

# **The Conclave**

## Programming 2

Final report for group project

Sapienza, Università di Roma  
Applied Computer Science and Artificial Intelligence  
Academic Year 2024/2025

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# Abstract

This paper describes the design and implementation of a desktop application that simulates the election of a pope inside a virtual Sistine Chapel. The simulation, written in Java 17, models each cardinal as an independent thread that roams a two-dimensional board, exchanges opinions, and eventually casts a vote. A dedicated *Conclave* thread orchestrates multiple scrutinies until a candidate reaches a two-thirds quorum, after which the graphical interface proclaims the *Habemus Papam*. We provide a class-level overview, detail the core algorithms—movement, persuasion, voting—and discuss performance characteristics and avenues for further work.

## Table of contents

Abstract	2
Table of contents	2
1. Introduction	3
2. Architectural Overview	4
3. Core Algorithms	6
4. Key Classes in Detail	9
Screenshots	12

# 1. Introduction

The papal conclave is a centuries-old procedure characterised by secrecy, debate and successive ballots. Modelling such a social process in software offers a sandbox for experimenting with collective decision making, graph diffusion and multithreaded coordination. The project analysed in this paper was produced as an academic exercise and is composed of roughly 1 500 lines of Java code.

A front-end built with *Swing* presents the Latin mass, a start menu, live vote counts and finally the elected pope; the back-end runs the stochastic simulation.

The goal of this report is:

1. Present the overall architecture and execution flow.
2. Detail the algorithms that drive movement, persuasion and voting.
3. Illustrate key implementation choices, highlighting strengths and areas for improvement.

## 2. Architectural Overview

The code is organised as follows

Layer	Package(s)	Principal classes	Responsibility
Simulation core	<code>project.conclave</code>	<i>Conclave, Cardinal, Board, Square, Opinion</i>	Domain logic, multithreading
Utilities	<code>project.util</code>	<i>Heap, Reader, Triplet, ConclaveSetupException</i>	Generic data structures & helpers
Presentation	<code>project.gui</code>	<i>Window + components (StartButton, Title, LeaderboardLabel, PopeLabel, ScrollingText)</i>	Swing UI

The entry point is *Main* in the root package, while a CSV file (passed at runtime) contains the roster.

### 2.1 Data model

#### *Cardinal*

Attributes capture **identity** (name, surname, id), **influence** (an integer weight), **position** (inner class *Position* storing x, y) and **state** (heap of opinions, availability flags, counters).

*Opinion* is a simple pair (id, value) representing the support a cardinal assigns to a candidate.

*Heap* implements a fixed-size **max-heap** of *Opinion* objects plus an auxiliary *positions* array that offers  $O(1)$  access to arbitrary nodes—crucial for quickly updating scores during debates.

*Triplet* records (name, surname, votes) and is comparable for leaderboard sorting.

The conclave unfolds on a *Board*—a square lattice of `size × size` *Squares*, each holding a list of resident cardinals. No walls exist; movement wraps inside bounds.

## 2.2 Thread Topology

```
Main Thread
├─ UI event-dispatch thread (Swing)
└─ Conclave (extends Thread)
    ├─ Cardinal[0] (Runnable)
    ├─ Cardinal[1]
    └─ ... Cardinal[n-1]
```

The *Conclave* thread spawns one worker per cardinal at each scrutiny and blocks until all have finished debating and voting. Synchronisation relies on `Thread.interrupt()` and polling plus explicit `wait()/notify()` on shared locks (`Main.data`, individual *Cardinal* monitors).

## 3. Core Algorithms

### 3.1 Movement in the Sistine Grid

Each cardinal starts at a random square (`Position.setRandom()`) and, while the *Conclave* marks the *debate* phase, executes `moveRandomly()`:

1. Build `availableDirections` based on current coordinates (edges forbid crossing the perimeter).
2. Choose one direction uniformly at random.
3. Atomically remove `this` from the old square, update `(x,y)`, and add to the new square.

The algorithm's critical sections are protected by synchronising on the target *Square* object, preventing inconsistent occupancy lists.

### 3.2 Informal Debates and Opinion Dynamics

When two cardinals share a square, a handshake determines who becomes *caller* vs *listener*. Both parties execute `exchangeInformation(Cardinal caller)`, whose heart is a

**probabilistic influence contest:**

```
threshold = myInfluence / (myInfluence +
callerInfluence) × 100
if random(0..99) <= threshold
    caller wins → listener updates
else    listener wins → caller updates
```

After a duel, the loser adopts (partially) the top opinion from the winner's heap, weighted by the latter's influence. Each participant carries an array `encounteredOpinions[cardinalId]` to avoid counting the same persuasion twice.

### 3.3 Voting and Scrutiny Cycle

At the end of the 3-second debate window:

1. *Conclave* interrupts every *Cardinal* thread.
2. Each cardinal's heap head (most valued candidate) is read and tallied into `votes[int id]`.
3. Results are pushed into `Main.data` as *Triplets*, the UI refreshes the leaderboard, and the majority check runs.

A cardinal becomes **pope** when their vote count satisfies

`votes[i] ≥ 2 × floor(cardinals.size / 3)`

—an implementation of the canonical  $\frac{2}{3}$  super-majority. If unmet, all opinion heaps persist, new *Cardinal* threads are spawned, and the process reiterates.

### 3.4 Quorum Calculation

Canon law prescribes that a new pontiff must be elected by at least two thirds of the cardinal electors present. The implementation translates this rule into integer arithmetic:

```
int target = (cardinals.size() / 3) * 2;    // 2/3
without floating point
if (votes[pope] >= target) {
    /* success */
}
```

Because `int` division truncates toward zero, the expression is correct for every cardinal count that is a multiple of three and errs on the safe side (requires one extra vote) otherwise.

