

CS201- Lecture 7

IA32 Data Access and Operations

Part II

RAOUL RIVAS

PORTLAND STATE UNIVERSITY

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Announcements

Address Computation Examples

<code>%rdx</code>	<code>0xf000</code>
<code>%rcx</code>	<code>0x0100</code>

- `movq 8(%rdx), %rax`
- `movq (%rdx,%rcx), %rax`
- `movq (%rdx,%rcx,4), %rax`

Expression	Address Computation	Address
<code>0x8(%rdx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%rdx,%rcx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%rdx,%rcx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%rdx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

Compute Effective Address

- **`leaq Src, Dst`**
 - *Src* is address mode expression
 - Set *Dst* to address denoted by expression
- **Uses**
 - Computing addresses without a memory reference
 - E.g., translation of `p = &x[i];`
 - Computing arithmetic expressions of the form $x + k*y$
 - $k = 1, 2, 4, \text{ or } 8$
- **Example**

```
long m12(long x)
{
    return x*12;
}
```

Converted to ASM by compiler:

```
leaq (%rdi,%rdi,2), %rax # t <- x+x*2
salq $2, %rax           # return t<<2
```

Basic Arithmetic Operations

- Two Operand Instructions:

<i>Format</i>	<i>Computation</i>	
<code>addq</code>	<code>Src, Dest</code>	<code>Dest = Dest + Src</code>
<code>subq</code>	<code>Src, Dest</code>	<code>Dest = Dest - Src</code>
<code>imulq</code>	<code>Src, Dest</code>	<code>Dest = Dest * Src</code>
<code>salq</code>	<code>Src, Dest</code>	<code>Dest = Dest << Src</code> <i>Also called shlq</i>
<code>sarq</code>	<code>Src, Dest</code>	<code>Dest = Dest >> Src</code> <i>Arithmetic</i>
<code>shrq</code>	<code>Src, Dest</code>	<code>Dest = Dest >> Src</code> <i>Logical</i>
<code>xorq</code>	<code>Src, Dest</code>	<code>Dest = Dest ^ Src</code>
<code>andq</code>	<code>Src, Dest</code>	<code>Dest = Dest & Src</code>
<code>orq</code>	<code>Src, Dest</code>	<code>Dest = Dest Src</code>

- Watch out for argument order!
- Notice the accumulator format
 - What's the C equivalent of this accumulator operators?

Unary Operators

- One Operand Instructions

<code>incq</code>	<i>Dest</i>	$Dest = Dest + 1$
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<code>decq</code>	<i>Dest</i>	$Dest = Dest - 1$
-------------------	-------------	-------------------

<code>negq</code>	<i>Dest</i>	$Dest = - Dest$
-------------------	-------------	-----------------

<code>notq</code>	<i>Dest</i>	$Dest = \sim Dest$
-------------------	-------------	--------------------

Translating from C

```
long easyOp(long x, long y)
{
    long rval;

    x = x + 2;
    y = y * x;
    rval = y;
    return rval;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y

Translating from C

```
long easyOp(long x, long y)
{
    long rval;

    x = x + 2;
    y = y * x;
    rval = y;
    return rval;
}
```

```
easyOp:
    addq    $2, %rdi
    imulq   %rdi, %rsi
    movq    %rsi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	rval

If you find accumulator
arithmetic tricky try writing
some C programs in this way

Accumulators

- Can we do the same in Accumulator and in Non-Accumulator architectures?
 - Yes. They are **Turing Equivalent**
 - Computer P can simulate computer Q and computer Q can simulate computer P
 - Any program that runs on P can be rewritten to run in Q



IBM 701
38 bit single
accumulator
register

Arithmetic Expression Example

```
long arith
(long x, long y, long z)
{
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
    movq    %rdi, %rax
    addq    %rsi, %rax
    addq    %rdx, %rax
    movq    %rdi, %rcx
    addq    $4, %rcx
    imulq   $48, %rsi
    addq    %rsi, %rcx
    imulq   %rcx, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z

Arithmetic Expression Optimization

```
long arith
(long x, long y, long z)
{
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
    leaq    (%rdi,%rsi), %rax
    addq    %rdx, %rax
    leaq    (%rsi,%rsi,2), %rdx
    salq    $4, %rdx
    leaq    4(%rdi,%rdx), %rcx
    imulq   %rcx, %rax
    ret
```

Interesting Instructions

- **leaq**: address computation
- **salq**: shift
- **imulq**: multiplication
 - But, only used once

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z

Arithmetic Expression Optimization

```
long arith
(long x, long y, long z)
{
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

arith:

```
leaq    (%rdi,%rsi), %rax    # t1
addq    %rdx, %rax          # t2
leaq    (%rsi,%rsi,2), %rdx
salq    $4, %rdx            # t4
leaq    4(%rdi,%rdx), %rcx   # t5
imulq   %rcx, %rax          # rval
ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	t1, t2, rval
%rdx	t4
%rcx	t5

Arithmetic tricks

- Shift to replace multiplication
 - `salq $4, %rdx`
- Use of LEA to replace complex arithmetic
 - `leaq (%rsi,%rsi,2), %rdx`
- XOR to replace initialization to zero
 - `xorq %rax, %rax`

What's Next?

- At this point we can write some simple Assembly functions
 - Read and Write from Memory
 - Variables, Pointers, Arrays, Structures
 - Logic and Arithmetic Operations
- What's missing?
 - Comparisons
 - Control Statements
 - If/Else, While, Do/While
 - Function Calls

Summary

- LEA instruction is used to compute addresses without a memory reference
- LEA is also used to optimize certain patterns of arithmetic computations
- Arithmetic Operations use an accumulator based notation