

TD1 - Processes and signals

Exercise 1 - Who am I?

```
int main ()
{
    pid_t child = fork();
    if (child == 1) {
        perror("fork() error");
        exit(1);
    }

    // à ajouter pour la question (b)
    // printf("My PID is %d.\n", getpid());

    if (child == 0) {
        printf("Child process: my PID is %d.\n", getpid());

        // à enlever pour la question (b)
        exit(0);
    }

    printf("Now my PID is %d.\n", getpid());
    exit(0);
}
```

a. What is the output of this program ?

This program outputs the child and parent PID, like so:

```
Child process: my PID is 4321.
Now my PID is 1234.
```

The order of the prints being arbitrary.

b. What is the output, if we add line 10 and remove line 16 ?

This program outputs the child PID 3 times and the parent PID twice, because the child process doesn't terminate until the last line, like so:

```
My PID is 4321.
My PID is 1234.
Child process: my PID is 4321.
Now my PID is 4321.
Now my PID is 1234.
```

The order of the prints being arbitrary.

Exercise 2 - Bob

```
int sunny_days_in_London(int year) { /* 6 heures de calcul */ }
int second_year_presence(int day) { /* 6 heures de calcul */ }

int main ()
{
    int captains_age = 0;
    pid_t cpid = fork();

    if (cpid != 0)
        captains_age += sunny_days_in_London(1305);
    else
        captains_age += second_year_presence(1305);

    printf(
        "Le capitaine a %d ans et mon PID est %d\n",
        captains_age,
        (int) getpid()
    );

    return 0;
}
```

Explain Bob's mistake and propose a way to solve the problem using `exit()` and `wait()`. Detail the actions of both processes.

Bob didn't make the child process exit after doing the calculation, so there are going to be 2 different prints.

Also, Bob forgot to wait for the child process to end before printing out the result, and to check for an exit code to check for errors.

This program is going to output two lines, containing the 2 different calculations.

To solve the problem, we could add `exit(age);` in the `if`, and a `wait()` before the `printf` statement.

Corrected program

```
int sunny_days_in_London(int year) { /* 6 heures de calcul */ }
int second_year_presence(int day) { /* 6 heures de calcul */ }

int main ()
{
```

```
int captains_age = 0;
int status;

pid_t cpid = fork();

if (cpid != 0) {
    captains_age += sunny_days_in_London(1305);
    status = 0;
    exit(captains_age); // or return age;
} else
    captains_age += second_year_presence(1305);

wait(&status);
WIFEXITED(&status) {
    captains_age += WEXITSTATUS(&status);
}

printf(
    "Le capitaine a %d ans et mon PID est %d\n",
    captains_age,
    (int) getpid()
);

return 0;
}
```

Exercise 3 - Firefork()

```
int main(int argc, char ** argv)
{
    pid_t id;
    int i, N = 0;

    if (argc > 1) N = atoi(argv[1]);

    for (i = 0; i < N; i++) {
        id = fork();

        if (id == 1) {
            perror("fork() error"); exit(1);
        }

        printf(
            "I am %d, son of %d.\n",
            getpid(),
            getppid()
        );
    }

    printf("%d out.\n", getpid());
}
```

```
    return 0;
}
```

Give a generic formula for the number of processes depending on N.

This program doesn't check if you're the parent or a child process : so this is going to fork from the parent and the childs. Now let's make an example :

Let's say we execute `./fforks 3` :

`atoi('3') = 3` so the loop is going to go 3 times.

1st iteration :
fork 1 child

2nd iteration :
Both the parent and the child fork a process.
We now have 4 total processes.

3rd iteration :
Each 4 of the processes fork a child.
We now have a total of 8 processes.

You've probably noticed the pattern : this program makes $O(N)$ child processes.

This type of program is called a **fork bomb**.

Exercise 4 - world! Hello,

```
int main ()
{
    pid_t id = fork();

    if (id == 0) {
        printf("Hello, ");
        exit(0);
    }

    // <-- TROU

    printf("world!");
    return 0;
}
```

We want to complete this program to make sure that the words are always printed in the right order, independently of scheduling decisions.

The system call `int sleep(int n)` puts the program in the waiting state for `n` seconds. Can we complete the program by calling `sleep(1)`? `sleep(100)`? Explain your answer.

Completing the program by adding a `sleep` system call might work, but is totally unreliable, no matter the amount of seconds slept.

