

# Performance Evaluation in Networks

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## 0 Introduction

### 0.1 Generalities

Performance of computer and communication systems/networks. Objectives:

- Observation
- Prediction
- Control & Optimization

Background needed:

- Not much, try to give always the specification of the system we will use

Systems can be architecture/hardware, code/software, communication network (included distributed systems) and logistic/industrial processes, etc.

The metrics will be:

- Speed, bandwidth, delay, load, losses.
- Worst case, average, ..

Use of either a mathematical analysis of an abstract model or a simulation (math model, scale model). Comparisons with experiments/measures on a real system (statistics).

Objectives:

- Designing and analysing mathematical models (probability assumptions)
- Improve knowledge in probability/statistical tool
- Improve knowledge about communication networks
- practising simulation/statistic with a simulation/statistic software

Why are there so many probabilities?

- Sometimes intrinsic in the system (i.e. noise in the core of the system)
- The users themselves introduce probabilities (Ethernet protocol and randomize algorithms)
- Probabilistic is a good way to have a good "rough estimation"
- Usually works well in practice

Use of statistics: Do I need to do more experiments? Do I need to run them longer (asymptotical behaviour)?

Classical questions: Overload? Average waiting time?

## Bibliography:

- Performance Evaluation of Computer and Communication systems
- Introduction aux probabilités et à la statistique
- The Art of Computer Systems Performance Analysis
- <http://perfeval.epfl.ch/>
- <http://proweb>.

## 0.2 Refresher about Probability

*Cf first year*

Usual working hypothesis:

- A list  $\Omega$  of the outcomes and a measure of their occurrences via the measures of events.

**Definition 1.**  $\Omega$  is the universe,  $\mathcal{F}$  a set of subsets of  $\Omega$ ,  $\mathbb{P}$  function from  $\Omega$  to  $\mathbb{R}$ .

Vocabulary:

- $\omega \in \Omega$  is an outcome
- $A \in \mathcal{F}$  is an event
- $\omega$  realise  $A$  if  $\omega \in A$

Recalls:

- Conditional probabilities
- Independence of events
- Law of total probabilities
- General/Real Random Variable
- Cumulative distribution function:  $F_x(x) = \mathbb{P}(X \leq x)$  ( $\lim_{x \rightarrow -\infty} F(x) = 0$  and  $\lim_{x \rightarrow +\infty} F(x) = 1$ ,  $F$  non decreasing,  $F$  right continuous)
- Discrete/Continuous r.v.<sup>1</sup>
- Random vectors & Joint distribution
- Expectation
- Composition of discrete/continuous r.v. (if  $g$  is Lebesgue integrable, then  $\mathbb{E}(g(X)) = \int_x g(x)f(x)dx$ )
- Same for vectors
- Moments of a r.v. (generalisation of expectation and standard deviation)

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<sup>1</sup>Random Variable

- Generating functions (Probabilities  $G_X(s) = \mathbb{E}(s^x)$ , moments  $M_X(t) = \mathbb{E}(e^{tX})$ , characteristic  $\Phi_X(t) = \mathbb{E}(e^{itX})$ )
- Markov, Bienaymé-Tchebychev, Jensen (convex function), Hölder, Minkowski
- Chernoff and Hoeffding
- Convergence (in law, in proba, almost sure)
- Central limit theorem
- Law of large numbers

## 1 Random number generation

cf Webpage

## 2 Markov chains

Recalls

## 3 Martingales and stuff

## 4 Continuous time Markov chains

**Definition 2.** A process  $(X_t)_{t \in \mathbb{R}_+}$  over a space of states  $E$  is a homogeneous continuous Markov chain if:

$$\begin{aligned} \forall t_1 < \dots < t_n \in \mathbb{R}_+, \mathbb{P}(X_{t_n} = i_n | X_{t_{n-1}} = i_{n-1}, \dots, X_{t_1} = i_1) &= \mathbb{P}(X_{t_n} = i_n | X_{t_{n-1}} = i_{n-1}) && \text{(Markov)} \\ &= \mathbb{P}(X_{t_n - t_{n-1}} = i_n | X_0 = i_{n-1}) && \text{(Homogeneous)} \end{aligned}$$

