2, \dots, e_n$  in the new coordinate system.

# Change of Coordinate System

Consider a step of simultaneous iteration with the standard basis vectors  $e_1, e_2, \dots, e_n$ .

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# QR Algorithm

Simultaneous iteration with a change of coordinate system

$$A_i = Q_{i+1} R_{i+1}, \quad A_{i+1} = Q_{i+1}^* A_i Q_{i+1},$$

for  $i = 1, 2, \dots$



# Example

$i$	$A$		$i$	$A$
0	$\begin{bmatrix} 4 & 2 \\ 3 & 3 \end{bmatrix}$			
1	$\begin{bmatrix} 6.0400 & 0.2800 \\ -0.7200 & 0.9600 \end{bmatrix}$		5	$\begin{bmatrix} 6.0001 & 0.9995 \\ -0.0005 & 0.9999 \end{bmatrix}$
2	$\begin{bmatrix} 6.0205 & -0.8832 \\ 0.1168 & 0.9795 \end{bmatrix}$		6	$\begin{bmatrix} 6.0000 & -0.9999 \\ 0.0001 & 1.0000 \end{bmatrix}$
3	$\begin{bmatrix} 6.0038 & 0.9807 \\ -0.0193 & 0.9962 \end{bmatrix}$		7	$\begin{bmatrix} 6.0000 & 1.0000 \\ -0.0000 & 1.0000 \end{bmatrix}$
4	$\begin{bmatrix} 6.0006 & -0.9968 \\ 0.0032 & 0.9994 \end{bmatrix}$			

# Francis's Algorithm

- J.G. Francis. The QR transformation, part 1. Computer J., 4:265-272, 1961.
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  - Definition and Existence of Eigenvalues
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  - Arnoldi Process
  - Symmetric Lanczos Process
  - Orthogonal Polynomials and Numerical Integration

# The Problem of Sparsity

- A sparse matrix is one in which the vast majority of their entries are zero.
- We may be limited by storing constraints.
- Similarity transformations are a bad idea.

# The Problem of Sparsity

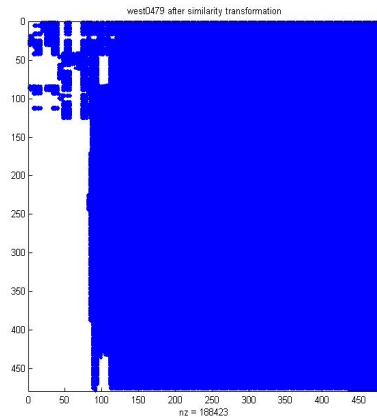
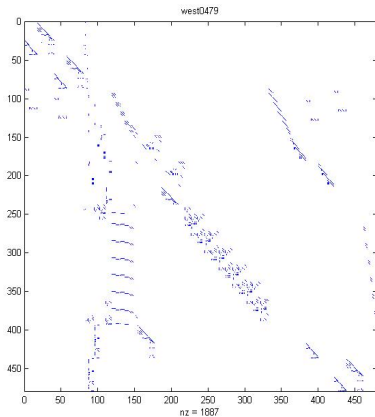
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# Example



# The Idea behind Arnoldi's Process

Given  $x \in \mathbb{C}^n \setminus \{0\}$ , the power method forms the Krylov sequence of vectors

$$x, Ax, A^2x, \dots$$

- At the  $k^{\text{th}}$  iteration we only have  $A^k x$ .
- After  $k$  steps in the Arnoldi process we have

$$x, Ax, \dots, A^k x.$$

- These vectors form a basis for the Krylov subspace

$$K_{k+1}(A, x) = \text{span} \{x, Ax, \dots, A^k x\}.$$

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# We want Orthonormal vectors

On the first step of the Arnoldi process we take  $q_1 = \frac{x}{\|x\|}$ .

- On the second step we take

$$\hat{q}_2 = Aq_1 - \langle Aq_1, q_1 \rangle q_1, \quad q_2 = \frac{\hat{q}_2}{\|\hat{q}_2\|}.$$

- On subsequent steps we take

$$\hat{q}_{k+1} = Aq_k - \sum_{j=1}^k q_j h_{jk}, \quad q_{k+1} = \frac{\hat{q}_{k+1}}{h_{k+1,k}},$$

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# Matrix Representation of Arnoldi Process

We can write the previous steps as

$$Aq_k = \sum_{j=1}^{k+1} q_j h_{jk}, \quad k = 1, 2, 3, \dots$$

- We can write this as

$$A[q_1 \cdots q_m] = [q_1 \cdots q_m q_{m+1}] H_{m+1,m}$$

- Where

$$H_{m+1,m} = \begin{bmatrix} h_{11} & h_{12} & \cdots & h_{1m} \\ h_{21} & h_{22} & \cdots & h_{2m} \\ & h_{32} & & \vdots \\ & & \ddots & h_{mm} \\ & & & h_{m+1,m} \end{bmatrix}.$$

