

- use case
 - query result instead of [[SQL]] table
 - aggregation value instead of scalar value
- subqueries can be stored as variable
 - WITH VariableName AS (SELECT ...)

Modularization with WITH C AS (SELECT ...)

- common use cases:

- **Subqueries w/ IN**

- Check containment of values in result set of sub query

```
SELECT Product, Quantity, Price
FROM Sales
WHERE Product NOT IN(
  SELECT Product FROM Sales
  GROUP BY Product
  HAVING sum(Quantity*Price)>1e6)
```

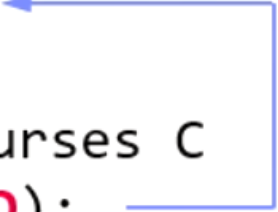
- **Other subqueries**

- **EXISTS**: existential quantifier $\exists x$ for correlated subqueries
- **ALL**: comparison (w/ universal quantifier $\forall x$)
- **SOME/ANY**: comparison (w/ existential quantifier $\exists x$)

Correlated and Uncorrelated Subqueries

- correlated if queries depend on each other
 - subquery executed for each tuple of outer query
 - * nested for loop
 - inefficient

```
SELECT P.Fname, P.Lname
FROM Professors P,
WHERE NOT EXISTS(
  SELECT * FROM Courses C
  WHERE C.PID=P.PID);
```



- uncorrelated if subquery executed once

```

SELECT P.Fname, P.Lname
FROM Professors P,
WHERE P.PID NOT IN(
    SELECT PID FROM Courses);

```

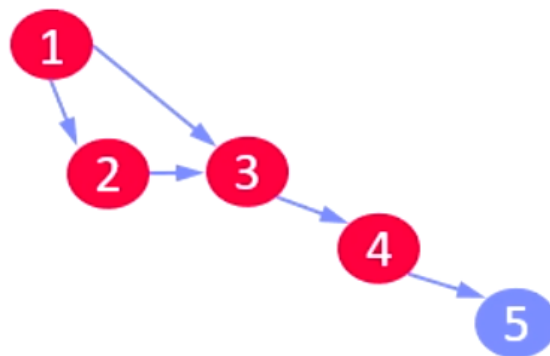
- correlated queries may be unnested (de-correlation)
 - can also be improved by “only” executing subquery for each distinct value

Recursive Queries

- terminates when recursive query returns empty table
 - **Approach**
 - **WITH RECURSIVE** <name> (<arguments>)
 - Compose recursive table from **non-recursive term**, **union all/distinct**, and **recursive term**
 - Terminates when recursive term yields empty result

▪ Example

- Courses(CID, Name),
Precond(pre REF CID, suc REF CID)
- Dependency graph (pre→suc)



```

WITH RECURSIVE rPrereq(p,s) AS(
    (SELECT pre, suc
     FROM Precond WHERE suc=5)
    UNION DISTINCT
    (SELECT B.pre, B.suc
     FROM Precond B, rPrereq R
     WHERE B.suc = R.p)
)
SELECT DISTINCT p FROM rPrereq

```

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[[Correlation]]