QUALITY: FOR(;;) {TEST; SPECIFY; CODE}

ILLUSTRATED WITH THE SIMPLE PROBLEM OF TESTING THE EQUIVALENCE OF CIRCULAR LISTS

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 $A(i: 0 \le i < N)$ $B(i: 0 \le i < N)$ A. i = A. (i + N)

```
@Test
public void testSetGetModulo() {
    CL a = new CL(3);
    a.set(0, 7); a.set(4, 8); a.set(2, 9);
    assertEquals(8,a.get(1));
}
```

```
@Invariant("s>=0")
public class CL {
    * s: size of the circular list (CL)
    public final int s;
    * a: array used to model a CL
    private int a[];
   CL(int _size){
       s = _size;
       a = new int[s];
```

```
/**
  * set the value of an element of the CL
  * @param _i index in CL understood modulo the size of the CL
  * @param _x value of the new element
  */
  @Requires("_i >= 0")
  public void set(int _i, int _x){
    a[(_i%s)] = _x;
}
```

```
/**
 * get a value from the CL
 *
 * @param _i index in CL understood modulo the size of the CL
 * @return the value at index _i modulo the size of the CL
 */
 @Requires("_i >= 0")
 public int get(int _i){
    return a[(_i%s)];
}
```

Set of rotations of A:

$$RA.i = A(k: i \le k < i + N)$$

$$RA.i = RA.(i + N)$$

$$\#RA \le N$$

```
@Test
public void testRot() {
    CL a = new CL(3);
    a.set(0, 7); a.set(1, 8); a.set(2, 9);
    int[] expected = {8, 9, 7};
    assertArrayEquals(expected, a.rot(1));
}
```

```
/**
 * rotate a CL by an offset of _i
 * @param _i
 * @return an array that represents the rotation of the CL
@Requires("_i >= 0")
public int[] rot(int _i){
   int[] res = new int[s];
   for(int j=_i, k=0; j < (_i+s); j++, k++){
       res[k] = get(j);
   return res;
```

Postcondition:

 $R: res \equiv (\exists i, j :: RA.i = RB.j)$

```
@Test
public void testRot_eq() {
    CL a = new CL(3);
    a.set(0,7); a.set(1,9); a.set(2,8);
    CL b = new CL(3);
    b.set(0,9); b.set(1,8); b.set(2,7);
    assertTrue(CL.rot_eq(a, 0, b, 2));
}
```

```
/**
  * @param _a CL
  * @param _i offset used to rotate _a
  * @param _b CL
  * @param _j offset used to rotate _b
  * @return Is the _ith-rotation of CL _a equals to the _jth-rotation of CL _b?
  */
@Requires({"_i >= 0", "_j >= 0"})
public static boolean rot_eq(CL _a, int _i, CL _b, int _j){
    return Arrays.equals(_a.rot(_i), _b.rot(_j));
}
```

 $R: res \equiv (\exists i, j :: RA.i = RB.j)$

Naïve solution:
Compare A with each element of RB

```
@Test
public void testCL_eq() {
    CL a = new CL(3);
    a.set(0,7); a.set(1,9); a.set(2,8);
    CL b = new CL(3);
    b.set(0,9); b.set(1,8); b.set(2,7);
    assertTrue(CL.eq(a, b));
}
```

