

Ejercicio 13

Correctitud

- $Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c)$
- $P_c \longrightarrow wp(\text{ciclo}, Q_c)$
- $Q_c \longrightarrow wp(\text{codigo posterior al ciclo}, Post)$

Por monotonía sabemos que $Pre \longrightarrow wp(\text{programa completo}, Post)$

Teorema del invariante

- $P_c \longrightarrow I$
- $(I \wedge \neg B) \longrightarrow Q_c$
- $\{I \wedge B\}$ ciclo $\{I\}$
- $\{I \wedge B \wedge (v_0 = f_v)\}$ ciclo $\{f_v < v_0\}$
- $(I \wedge f_v \leq 0) \longrightarrow \neg B$

Demostración

Datos

- $Pre \equiv \text{true}$
- $Post \equiv r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)]$
- $P_c \equiv i = 0 \wedge j = |s| - 1 \wedge r = \text{true}$
- $Q_c \equiv r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)]$
- $I \equiv (0 \leq i \leq |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])$
- $B \equiv i < |s|$
- $S1 \equiv \text{if } s[i] \neq s[j] \text{ then } r := \text{false} \text{ else skip fi}$
- $S2 \equiv i := i + 1$
- $S3 \equiv j := j - 1$
- $ciclo \equiv S1; S2; S3;$
- $f_v \equiv |s| - i$

$Pre \longrightarrow wp(\text{codigo previo al ciclo}, P_c)$

$wp(i := 0, wp(j := |s| - 1, wp(r := \text{true}, i = 0 \wedge j = |s| - 1 \wedge r = \text{true}))) \equiv$
 $0 = 0 \wedge |s| - 1 = |s| - 1 \wedge \text{true} = \text{true} \equiv \text{true}$
 $Pre \longrightarrow \text{true} \equiv \text{true} \longrightarrow \text{true}$

$Q_c \longrightarrow wp(\text{codigo posterior al ciclo}, Post)$

$wp(\text{codigo posterior al ciclo}, Post) \equiv$

$r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)]$

$Q_c \longrightarrow wp(\text{codigo posterior al ciclo}, Post) \equiv$

$r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)] \longrightarrow$

$r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)] \equiv \text{true}$

$$P_c \longrightarrow wp(ciclo, Q_c)$$

$$\blacksquare P_c \longrightarrow I$$

$$i = 0 \wedge j = |s| - 1 \wedge r = \text{true} \longrightarrow (0 \leq i \leq |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \equiv \text{true}$$

$$\blacksquare (I \wedge \neg B) \longrightarrow Q_c$$

$$(0 \leq i \leq |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \wedge i \geq |s| \longrightarrow r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)] \equiv$$

$$(i = |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \longrightarrow r = \text{true} \leftrightarrow (\forall i : \mathbb{Z}) 0 \leq i < |s| \longrightarrow_L s[i] = s[|s| - (i + 1)] \equiv \text{true}$$

$$\blacksquare \{I \wedge B\} \text{ ciclo } \{I\}$$

$$wp(S1; S2; S3, I) \stackrel{Ax3}{\equiv} wp(S1, wp(S2, wp(S3, I)))$$

$$\bullet wp(S3, I) \equiv wp(j := j - 1, I)$$

$$\{(0 \leq i \leq |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])\}$$

$$\bullet wp(S2, wp(S3, I)) \equiv wp(i := i + 1, wp(S3, I)) \equiv$$

$$\{(0 \leq i + 1 \leq |s| \wedge j = |s| - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])\}$$

$$\bullet wp(S1, wp(S2, wp(S3, I))) \equiv$$

$$wp(\text{if } s[i] \neq s[j] \text{ then } r := \text{false} \text{ else } skip \text{ fi}, (0 \leq i + 1 \leq |s| \wedge j = |s| - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])) \stackrel{Ax4}{\equiv}$$

$$(0 \leq i < |s| \wedge 0 \leq j < |s|) \wedge_L ($$

$$(s[i] \neq s[j] \wedge ((0 \leq i + 1 \leq |s| \wedge j = |s| - i)) \wedge_L (\text{false} = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]))) \vee$$

$$(s[i] = s[j] \wedge (0 \leq i + 1 \leq |s| \wedge j = |s| - i)) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])))$$

$$(0 \leq i \leq |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \wedge s < |s| \equiv$$

$$(0 \leq i < |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \longrightarrow$$

$$(0 \leq i < |s| \wedge 0 \leq j < |s|) \wedge_L ($$

$$(s[i] \neq s[j] \wedge (0 \leq i + 1 \leq |s| \wedge j = |s| - i)) \wedge_L (\text{false} = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])) \vee$$

$$(s[i] = s[j] \wedge (0 \leq i + 1 \leq |s| \wedge j = |s| - i)) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]))) \equiv$$

$$(0 \leq i < |s| \wedge j = |s| - 1 - i) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)]) \longrightarrow$$

$$(0 \leq i < |s| \wedge 0 \leq j < |s|) \wedge_L ($$

$$(s[i] = s[j] \wedge (0 \leq i + 1 \leq |s| \wedge j = |s| - i)) \wedge_L (r = \text{true} \leftrightarrow (\forall k : \mathbb{Z}) 0 \leq k < i \longrightarrow_L s[k] = s[|s| - (k + 1)])) \equiv \text{true}$$

$$\blacksquare \{I \wedge B \wedge (v_0 = f_v)\} \text{ ciclo } \{f_v < v_0\}$$

$$wp(j : j - 1, j < v_0) \equiv j - 1 < v_0$$

$$wp(i : i + 1, j - 1 < v_0) \equiv j - 1 < v_0$$

$$wp(\text{if } s[i] \neq s[j] \text{ then } r := \text{false} \text{ else } skip \text{ fi}, j - 1 < v_0) \equiv 0 \leq j | s| \wedge j - 1 < v_0$$

$$I \wedge B \wedge v_0 = j \longrightarrow 0 \leq j | s| \wedge j - 1 < v_0 \equiv \text{true}$$

$$\blacksquare (I \wedge f_v \leq 0) \longrightarrow \neg B$$

$$I \wedge j \leq 0 \longrightarrow i < |s| \equiv \text{true}$$