**Notation:**  $p_{ij}(t) = \mathbb{P}(X_t = j | X_0 = i) \rightarrow P_t = (p_{ij}(t))_{i,j \in E}$

**Kolmogorov equations:**  $P_{t+s} = P_s P_t, \forall s, t \in \mathbb{R}_+$  (same proof)

**Discrete case:**  $P_n = P^n$  with  $P = P_1 = \mathbb{P}(X_1 = j | X_0 = i)$  ("kernel")

**Continuous case:**  $P_t \xrightarrow{t \rightarrow 0} ?$

**Classical assumption:**  $P_t \xrightarrow{t \rightarrow 0} Id$  i.e.  $\forall i, j, p_{ij}(t) \xrightarrow{t \rightarrow 0} 0$  if  $i \neq j$  and  $p_{ii}(t) \xrightarrow{t \rightarrow 0} 1$   
("standard") ("continuity") ("differentiability at 0")

$$\begin{aligned} \exists Q, \frac{P_t - Id}{t} &\xrightarrow{t \rightarrow 0} Q \\ \text{i.e. } \forall i \neq j, \frac{p_{ij}(t)}{t} &\rightarrow q_{ij} \\ \forall i, \frac{p_{ii}(t) - 1}{t} &\rightarrow q_{ii} \end{aligned}$$

$Q$  = infinitesimal generator = "kernel"

**Remarks about coeff of  $Q$ :**

- $\forall i \neq j, q_{ij} \geq 0$
- 

$$\begin{aligned} \forall t \in \mathbb{R}_+, \sum_{j \in E} p_{ij}(t) &= 1 \\ \Rightarrow \forall t > 0, \frac{\sum_{j \in E} P_{ij}(t) - 1}{t} &= 0 \\ (t \rightarrow 0) \Rightarrow \sum_{j \in E} q_{ij} &= 0 \\ \Rightarrow q_{ii} &\leq 0 \end{aligned}$$

**Vocabulary:**  $q_{ij}$  = transition *rate* from  $i$  to  $j$

**Theorem 1** (Kolmogorov). *With previous assumptions,*

$$\forall t \in \mathbb{R}, \frac{dP_t}{dt} = QP_t$$

*Proof.* Use  $P_{t+s} = P_t P_s$ , more precisely:

$$\begin{aligned} p_{ij}(t+h) - p_{ij}(t) &= \sum_k p_{ik}(h)p_{kj}(t) - p_{ij}(t) \\ &= \sum_{k \neq i} p_{ik}(h)p_{kj}(t) - (p_{ii}(h) - 1)p_{ij}(t) \end{aligned}$$

Divide by  $h$  and make  $h$  tend to 0:

$$\begin{aligned} \lim_{h \rightarrow 0} \frac{p_{ij}(t+h) - p_{ij}(t)}{h} &= \sum_{i \neq k} q_{ik}p_{kj}(t) + q_{ii}p_{ij}(t) \\ &= \text{coeff } (i, j) \text{ of } QP_t \end{aligned}$$

□

**Theorem 2** (Kolmogorov bis).

$$\frac{dP_t}{dt} = P_t Q$$

**Corollary 1.** *Since we know that  $P_0 = I$  the identity matrix:*

$$P_t = e^{tQ} \stackrel{\text{def}}{=} \sum_{n \in \mathbb{N}} \frac{(tQ)^n}{n!}$$

**Remark:** all the work on finite matrices apply to those infinite matrices up to defining/using matrix norms (many norms defined by combinations of sup,  $\sum$ )

**Alternate description:** "jump process". 1 Trajectory of HMC = alternation of *waiting times* and jumps.

**Waiting times:** suppose you are at state  $i$  at time 0.  
Consider

$$\begin{aligned}
W &= \text{the time when you leave } i \\
&= \inf\{t \geq 0 | X_t \neq i\} \\
\mathbb{P}(W > s+t | W > s) &=? \\
\{W > s\} &= \{\forall 0 \leq u \leq s, X_u = i\} \\
&= \mathbb{P}(\underbrace{\forall 0 \leq u \leq s+t, X_u = i}_{s \leq u \leq s+t} | \underbrace{\forall 0 \leq u \leq s, X_u = i}_{X_s = i}) \\
&= \mathbb{P}(W > t) \rightarrow \text{Memoryless law} \\
&= \exp(\underbrace{\lambda}_{?}) \\
\mathbb{P}(W > t) &= e^{-\lambda t} \underset{t \text{ small}}{\simeq} p_{ii}(t) = \mathbb{P}(X_t = i | X_0 = i) \\
&\stackrel{\text{def}}{=} 1 + Q_{ii}t + o(t)
\end{aligned}$$

**Jump probability:** supp. that you are at state  $i$  at time 0.

Consider an interval  $[s, s+t[$  such that  $s \leq W < s+t$ , and suppose that the chain jumps only in this interval,

$$\mathbb{P}(\text{jump to } i | \text{it jumps}) \underset{t \text{ small}}{\simeq} \frac{p_{ij}(t)}{1 - p_{ii}(t)} = \frac{q_{ij}t + o(t)}{-q_{ii}t + o(t)} \xrightarrow{t \rightarrow 0} \frac{q_{ij}}{-q_{ii}}$$

## Evolution/asymptotic behaviour of HMC

**Invariant measure/distribution:**  $\pi \in \mathbb{R}_+^E$  satisfying  $\forall t \geq 0, \pi P_t = \pi$  (and proba dist, i.e.  $\sum_{i \in E} \pi_i = 1$ )