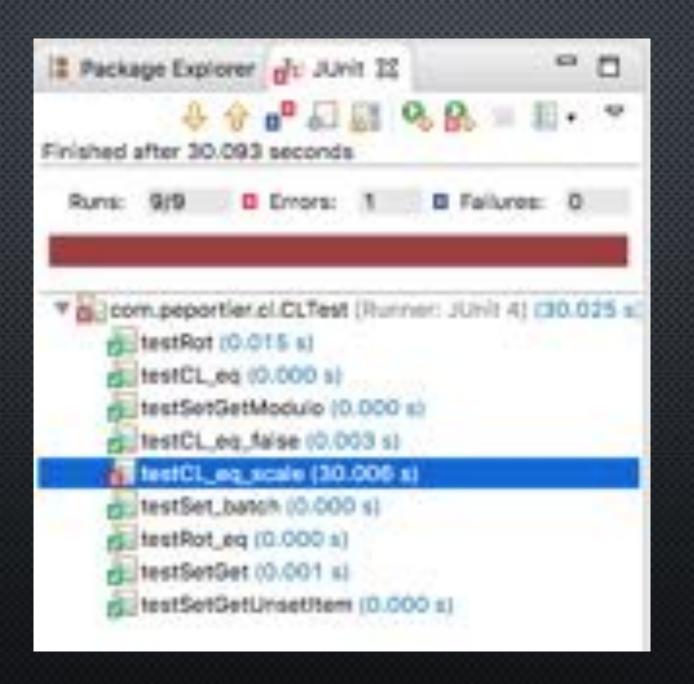
```
/**
 * @param _a a CL
 * @param _b a CL
 * @return Are the lexicographically sorted versions of a and b are equal?
@Requires("_a.s == _b.s")
public static boolean eq(CL _a, CL _b) {
   for(int i=0; i<_a.s; i++){
      if(Arrays.equals(_a.a, _b.rot(i))) return true;
   return false;
```

EXAMPLE OF A COFOJA ERROR



The naı̈ve solution doesn't scale: $O(N^2)$

```
@Test(timeout=30000)
public void testCL_eq_scale() {
   int SIZE = 100000;
   int[] a_vals = new int[SIZE];
   for (int i = 0; i < SIZE; i++) {
       a_vals[i] = i;
   }
   CL a = new CL(SIZE, a_vals);
   int[] b_vals = a.rot(SIZE/2);
   CL b = new CL(SIZE, b_vals);
   assertTrue(CL.eq(a, b));
}</pre>
```



RA and RB are either disjoint or equal

$$(\forall k, i, j :: RA.i \equiv RB.j \Rightarrow RA.(i + k) = RB.(j + k))$$

It is sufficient to compare canonical elements

AA: first element of RA sorted lexicographically

R can be solved by computing AA and BB

To find the solution $res \equiv true$, we can: (a) Find a pair (i,j) such that RA.i = RB.j

To find the solution $res \equiv false$, we can: (b) Observe $AA \neq BB$ To find the solution $res \equiv true$, we can: (a) Find a pair (i, j) such that RA.i = RB.j

Weakening (a) into a loop invariant:

$$P: 0 \le h \land (\forall k: 0 \le k < h: RA.i.k = RB.j.k)$$

$$P \land h \ge N \Rightarrow true \equiv (\exists i, j :: RA.i = RB.j)$$

```
@Test
public void testCL_inv_P() {
    CL a = new CL(3);
    a.set(0,7); a.set(1,9); a.set(2,8);
    CL b = new CL(3);
    b.set(0,9); b.set(1,8); b.set(2,7);
    assertTrue(CL.inv_P(2, 0, 2, a, b));
}
```

```
@Requires({"_h >= 0", "_i >= 0", "_j >= 0"})
public static boolean inv_P(int _h, int _i, int _j, CL _a, CL _b) {
    boolean res = true;
    for (int k = 0 ; k < _h ; k++)
        res &= _a.rot(_i)[k] == _b.rot(_j)[k];
    return res;
}</pre>
```

To find the solution $res \equiv false$, we can: (b) Observe $AA \neq BB$

Weakening (b) into a loop invariant:

$$QA: 0 \le i \quad \land \quad (\forall k: 0 \le k < i: RA.k > BB)$$

$$QA \wedge i \geq N \Rightarrow false \equiv (\exists i, j :: RA.i = RB.j)$$

Symmetrically for *QB* The two CL differ if:

