

# Ejercicio 14

## Correctitud

- $Pre \longrightarrow wp(\text{codigo previo a los ciclos}, P_c)$
- $P_c \longrightarrow wp(\text{ciclo1; ciclo2}, Q_c)$
- $Q_c \longrightarrow wp(\text{codigo posterior a los ciclos}, Post)$

Por monotonía sabemos que  $Pre \longrightarrow wp(\text{programa completo}, Post)$

## Demostración

### Datos

- $Pre \equiv |r| = |a| + |b| \wedge r = R_0$
- $Post \equiv |r| = |R_0| \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |a| \longrightarrow_L r[j] = a[j]) \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |b| \longrightarrow_L r[j + |a|] = b[j])$
- $P_c \equiv |r| = |a| + |b| \wedge r = R_0 \wedge i := 0$
- $Q_c \equiv |r| = |R_0| \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |a| \longrightarrow_L r[j] = a[j]) \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |b| \longrightarrow_L r[j + |a|] = b[j])$

$Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c)$

$$\begin{aligned} wp(i := 0, |r| = |a| + |b| \wedge r = R_0 \wedge i := 0) &\equiv \\ |r| = |a| + |b| \wedge r = R_0 & \\ Pre \longrightarrow \text{true} &\equiv \\ |r| = |a| + |b| \wedge r = R_0 \longrightarrow |r| = |a| + |b| \wedge r = R_0 &\equiv \text{true} \end{aligned}$$

$Q_c \longrightarrow wp(\text{codigo posterior al ciclo}, Post)$

$$Q_c \longrightarrow wp(\text{codigo posterior al ciclo}, Post) \equiv$$

$$\begin{aligned} |r| = |R_0| \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |a| \longrightarrow_L r[j] = a[j]) \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |b| \longrightarrow_L r[j + |a|] = b[j]) &\longrightarrow \\ |r| = |R_0| \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |a| \longrightarrow_L r[j] = a[j]) \wedge ((\forall j : \mathbb{Z}) 0 \leq j < |b| \longrightarrow_L r[j + |a|] = b[j]) &\equiv \text{true} \end{aligned}$$

$P_c \longrightarrow wp(\text{ciclo1; ciclo2}, Q_c)$

ciclo1.pdf ciclo2.pdf