**Theorem 3.**

$$\forall t \geq 0, \pi P_t = \pi \Leftrightarrow \pi Q = 0$$

*Proof.*

$$\begin{aligned}
\pi Q = 0 &\Leftrightarrow \forall n \in \mathbb{N}^*, \pi Q^n = 0 \\
&\Leftrightarrow \forall t \in \mathbb{R}_+, \pi \sum_{n=0}^{+\infty} \frac{t^n Q^n}{n!} = \pi
\end{aligned}$$

□

**Irreducible HMC:** rate transition graph

$$= \begin{cases} \text{vertices} & \text{states } E \\ \text{directed edges } i \rightarrow j & \text{if } q_{ij} > 0 \end{cases}$$

irreducible  $\stackrel{\text{def}}{=}$  graph strongly connected.

**Theorem 4** (Convergence result). *Suppose that an HMC is irreducible and admits an invariant proba distribution  $\pi$ .*

$$\lim_{t \rightarrow \infty} P_{ij}(t) = \pi_j$$

Moreover, let  $f : E \rightarrow \mathbb{R}$  such that  $\sum_{i \in E} \pi_i |f(i)| < \infty$ .

Then

$$\lim_{t \rightarrow +\infty} \frac{1}{t} \int_0^t f(X_s) ds = \sum_{i \in E} \pi_i f(i)$$

## 5 Queues

Queue = buffer + server

### Classical questions:

- Average waiting time for a client
- Proportion of time the server is busy
- Distribution of nb of clients waiting in the queue

**Kendall notation** (for single queue)

$$\underbrace{A}_{\substack{\text{law of interarrival} \\ \text{of clients}}} / \underbrace{B}_{\substack{\text{law of service} \\ \text{time (time to} \\ \text{serve one client)}}} / \underbrace{S}_{\substack{\text{nb of} \\ \text{process}}} (/ \underbrace{N}_{\substack{\text{nb of} \\ \text{buffers}}})$$

All are usually assumed independent.

### Classical laws:

- $M$  (Markov) = exponential law
- $D$  (Deterministic) = constant time
- $E$  (Erlang) = Erlang law
- $G$  (General) = arbitrary law

**Example:** analysis of the  $M(\lambda)/M(\mu)/1$  queue.

$$X_t = \text{size of the queue at time } t$$

### Different behaviour:

- $X_t = 0 \rightarrow$  new client will arrive within time  $\sim \exp(\lambda)$
- $X_t \geq 1 \rightarrow$  one client is being served within time  $\sim \exp(\mu)$  and one new client will arrive within time  $\sim \exp(\lambda)$

### 5.1 M/M/1 queue

$(N_t)_{t \in \mathbb{R}_+}$  is a continuous time HMC.

**Search for invariant proba distribution**  $\pi$   $\pi Q = \pi \Rightarrow \forall n \in \mathbb{N}, \pi_n = \left(\frac{\mu}{\lambda}\right)^n \pi_0$  An invariant proba measure exists iff  $\sum_{n=0}^{+\infty} \pi_n = 1 \Leftrightarrow \frac{\mu}{\lambda} < 1$ . Then  $\pi_n = \left(\frac{\lambda}{\mu}\right)^n \left(1 - \frac{\lambda}{\mu}\right)$  (geometric law).  
under the assumption that the system is at the permanent/stationary state/regime, i.e.  $N_0$  follows the invariant distribution  $\pi$ , let's study some performance parameters.

1. Average nb of packets in the system (we assume that we have the distribution  $\pi$  at time 0)

$$\mathbb{E}_\pi(N_t) = \sum_{n=0}^{+\infty} n \pi_n = \frac{\frac{\mu}{\lambda}}{1 - \frac{\mu}{\lambda}} \leftarrow \text{mean for geometric law}$$

2. Let  $W$  = waiting time before being served

Two cases:

- arriving in an empty queue  $\Leftrightarrow$  being served immediately  
 $\mathbb{P}(W = 0) = \mathbb{P}_\pi(N_t = 0) = \pi_0 = 1 - \frac{\mu}{\lambda}$
- arriving in a non-empty queue: if there is a packet (with probability  $\pi_n$ ) you must wait or sum of  $n$  independent exponential laws  $Exp(\lambda)$   
 $\rightarrow$  density  $f(t)$  for  $W$ :

$$\begin{aligned} f(t) &= \sum_{n=1}^{+\infty} \underbrace{\pi_n}_{\left(\frac{\mu}{\lambda}\right)^n \left(1 - \frac{\mu}{\lambda}\right)} \times \underbrace{\text{density of a sum of independent } Exp(\lambda)}_{e^{-\lambda t} \frac{\lambda^n t^{n-1}}{(n-1)!}} \\ &= e^{-\lambda t} \left(1 - \frac{\mu}{\lambda}\right) \underbrace{\mu \sum_{n=1}^{+\infty} \frac{(\mu t)^{n-1}}{(n-1)!}}_{e^{\mu t}} \\ &= e^{-(\lambda - \mu)t} (\lambda - \mu) \frac{\mu}{\lambda} \end{aligned}$$