$$QA \wedge QB \wedge (i \geq N \vee j \geq N)$$

```
@Test
public void testCL_inv_Q() {
    CL a = new CL(3);
    a.set(0,7); a.set(1,9); a.set(2,8);
    CL b = new CL(3);
    b.set(0,9); b.set(1,8); b.set(2,8);
    assertTrue(CL.inv_Q(2, b, a));
}
```

```
@Test
public void testCL_smallest() {
    CL a = new CL(3);
    a.set(0,9); a.set(1,8); a.set(2,8);
    int[] expected = {8, 8, 9};
    assertArrayEquals(expected, a.smallest());
}
```

```
public int[] smallest() {
   int aa[] = a.clone();
   for (int k = 1 ; k < s ; k++) {
       int rot_k[] = rot(k);
       boolean new_smallest = false;
       for (int i = 0; i < s; i++) {
               (aa[i] == rot_k[i]) {
              // do nothing
          } else if (aa[i] > rot_k[i]) {
              new_smallest = true; break;
          } else if (aa[i] < rot_k[i]) {</pre>
              new_smallest = false; break;
          } else assert false;
       if (new_smallest) aa = rot_k.clone();
   return aa;
```

```
@Requires(\{"_i >= 0", "_a.s == _b.s"\})
public static boolean inv_Q(int _i, CL _a, CL _b) {
   boolean res = true;
   int bb[] = _b.smallest();
   for (int k = 0; k < _i; k++) {
       int rot_k_of_a[] = _a.rot(k);
       for (int j = 0; j < _a.s; j++) {
           if (rot_k_of_a[j] == bb[j]) {
              if (j == (_a.s - 1)) { res = false; break; }
           } else {
              if (rot_k_of_a[j] > bb[j]) {res &= true; break; }
              else { res = false; break; }
       if (res == false) break;
   return res;
```

We have the sketch of a solution...

```
@Requires("_a.s == _b.s")
@Ensures("result == CL.eq(_a,_b)")
public static boolean eq2(CL _a, CL _b) {
   boolean res = false;
   int h,i,j;
   h = i = j = 0;
   assert CL.inv_P(h, i, j, _a, _b);
   assert CL.inv_Q(i, _a, _b);
   assert CL.inv_Q(j, _b, _a);
   while (h<_a.s && i<_a.s && j<_a.s) {
       assert (h+i+j) <= (3*_a.s - 3);
       // increase h+i+j while maintaining P & QA & QB
   assert CL.inv_P(h, i, j, _a, _b);
   assert CL.inv_Q(i, _a, _b);
   assert CL.inv_Q(j, _b, _a);
   assert (h>=_a.s) || (i>=_a.s) || (j>=_a.s);
   if (h >= _a.s) res = true;
   else if (i>=_a.s || j>=_a.s) res = false;
   return res;
```

$$P: 0 \le h \land (\forall k: 0 \le k < h: RA.i.k = RB.j.k)$$

When RA.i.h = RB.j.h (i.e., A.(i + h) = B.(j + h)):

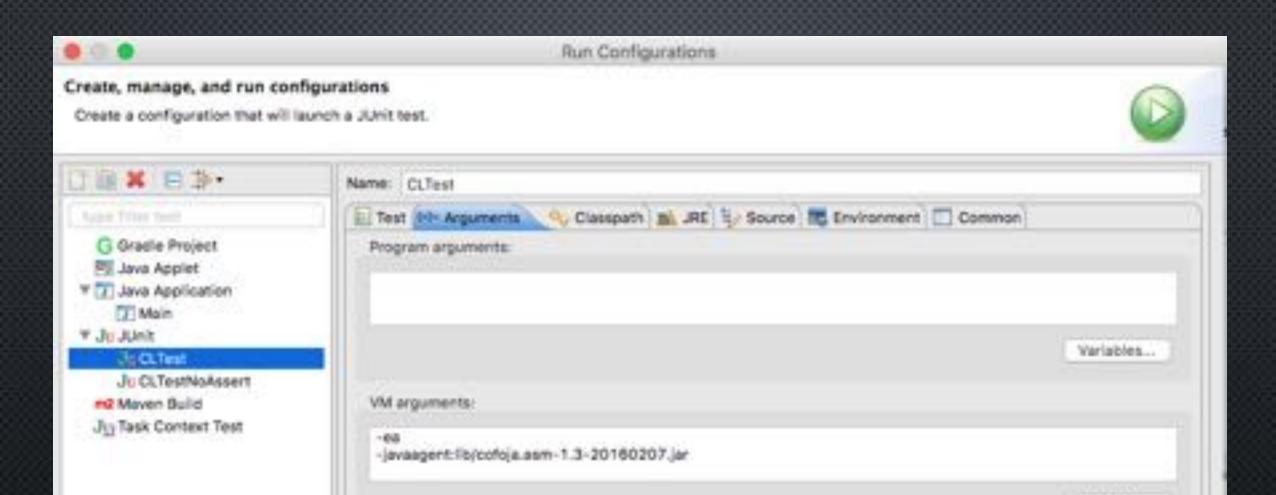
 $h \leftarrow h + 1$ will maintain P and assure progress.

```
P: 0 \leq h
                     (\forall k: 0 \leq k < h: RA.i.k = RB.j.k)
                 \land \quad (\forall \, k : 0 \le k < i : RA.k > BB)
QA: 0 \leq i
        RA.i.h > RB.j.h \land P
     \Rightarrow {def. P, 0 \le p \le h}
        RA.i.h > RB.j.h \land (\forall k : p \le k < h : RA.i.k = RB.j.k)
     ⇒ {lexicographic ordering}
        RA.(i + p) > RB.(j + p)
     \Rightarrow {def. BB}
        RA.(i+p) > BB
        QA \land P \land A.(i+h) > B.(j+h)
     ⇒ {see above}
        QA \land (\forall p: 0 \le p \le h: RA.(i+p) > BB)
     = {renaming the bounded variable: i + p = k}
        QA \land (\forall k : i \leq k \leq i + h : RA.k > BB)
     = \{ \text{def. } QA \}
        QA[i \setminus i + h + 1]
```

