

# Project Template

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December 7, 2016

## 1 Introduction

The goal of our project is to create a routing system for bicycle trips that takes into account more than just distance. Current map applications like Google often route cyclists through busy streets and intersections, resulting in unnecessarily dangerous trip and sometimes accidents. The algorithm we chose to use for this problem is A\* search with a variety of cost functions and heuristics incorporating distance, safety of a given route, and changes in elevation (also important for biking but not information Google maps currently considers).

Additionally, we considered the problem of finding a central meeting place for two or more cyclists starting from different positions. To solve this optimization problem, we used k-beam search and simulated annealing.

The graph search problem is directly analogous to techniques and ideas from the class: using intersections as nodes in the graph and road segments connecting intersections as paths, we are able to construct the city maps as directed graphs. From there, we can use data generated from a variety of sources to come up with approximate costs associated with paths and related heuristics at nodes.

The optimization techniques we use are also direct applications of algorithms from the optimization portion of the course.

Problem-specific adaptations were largely limited to the collection and interpretation of relevant data and the development of a cost function and relevant heuristics.

## 2 Background and Related Work

While our problem was fairly straightforward, we did do some research to get ideas for potential cost attributes and data sources. Our initial formulation of the problem, for example, did not include elevation differencing, but reading previous studies led us to believe that this was an important criterion. [1].

## 3 Problem Specification

The problem that we are solving is of the general class of graph search problems. Given a strongly connected graph, we seek to minimize some cost function  $cost(\sum_{i=0}^n path_i)$  given that  $path_0 \in Connections_{startingnode}$

$path_i, path_{i+1} \in Connections_{node_j}$  for some  $j$ ,  
 $path_i, path_{i-1} \in Connections_{node_k}$  for some  $k$ , and  
 $path_n \in Connections_{targetnode}$

Specifically for our problem, the cost function is:  
 $cost(path_i) = \alpha_i * length_i + \beta_i * abs(\Delta elevation_i)$

In the optimization portion of our project, we seek to minimize the total cost to all parties:  
 $min(\sum_j^n cost_j)$  where  $cost_j$  is the cost to participant  $j$ .

## 4 Approach

Data Structures: Our primary data structures are an intersection graph, which is stored as a Python dictionary, and a connection dictionary, which is also stored as a Python dictionary. The intersection graph maps nodes (intersections) by id to a list of paths (road segments) from that node. The connection dictionary maps a connection (road segment) to its source node, sink node, and various cost parameters (in our case distance, number of bicycle crashes on that road segment, and change in elevation over that road segment).

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### Algorithm 1 A-Star Search

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```

function A-STAR-SEARCH(graph, startnode, targetnode)
  node  $\leftarrow$  a node with STATE = startnode
  PATH-COST  $\leftarrow$  heuristic(startnode, targetnode)
  frontier  $\leftarrow$  a priority queue ordered by PATH-COST + heuristic(node, targetnode) with node
  as the only element
  explored  $\leftarrow$  an empty set
  loop
    if EMPTY?(frontier) then
      return failure
    end if
    node  $\leftarrow$  POP(frontier) /*chooses the lowest cost+heuristic node in frontier */
    if node == targetnode then
      return SOLUTION(node)
    end if
    add node.STATE to explored
    for each path in PATHS(node) do
      child  $\leftarrow$  child-node(node, path)
      if child.STATE is not in explored or frontier then
        frontier  $\leftarrow$  insert(child, frontier)
      else if child.STATE is in frontier with higher PATH-COST + heuristic then
        replace that frontier node with child
      end if
    end for
  end loop
end function

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**Algorithm 2** Simulated Annealing Meeting Spot

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```
function SIMULATED ANNEALING MEETING SPOT(graph, startingpts, cost, heuristic)  
  if length(startingpts) < 2 then  
    return error  
  end if  
  current  $\leftarrow$  mean(startingpts).CLOSEST-NODE  
  temperature  $\leftarrow e^{10}$   
   $\gamma \leftarrow .5$  /*schedule to manage temperature */  
  while temperature >  $e^{-2}$  do  
    temperature  $\leftarrow$  temperature *  $\gamma$   
    next  $\leftarrow$  a randomly selected child of current  
    current.VALUE  $\leftarrow \sum_{pt \in \text{startingpts}} \mathbf{cost}(pt, \text{centroid})$   
    next.VALUE  $\leftarrow \sum_{pt \in \text{startingpts}} \mathbf{cost}(pt, \text{next})$   
     $\Delta E \leftarrow \text{next.VALUE} - \text{current.VALUE}$   
    if  $\Delta E > 0$  then  
      current  $\leftarrow$  next  
    else  
      current  $\leftarrow$  next with probability  $e^{\Delta E / \text{temperature}}$   
    end if  
  end while  
end function
```

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**Algorithm 3** K-Beam Search Meeting Spot

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```
function K-BEAM SEARCH MEETING SPOT( $k, graph, startingpts, cost, heuristic$ )  
  if length( $startingpts$ ) < 2 then  
    return error  
  end if  
  
   $\{candidatenodes\} \leftarrow node \forall node \in graph \text{ s.t.}$   
   $node.x \geq \min(startingpts.x) \& node.x \leq \max(startingpts.x) \&$   
   $node.y \geq \min(startingpts.y) \& node.y \leq \max(startingpts.y)$   
  
   $point_i \leftarrow$  a randomly selected  $node \in \{candidatenodes\} \forall i \leq k$   
   $best.VALUE \leftarrow \min_{i \leq k} (\sum_{pt \in startingpts} cost(pt, point_i))$   
  while True do  
     $\{nextcosts\} \leftarrow \sum_{pt \in startingpts} cost(pt, child_i) \forall i \leq k, child_i \in PATHS(point_i).endnode$   
     $point_i \leftarrow$  i-th least  $node \in \{nextcosts\} \forall i \leq k$   
     $next.VALUE \leftarrow \sum_{pt \in startingpts} cost(pt, point_1)$   
    if  $next.VALUE < best.VALUE$  then  
       $best \leftarrow next$   
    else  
      break  
    end if  
    return  $best$   
  end while  
end function
```

