# Assignment Unit 5

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1. Suppose the function $G : N \to N$ is defined recursively by setting $G(0) = 0$ and $G(n) = 1 + G(\lfloor n/2 \rfloor)$ for every positive integer n. Prove that for all $n \in \mathbb{N}$, we have $G(n) = 1 + \lfloor \log_2 n \rfloor$.

Hint: If you are struggling with the inductive step, try splitting the problem into even and odd cases.

2. Indicate, for each pair of expressions $(A, B)$ in the table below, whether $A$ is $O$, $\Omega$, or $\Theta$, of $B$. Assume all logarithms have base 2.

a. $A = \log(n!), B = \log(n^n)$

b. $A = 400n^{-3} + \sqrt{n} + \pi, B = 1 + \frac{\sin n}{\pi}$

c. $A = \log(n)^{\log(n)}, B = 4^{\log(n)}$

d. $A = 2^n, B = n!$

e. $A = n^{-5099}, B = \frac{1}{5099} \log n$

3. Write a program that uses recursion to calculate the nth Fibonacci number. Now write one that uses dynamic programming to do the same thing. Record your time and space efficiency for each approach using Big-O notation. (Note: for this problem a Python solution is preferred, although readable pseudocode is also acceptable)

4. Consider a weighted graph $G$ with $V$ vertices and $E$ edges, and suppose all edge weights are nonnegative. Prove using induction that the following algorithm can find the shortest path from a vertex "source" to any other vertex of the graph.

```
function Dijkstra(Graph, source):
    create vertex priority queue Q
    dist[source] ← 0
    add source to Q
    for each vertex v in Graph.Vertices:
        dist[v] ← INFINITY
        prev[v] ← UNDEFINED
        add v to Q
```

```
dist[source] ← 0

while Q is not empty:
    u ← vertex in Q with minimum dist[u]
    remove u from Q

    for each neighbor v of u still in Q:
        alt ← dist[u] + Graph.Edges(u, v)
        if alt < dist[v]:
            dist[v] ← alt
            prev[v] ← u

return dist, prev
```

Also use Big-O notation to find the time and space complexity of the algorithm in terms of $V$ and $E$. A short explanation will suffice.

5. Find a recurrence relation for the auxilliary space $S(n)$ used by mergesort on an array $n$. Use the recurrence to prove that the space complexity of merge sort is $\theta(n)$

Hint: After looking at the code, you might be tempted to think $S(n)$ satisfies the same recurrence as $T(n)$ from the tutorial. However, there is a slight difference that changes the space complexity.