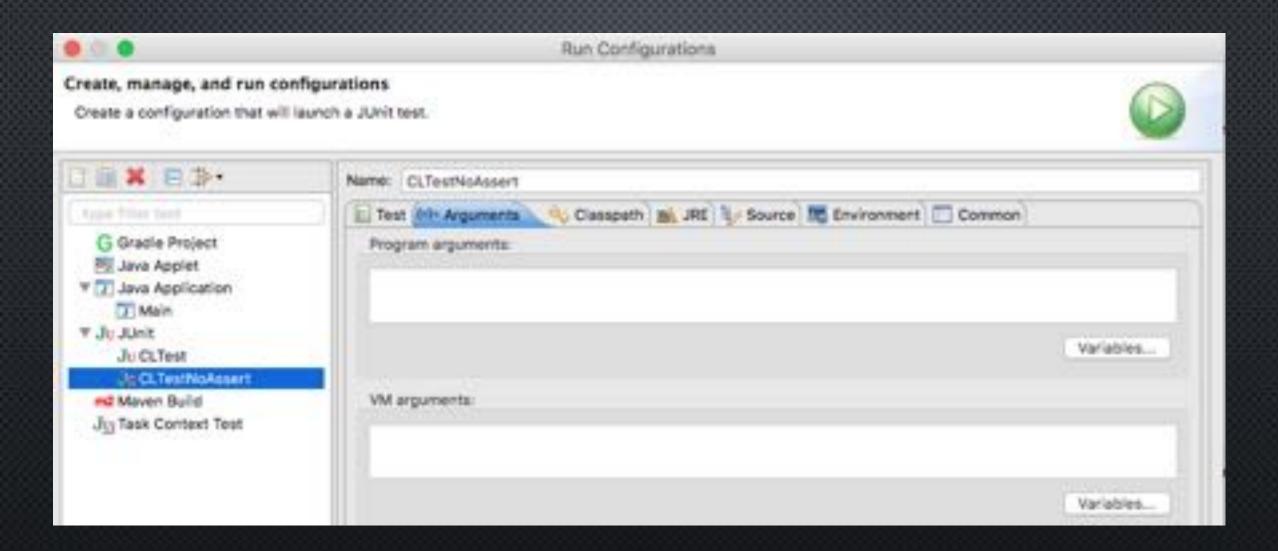
```
int progress;
while (h<_a.s && i<_a.s && j<_a.s) {
   assert (h+i+j) <= (3*_a.s - 3);
   progress = h+i+j;
   // increase h+i+j while maintaining P & QA & QB
   if (_a.get(i+h) == _b.get(j+h)) h = h+1;
   else if (_a.get(i+h) > _b.get(j+h)) {i = i+h+1; h=0;}
   else if (_b.get(j+h) > _a.get(i+h)) {j = j+h+1; h=0;}
   else assert false;
   assert CL.inv_P(h, i, j, _a, _b);
   assert CL.inv_Q(i, _a, _b);
   assert CL.inv_Q(j, _b, _a);
   assert h+i+j > progress;
```

```
@Test(timeout=60000)
public void testCL_eq2_scale() {
   int SIZE = 100000;
   int[] a_vals = new int[SIZE];
   for (int i = 0; i < SIZE; i++) {
       a_vals[i] = i;
    CL a = new CL(SIZE, a_vals);
   int[] b_vals = a.rot(SIZE/2);
   CL b = new CL(SIZE, b_vals);
    assertTrue(CL.eq2(a, b));
```