## 3.5 Termination and Resource Cleanup

Once the condition is met the `Conclave` thread breaks out of its loop, interrupts any lingering worker threads (belt&braces) and publishes the pope's full name via the static field `Main.pope`. The UI reads that field in `Window.displayPope()` and replaces the whole content pane with a bold "Habemus Papam" title and the elected name. In production one would also:

- shut down the Swing timer responsible for scrolling text;
- close the CSV scanner;
- interrupt the event-dispatch thread on a clean exit.

## 3.6 Complexity Analysis

*Movement and opinion updates are  $O(1)$  amortised; vote counting is  $O(n)$  per scrutiny.* The heap operations remain logarithmic, but with only one entry per cardinal, they behave effectively as constant time. The whole simulation therefore scales linearly with the number of participants.



## 4. Key Classes in Detail

### 4.1 Conclave

*Conclave* holds static references to the *Board* and `ArrayList<Cardinal>`. Its `run()` method:

```
initializeBoard(csvPath, boardSize);
while (true) {
    spawnCardinals();
    Arrays.fill(votes, 0);
    sleep(3000);
    interruptWorkers();
    waitWorkers();
    collectVotes();
    if (hasTwoThirdsMajority()) break;
}
announcePope();
```

Robustness: the constructor validates the CSV path and throws *ConclaveSetupException* on error.

### 4.2 Cardinal

Implements *Runnable*; the loop terminates on `Thread.interrupt()`.  
Concurrency primitives:

- A cardinal waits (`wait(random(250—750))`) when it has spoken to all neighbours to mitigate busy-waiting.
- Method-level synchronized blocks ensure that opinion exchanges are transactional.

The twin counters `cardinalsToTalkTo` and `cardinalsToListenTo.size()` implement a simple rendezvous protocol: a cardinal finishes only after delivering and receiving every scheduled message.

## 4.3 Heap

Functions: `add(id, value)`, `remove(id)`, `getTop()`. The heap stores one node per candidate and updates values in logarithmic time via `heapifyUp/heapifyDown`. The auxiliary `positions[ ]` solves *decrease-key* without a hash map.

## 4.4 Utility & GUI Components

*Reader* parses semicolon-separated values and skips the CSV header. *Window* encapsulates a `JFrame` and exposes high-level views: mass scroll, main menu, votes leaderboard and pope screen. Child components adopt monospace fonts for an *ad hoc* aesthetic. Synchronisation with the simulation uses the shared `Main.data` and `Main.menu` locks.

## 4.5 GUI Workflow and Styling

*Window* anchors the Swing hierarchy and opts for absolute positioning rather than layout managers—a valid choice given the fixed dimensions of 600×400px. Each view is composed like a slide:

- **Menu View** – shows a mono-spaced title and a `StartButton`.

When the user clicks, the listener performs `synchronized (Main.menu) { menu.notify(); }`, unblocking the main thread.

- **Mass View** – creates a `ScrollingText` pane with Latin verses of the *Exit omnes* mass, centers it horizontally, and starts a `javax.swing.Timer` that raises its `y` coordinate by 2px every 30ms, achieving a slow upward crawl.

- **Leaderboard View** – accepts `ArrayList<Triplet<String,Integer>>` and, after sorting, instantiates one `LeaderboardLabel` per entry. Labels maintain consistent alignment by padding vote counts to three digits (`String.format("%03d", votes)`).
- **Pope View** – emphasises the finale with a 32pt bold font. Finally, *Habemus Papam!*

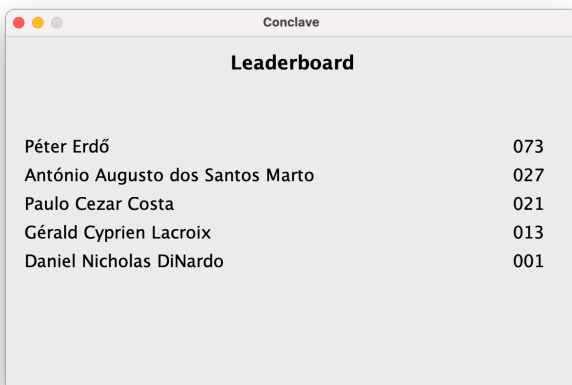
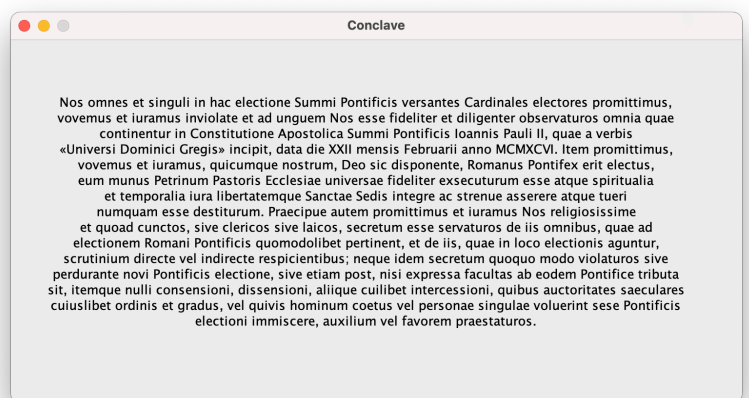
All custom components share the `DejaVu Sans Mono` typeface to impose a retro terminal vibe. Colours are OS defaults, achieving sufficient contrast without explicit RGB codes.

## Screenshots



<< The Conclave is set to start, it just needs the Camerlengo (you) to start it

>> The Exit Omnes Mass is conducted by the Camerlengo.  
No worries,  
we've got you covered!



<< Cardinals are influencing each other, the leaderboard updates!

>> White smoke comes from the Sistine Chapel! We got the new Pope! *Habemus Papam!*