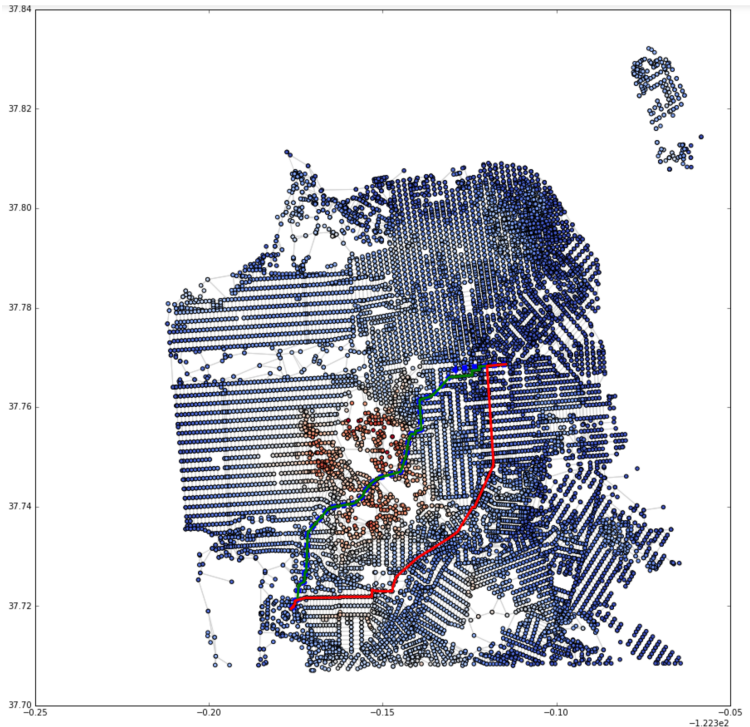
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## 5 Experiments

Analysis, evaluation, and critique of the algorithm and your implementation. Include a description of the testing data you used and a discussion of examples that illustrate major features of your system. Testing is a critical part of system construction, and the scope of your testing will be an important component in our evaluation. Discuss what you learned from the implementation.

- **Analysis:** To test our algorithms, we built graphs on GIS data from both Cambridge, MA and San Francisco, CA (supplemented with elevation and crash data collected from other sources as necessary). We ran both the single-path and meeting point algorithms for a variety of cost functions and start and end points and were pleased to see that the paths generated by our algorithms when taking into account safety and elevation factors do indeed seem to avoid busy intersections and steep hills in a reasonable way.

Figure 1: Paths between nodes in San Francisco. Blue path optimizes for distance only, green path avoids previous bicycle accidents, and red path minimizes altitude changes. Overlaid with an altitude plot of San Francisco, we can clearly see the red path avoiding a hill.



- Evaluation:
- Critique:

Score
Approach 1
Approach 2

Table 1: Description of the results.

### 5.1 Results

For algorithm-comparison projects: a section reporting empirical comparison results preferably presented graphically.

## 6 Discussion

Summary of approach and results. Major takeaways? Things you could improve in future work?

## A System Description

Appendix 1 A clear description of how to use your system and how to generate the output you discussed in the write-up. *The teaching staff must be able to run your system.*

## B Group Makeup

1. Nick Hoernle
  - (a) Creation of graph dictionary structure and A\*search algorithm
  - (b) Simulated annealing
2. Nikhila Ravi
  - (a) K-Beam Search
  - (b) Visualization and analysis of results of graph search algorithms
3. Anna Sophie Hilgard
  - (a) Construction of Datasets
  - (b) Research and Implementation of more complicated cost functions and heuristics

## References

- [1] Jan Hrnčíř, Palov Zilecky, Qing Song, and Michal Jakob. Practical multicriteria urban bicycle routing. *IEEE Transactions on Intelligent Transportation Systems*, PP(99):1–12, 2016.