

Demostración de correctitud del programa 1

Demostración de correctitud

- $Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c)$
- $P_c \longrightarrow wp(\text{ciclo}, Q_c)$
- $Q_c \longrightarrow wp(\text{codigo post ciclo}, Post)$

Por monotonia sabemos que $Pre \longrightarrow wp(\text{programa completo}, Post)$

Datos

- $Pre \equiv |s| > 0$
- $Post \equiv 0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j])$
- $P_c \equiv |s| > 0 \wedge i = 0 \wedge r = 0$
- $Q_c \equiv 0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j])$
- $I \equiv (0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j])$
- $B \equiv i < |s|$
- $f_v \equiv |s| - i$
- $S1 \equiv \text{if } s[i] \geq s[r] \text{ then } r := i \text{ else skip fi}$
- $S2 \equiv i := i + 1$
- $ciclo \equiv S1; S2$

$Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c)$

$$\begin{aligned} wp(i := 0, wp(r := 0, |s| > 0 \wedge i = 0 \wedge r = 0)) &\equiv |s| > 0 \\ Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c) &\equiv \\ |s| > 0 \longrightarrow |s| > 0 &\equiv \text{true} \end{aligned}$$

$Q_c \longrightarrow wp(\text{codigo post ciclo}, Post)$

$$\begin{aligned} Q_c \longrightarrow wp(\text{codigo post ciclo}, Post) &\equiv \\ 0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j]) &\longrightarrow \\ 0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j]) &\equiv \text{true} \end{aligned}$$

$P_c \longrightarrow wp(ciclo, Q_c)$

Teorema del invariante

- $P_c \longrightarrow I$
- $(I \wedge \neg B) \longrightarrow Q_c$
- $\{I \wedge B\} ciclo \{I\}$
- $\{I \wedge B \wedge f_v = v_0\} ciclo \{f_v < v_0\}$
- $(I \wedge f_v \leq 0) \longrightarrow \neg B$

$P_c \longrightarrow I$

- $P_c \equiv |s| > 0 \wedge i = 0 \wedge r = 0$
 - $I \equiv (0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j])$
- $$|s| > 0 \wedge i = 0 \wedge r = 0 \longrightarrow$$
- $$(0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j]) \equiv \text{true}$$

$$(I \wedge \neg B) \longrightarrow Q_c$$

- $Q_c \equiv 0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j])$
- $I \equiv (0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j])$
- $B \equiv i < |s|$
- $\neg B \equiv i \geq |s|$

$$(0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j]) \wedge i \geq |s| \equiv$$

$$i = |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j]) \longrightarrow$$

$$0 \leq r < |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < |s| \longrightarrow_L s[r] > s[j]) \equiv \text{true}$$

$$\{I \wedge B\} \text{ciclo}\{I\}$$

- $I \equiv (0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j])$
- $B \equiv i < |s|$
- $S1 \equiv \text{if } s[i] \geq s[r] \text{ then } r := i \text{ else skip fi}$
- $S2 \equiv i := i + 1$
- $\text{ciclo} \equiv S1; S2$

$$\text{wp}(\text{ciclo}, I) \equiv \text{wp}(S1; S2, I) \stackrel{Ax3}{\equiv} \text{wp}(S1, \text{wp}(S2, I))$$

$$\text{wp}(S2, I) \equiv$$

$$\text{wp}(i := i + 1, (0 \leq i \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j])) \equiv$$

$$(0 \leq i + 1 \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i + 1 \longrightarrow_L s[r] > s[j]) \equiv$$

$$\text{wp}(S1, \text{wp}(S2, I)) \equiv$$

$$\text{wp}(\text{if } s[i] \geq s[r] \text{ then } r := i \text{ else skip fi}, (0 \leq i + 1 \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i + 1 \longrightarrow_L s[r] > s[j])) \stackrel{Ax4}{\equiv}$$

$$\{(0 \leq i < |s| \wedge 0 \leq r < |s|) \wedge_L ($$

$$(s[i] \geq s[r]) \wedge 0 \leq i + 1 \leq |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < i \longrightarrow_L s[i] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) i \leq j < i + 1 \longrightarrow_L s[i] > s[j])) \vee$$

$$(s[i] < s[r]) \wedge (0 \leq i + 1 \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i + 1 \longrightarrow_L s[r] > s[j]))\}$$

$$I \wedge B \equiv (0 \leq i < |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i \longrightarrow_L s[r] > s[j]) \longrightarrow$$

$$(0 \leq i < |s| \wedge 0 \leq r < |s|) \wedge_L ($$

$$(s[i] \geq s[r]) \wedge 0 \leq i + 1 \leq |s| \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < i \longrightarrow_L s[i] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) i \leq j < i + 1 \longrightarrow_L s[i] > s[j])) \vee$$

$$(s[i] < s[r]) \wedge (0 \leq i + 1 \leq |s| \wedge 0 \leq r < |s|) \wedge_L ((\forall j : \mathbb{Z}) 0 \leq j < r \longrightarrow_L s[r] \geq s[j]) \wedge ((\forall j : \mathbb{Z}) r \leq j < i + 1 \longrightarrow_L s[r] > s[j])) \equiv \text{true}$$

$$\{I \wedge B \wedge f_v = v_0\} \text{ciclo}\{f_v < v_0\}$$

$$|s| - i = v_0 \longrightarrow |s| - i - 1 < v_0 \equiv \text{true}$$

$$(I \wedge f_v \leq 0) \longrightarrow \neg B$$

$$|s| - i \leq 0 \longrightarrow i \geq |s|$$