1. The Tree Monad

The Tree type defined in the lecture can actually be interpreted as a monad—we can think of a value of type Tree a as a computation delivering a tree of results, such as one might use in a search algorithm, for example. Write the Monad instance for this type.

You can check your definition by copying the following into the file Tree.hs, and adding your definition to it.

As well as a definition of the Tree type, this code defines how to generate random trees (arbitrary)¹, and how to simplify a tree in a failed test (shrink). It also defines property instances for testing the monad laws. These property instances use another module, MonadLaws, that you will need to copy into the file MonadLaws.hs.

¹ Leaves are given a higher weight than branches to ensure that random generation terminates.

```
prop_Assoc m (Fun _ f) (Fun _ g) =
  ((m >>= f) >>= g)
  ==
  (m >>= \x -> f x >>= g)
```

Add your definition of the Monad instance for Tree, load Tree. hs into ghci, make sure it type-checks, and then test the laws.

I'm not familiar with Haskell—how do I do that?

Install the Haskell Platform. Once you've done so, on Windows it's convenient to use WinGHCi, which you will find in the Haskell Platform folder in the Start menu. On Linux/Mac just type ghci in the shell.

To load a file into ghci, type either

```
:1 <filename without the .hs>
```

or, in WinGHCi, use the File->Load... menu item. Imported modules will also be loaded, and any compile-time errors will be reported.

Once you have loaded your code, you can evaluate expressions at the ghci prompt. To test the monad laws, just type

```
quickCheck prop_TreeLeftUnit
```

and so on.

2. The State Monad

Copy the following code into State.hs.

```
module State where
import Control.Monad
import Test.QuickCheck
newtype State s a = MkState {unState :: s -> (a,s)}
instance Monad (State s) where
  return x = MkState (\s -> (x,s))
 MkState f \gg g = MkState (\s -> let (a,s') = f s in
                                    unState (g a) s')
get :: State s s
get = undefined
put :: s -> State s ()
put s = undefined
(===) ::
  Eq a => State Integer a -> State Integer a -> Integer -> Bool
(f === g) s = unState f s == unState g s
prop get get =
  do x <- get
```

```
y <- get
  return (x,y)
===
do x <- get
  return (x,x)</pre>
```

It defines the state monad presented in the lecture, but using a newtype definition rather than a type synonym, so that it is accepted by GHC. Instead of a tick operation to increment the state, it defines a get operation to read the state, and a put operation to write the state. But these definitions are incomplete—finish them.

It is a little awkward to generate random values of the State type (we would need to generate values containing a QuickCheck Fun, and convert them into the State type as needed), so instead we use QuickCheck to formulate and test some properties of put and get. The (===) operator checks that two values of type State Integer a deliver the same value and final state, when invoked in the same state. It can be used to define properties such as prop_get_get, which says that two consecutive gets are equivalent to a single one. Check that this property passes (using quickCheck prop_get_get), and add and test properties relating puts and gets in either order, and relating two puts to one put.

3. Readers and Writers

The state monad lets us thread a state through a program, passing it in to each computation, and returning a new state after each one. Sometimes, though, we need only to do one of these things.

The Reader s monad lets us pass a state into a computation, but does not let us modify it:

```
newtype Reader s a = MkReader {unReader :: s -> a}
ask :: Reader s s
```

The Writer's monad lets us return state *from* a computation, but does not let us read it. Since we may write the state many times in a computation, we collect a *list* of state values:

```
data Writer s a = MkWriter [s] a
tell :: s -> Writer s ()
```

Write Monad instances for Reader s and Writer s, and definitions of ask and tell, and make sure they type-check.

The version of the Writer monad in the Haskell libraries is a little more general than this: is allows the state to be of any *monoid* type, rather than just lists; we need to be able to combine states (a binary operator) in (>>=), we need an "empty" state to use in return, and the monad laws demand that these form a monoid.

4. The List Monad

Haskell's list type is also a monad—think of it as representing computations that return zero, one, or more answers. Try to define a suitable Monad instance for lists²:

```
instance Monad [] where
  return x = ...
  xs >>= f = ...
```

Sadly, you cannot compile this definition yourself, because a monad instance for lists is already defined. Instead, you will need to define an isomorphic list type MyList, and write your instance for that. Copy the following code into MyList.hs, add a suitable Monad instance, and test the monad laws.

```
{-# LANGUAGE GeneralizedNewtypeDeriving #-}
module MyList where

import Control.Monad
import Test.QuickCheck
import MonadLaws

newtype MyList a = MkList {unList :: [a]}
   deriving (Eq, Show, Arbitrary)

prop_MyListLeftUnit = prop_LeftUnit :: PropLeftUnit MyList
prop_MyListRightUnit = prop_RightUnit :: PropRightUnit MyList
prop_MyListAssoc = prop_Assoc :: PropAssoc MyList
mymap f (MkList xs) = MkList (map f xs)

flat :: MyList (MyList a) -> MyList a
flat (MkList xs) = MkList (concat (map unList xs))
```

This code defines the MyList type, inheriting test data generation from the underlying list type via "deriving Arbitrary"—the pragma at the top enables this useful extension to Haskell 98. It defines monomorphic properties for testing the monad laws, and then a pair of useful auxiliary functions for writing a Monad instance. Add your own Monad instance, and test the stated properties, now. Try varying your definitions, and see whether the monad laws then fail.

Once you have experimented a little, read the next section of the exercise.

It turns out that there *is* another plausible way to write a Monad instance for MyList. If you found the expected definition, then the flat function will have played a central role in the definition of (>>=). This function just concatenates a list of lists to obtain a list of all the elements. But if we think of a list of lists as a matrix, then there is another natural way to extract a list of elements—we can just take the diagonal of the matrix! Add the following code to your file:

```
diag :: MyList (MyList a) -> MyList a
diag (MkList (MkList (x:xs):xss)) =
  mycons x (diag (MkList (map mydrop xss)))
```

 $^{^2}$ The parameterised list type in Haskell is written <code>[]</code>, so <code>[]</code> a is the same type as <code>[a]</code> .

```
diag _ = MkList []
mycons x (MkList xs) = MkList (x:xs)
mydrop (MkList (x:xs)) = MkList xs
mydrop (MkList []) = MkList []
```

Now, what if we replace the concatenation in (>>=) with the diagonal—do we get another, different monad?

It turns out that we have to modify the return function too; the right definition of return in this case should construct an *infinite* list.

```
return x = MkList (repeat x)
```

This in turn means that we may need to *compare* infinite lists as we test the monad laws. Of course, this will not terminate—so let us *replace* the definition of equality for MyList with an approximate equality that only compares the first 100 elements. This is good enough for testing the monad laws; we expect that any failures will be observable in far fewer than 100 elements. Add the following instance definition to your code, and remove the "Eq," from the deriving clause.

```
instance Eq a => Eq (MyList a) where
  MkList xs == MkList ys = take 100 xs == take 100 ys
```

Complete a Monad instance using diag. Is the result really a monad?