Variables...



http://www.eclemma.org/

```
1389
          @Requires("_a.s -- _b.s")
 139
          #Ensures("result == (L.eq(_0,_b)")
          public static boolean eq2(CL _a, CL _b) {
 1.48
             boolean res = false;
 141
              int h,i,j;
 142
143
             h - 1 - 1 - 0:
              assert (L.inv_P(h, i, j, _a, _b);

⊕144.

9145
             assert CL.inv_QCi, _a, _b);
146
              dissert (L.inv_Q(j, _b, _a);

    147

              while (ha_a.s && ia_a.s && ja_a.s) {
9148
                 dssert (h+i+j) ex (3*_a.s - 3);
                 // increase hais; while mointaining P & QA & QB
1.49
158
                 if (_a.get(i+h) == _b.get(j+h)) h = h+1;
                 else if (_a.get(i+h) > _b.get(j+h)) {i = i+h+1; h-0;}
9151
                  else if (_b.get(j+h) > _a.get(i+h)) [j = j+h+1; h=8;}
9152
                 alse assert false o 1 of 2 branches missed.
4153
154
              assert CL. inv_P(h, 1)
9155
$156
              assert CL. (mv_Q(i, _a, _b))
              assert (L.inv_Q(j, _b, _a);
157
9158
              assert (h-_a.s) || (i--a.s) || (j--a.s);
159
              if (h >= _a.s) res = true;
              else if (in-a.s II ja-a.s) res - false;
9160
 161
              return res;
 162
```

```
8.60
 118-
          #Requires({"_i >- 0", "_o.s -- _b.s"})
          public static boolean inv_Q(int _i, CL _a, CL _b) {
 110
 120
              boolean res - true;
              int bb[] = _b.get_copy_err();
 121
              Arroys.sort(bb);
 122
              for (int k = 8 ; k < _l ; k++) [
123
                  int rot_k_of_o[] - _a.rot(k);
124
                  for (int ) = 0 ; ] = _0.5 ; ]++) {
@125
                      if (rot_k_of_a[j] \rightarrow bb[j]) {
@126
                          if (i - (.e.s - 1)) i res - false; break; }
$122
                       } else {
128
Q129
                           if (rot_k_of_o[j] > bb[j]) {res &= true; break; }
                           else { res - false; breck; }
 130
 131
 132
9133
                   if (res -- false) break;
 134
 135
              return resi
 136
```

Shiloach, Yossi.

"A fast equivalence-checking algorithm for circular lists." Information Processing Letters 8.5 (1979): 236-238.

Gasteren, Antonetta JM.
"On the shape of mathematical arguments."
Vol. 445. Springer Science & Business Media, 1990.

https://en.wikipedia.org/wiki/Lexicographically_minimal string_rotation

LINKS

- HTTP://DOCS.ORACLE.COM/JAVASE/6/DOCS/TECHNOTES/GUIDES/LANGUAGE/ASSERT.HTML #USAGE-CONDITIONS
- HTTPS://GITHUB.COM/NHATMINHLE/COFOJA
- HTTPS://WWW.EIFFEL.COM/VALUES/DESIGN-BY-CONTRACT/
- HTTPS://GITHUB.COM/JUNIT-TEAM/JUNIT4/WIKI/ASSERTIONS
- http://www.vogella.com/tutorials/JUnit/article.html#using-junit-integrated-
 http://www.vogella.com/tutorials/JUnit/article.html#using-junit-integrated-
 http://www.vogella.com/tutorials/JUnit/article.html#using-junit-integrated-