We have the law of  $W$ , let's compute  $\mathbb{E}_\pi(W)$

$$\begin{aligned} \mathbb{E}_\pi(W) &= \int_{t=0}^{+\infty} t f(t) dt = \int_0^{+\infty} t e^{-(\lambda - \mu)t} (\lambda - \mu) \frac{\mu}{\lambda} dt \\ &= \frac{\mu}{\lambda} \int_0^{+\infty} t \underbrace{(\lambda - \mu) e^{-(\lambda - \mu)t}}_{\text{density for } Exp(\lambda - \mu)} dt \\ &= \frac{\mu}{\lambda} \frac{1}{\lambda - \mu} \end{aligned}$$

Average time before leaving the system  $\bar{T}$  i.e. spent in the system =  $\mathbb{E}_\pi(W) + \underbrace{\frac{1}{\lambda}}_{\text{average waiting time}} = \frac{\mu}{\lambda} \frac{1}{\lambda - \mu} + \frac{1}{\lambda} = \frac{1}{\lambda - \mu}$

**Remark**

$$\underbrace{\bar{N}}_{\text{average size of the queue}} = \underbrace{\mu}_{\text{input rate}} \underbrace{\bar{T}}_{\text{average time spent in the queue}} \quad (\text{Little's law})$$

## 5.2 M/M/ $\infty$ queue

Behaviour of  $N_t$  = number of packet in the system?

**Search for an invariant distribution  $\pi$**  Substitution:

$$\pi_n = \frac{1}{n!} \left(\frac{\mu}{\lambda}\right)^n \pi_0$$

There is always an invariant distribution because  $\sum_{n=0}^{+\infty} \pi_n = \pi_0 \sum_{n=0}^{+\infty} \frac{1}{n!} \left(\frac{\mu}{\lambda}\right)^n = \pi_0 e^{\frac{\mu}{\lambda}}$



$$\Rightarrow \forall n \in \mathbb{N}, \pi_n = e^{-\frac{\mu}{\lambda}} \frac{1}{n!} \left(\frac{\mu}{\lambda}\right)^n$$

Under the stationary assumption

1.  $\mathbb{E}_\pi(N_t) = \bar{N}$  = mean value for Poisson laws of parameter  $\frac{\lambda}{\mu} = \frac{\lambda}{\mu}$
2.  $\mathbb{E}_\pi(\underbrace{\text{time spent in the system}}_{\text{service time}}) = \text{mean value for } Exp(\lambda) = \frac{1}{\lambda}$ .

**Remark**

$$\bar{N} = \mu \bar{T}$$

## A useful relation: Little's law

In many cases,

average nb of guys in the system = average arrival rate  $\times$  average time spend in the system

$$\bar{N} = \mu \bar{T}$$

## Examples

1. M/M/1 queue in the stationary state
2. M/M/ $\infty$  in the stationary state
3. General deterministic case with finite horizon

**Definition 3** (Pseudo-inverse).

$$f^{(-1)}(n) = \inf\{t \mid f(t) \geq n\}$$

**Definition 4** (Definitions). •  $A(t)$  = number  $n \in \mathbb{N}$  of packets entering between time 0 and time  $t \in \mathbb{R}_+$

- $B(t)$  = number  $n \in \mathbb{N}$  of packets leaving between time 0 and time  $t \in \mathbb{R}_+$
- $N(t) = A(t) - B(t)$  = number of packet in the system at time  $t$
- $T(n) = B^{(-1)}(n) - A^{(-1)}(n)$  = time spent in the system by the  $n$ -th packet

**”Averages:”**

- “finite horizons” = on  $[0, t]$
- $\mu = \frac{A(t)}{t}$  (arrival rate)
- $\bar{T} = \frac{1}{A(t)} \sum_{i=1}^{A(t)} T(i)$  (average time spent in the system)
- $\bar{N} = \frac{1}{t} \int_{x=0}^t N(x) dx$  (average number of packets).

**Theorem 5** (Little's theorem (1961)). If  $A(t) = B(t)$

$$\bar{N} = \mu \bar{T}$$

*Proof.* Graphical proof, interpreting integrals as areas.

$$\underbrace{\frac{1}{t} \int_{x=0}^t N(x) dx}_{\bar{N}} = \underbrace{\frac{A(t)}{t}}_{\mu} \underbrace{\frac{1}{A(t)} \sum_{i=1}^{A(t)} T(i)}_{\bar{T}}$$

□