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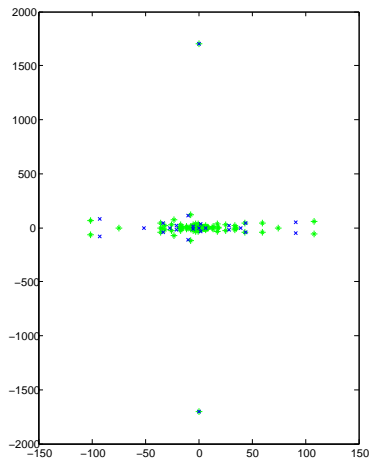
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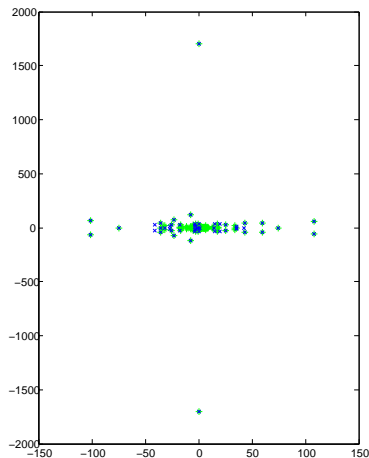
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# Degree 25 Eigenvalue Approximations



# Degree 50 Eigenvalue Approximations



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  - **Symmetric Lanczos Process**
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# Arnoldi Simplified

Let  $A \in \mathbb{R}^{n \times n}$  be symmetric and consider the matrix representation of  $A$  with respect to  $(q_j)_{j=1}^m$ .

$$H_{m+1,m} = \begin{bmatrix} h_{11} & h_{12} & \cdots & h_{1m} \\ h_{21} & h_{22} & \cdots & h_{2m} \\ & h_{32} & & \vdots \\ & & \ddots & h_{mm} \\ & & & h_{m+1,m} \end{bmatrix}.$$



# Arnoldi Simplified

We can write  $H$  as follows

$$H_{m+1,m} = \begin{bmatrix} \alpha_1 & \beta_2 & & \\ \beta_2 & \alpha_2 & \ddots & \\ & \beta_3 & \ddots & \beta_m \\ & & \ddots & \alpha_m \\ & & & \beta_{m+1} \end{bmatrix},$$

where  $\alpha_k = \langle Aq_k, q_k \rangle$  and  $\beta_k = \langle Aq_k, q_{k-1} \rangle$ .

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# A functional space example

Let  $T : \mathcal{H} \rightarrow \mathcal{H}$  be the linear operator  $Tf(x) = xf(x)$ .

- $\mathcal{H} = L^2(-1,1)$  with inner product

$$\langle f, g \rangle = \int_{-1}^1 f(x)g(x)dx,$$

where  $f$  and  $g$  are real valued Lebesgue integrable functions over the interval  $[-1, 1]$ .

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# Space of polynomials

Consider the Krylov sequence starting from  $1 \in \mathcal{H}$

$$1, x, x^2, \dots$$

- The Symmetric Lanczos process will provide an orthonormal basis for this space!
- Starting from  $q_1 \in \mathcal{H}$  such that  $\|q_1\| = 1$ , we have

$$\beta_{k+1}q_{k+1} = (x - \alpha_k)q_k - \beta_k q_{k-1}.$$

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# Legendre Polynomials

Let  $p_1(x) = \frac{1}{\sqrt{2}}$ .

- $\beta_2 p_2(x) = (x - \alpha_1) p_1(x),$
- $\frac{1}{\sqrt{3}} p_2(x) = \frac{1}{\sqrt{2}} x \Rightarrow p_2(x) = \sqrt{\frac{3}{2}} x.$
- $\beta_3 p_3(x) = (x - \alpha_2) p_2(x) - \beta_2 p_1(x),$
- $\frac{2}{\sqrt{15}} p_3(x) = \sqrt{\frac{3}{2}} (x^2 - 1) \Rightarrow p_3(x) = \sqrt{\frac{5}{2}} \left( \frac{3}{2} x^2 - \frac{1}{2} \right).$



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# Legendre Polynomials

We are building the following infinite matrix

$$H = \begin{bmatrix} 0 & \frac{1}{\sqrt{3}} & & & \\ \frac{1}{\sqrt{3}} & 0 & \frac{2}{\sqrt{15}} & & \\ & \frac{2}{\sqrt{15}} & 0 & \frac{3}{\sqrt{35}} & \\ & & \frac{3}{\sqrt{35}} & 0 & \ddots \\ & & & \ddots & \ddots \end{bmatrix}.$$

# Numerical Integration

Suppose we want to compute

$$I = \int_{-1}^1 x \sin(x) dx.$$

- We can approximate this integral by Gauss Quadrature

$$Q_m = 2 \sum_{i=1}^m f(x_i) w_i$$

where  $f(x) = x \sin(x)$ ,  $x_i$  are interpolation points, and  $w_i$  are weights.

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- The correct interpolation points and weights are  $x_i = d_{ii}$  and  $w_i = u_{1i}^2$ .
- $Q_4 = 2(f(d_{11})u_{11}^2 + f(d_{22})u_{12}^2 + f(d_{33})u_{13}^2 + f(d_{44})u_{14}^2) = 0.60234$ .



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