Can we complete the program with a `wait()` call? Substantiate your answer.

Yes, a `wait` system call is the right option here : it allows the parent process to wait for a `SIGCHILD` signal, which is sent when the child process has finished execution.

It is reliable and this way, we can reliably execute our program in a organized and predictable way.

Exercise 5 - To infinity in eight seconds.

The following program increments a counter of type `unsigned int` and stops after the maximal value is reached and the counter goes back to zero.

```
volatile unsigned int count = 0;
volatile int step = 1;
//
// < TROU 1
//

int main() {
    //
    // < TROU 2
    //
    for (count = 1; count > 0; count += step);
    return 0;
}
```

We declare the variables `count` and `step` as `volatile` to prevent compiler optimisation (which otherwise might remove the loop entirely).

Complete the code so that `count` is reset to 1 when the process receives the `SIGUSR1` signal.

```
volatile unsigned int count = 0;
volatile int step = 1;

void hnd1(int sig) { // signal handler
    count = 1;
}

int main() {
    // bind SIGUSR1 to handler hnd1
    signal(SIGUSR1, &hnd1);

    for (count = 1; count > 0; count += step);
    return 0;
}
```

Complete the code so that the increment step changes sign when the process receives SIGUSR2.

```
volatile unsigned int count = 0;
volatile int step = 1;

void hnd1(int sig) { // signal handler
    count = 1;
}

void hnd2(int sig) {
    step = -step;
}

int main() {
    // bind SIGUSR1 to handler hnd1
    signal(SIGUSR1, &hnd1);
    signal(SIGUSR2, &hnd2);

    for (count = 1; count > 0; count += step);
    return 0;
}
```

To ensure that our program behaves correctly, we want to watch the counter. A new system call comes in handy:

```
#include <unistd.h>
unsigned int alarm(unsigned int nb_sec);
```

The system call `alarm()` sets a timer that will send a `SIGALRM` signal to the calling process `nb_sec` seconds later. Any `alarm()` call cancels and replaces the previously set alarm; `alarm(0)` cancels the current alarm and does not set a new one.

Complete the code so that the current value of count is printed every second.

```
#include <unistd.h>

volatile unsigned int count = 0;
volatile int step = 1;

void hnd1(int sig) { // signal handler
    count = 1;
}

void hnd2(int sig) {
    step = -step;
}
```

```

void hnd3(int sig) {
    printf("Count : %d", count);
    alarm(1);
}

int main() {
    // bind SIGUSR1 to handler hnd1
    signal(SIGUSR1, &hnd1);
    signal(SIGUSR2, &hnd2);
    signal(SIGALRM, &hnd3);

    alarm(1);

    for (count = 1; count > 0; count += step);
    return 0;
}

```

Exercise 6 - Et mon courroux, coucou!

Let us go back to [Exercise 4](#). As we know, the parent process receives the `SIGCHLD` signal whenever its child process terminates. Let us see if we can use this to ensure the correct order of execution.

The primitive `int pause(void)` puts the calling process to sleep until it receives a signal which is neither ignored nor blocked. Can we complete the program with a `pause()` call? Explain your answer.

No, because with the program as is, the `SIGCHLD` signal will get ingored by the `pause()` system call.

Can we complete the program with a `pause()` call and a simple handler for `SIGCHLD`, for example, a function that does nothing?

Yes, but we need to make sure that our handler is defined before the `fork()`.

```

void hnd(int sig) {}

int main()
{
    signal(SIGCHLD, hnd);
    // rest of program
}

```

Bob is not happy with our answers. What if we used our signal handler to avoid calling `pause()` if the signal has already arrived? To do that, we would need a global `flag` variable:

```

volatile int flag = 0;
void handler_chld (int sig) { flag = 1; }

```

```
int main ()
{
    signal(SIGCHLD, handler_chld);
    if (fork() == 0) { printf("Hello"); exit(0); }
    while (!flag) pause();
    printf(" world");
    return 0;
}
```

Why is it important to install the handler before calling `fork()`?

Because otherwise the handler might not be defined in the parent process when the child process exits.

Why is it important to test the flag in a `while` loop?

It's important to do so because if the parent evaluates `flag` before the child process has ended it will get past the instruction and exit.

Does Bob's program guarantee that the words are printed in the right order?

Yes, `while(!flag)` waits for the end of the child process before resuming execution. It is however very demanding for the machine : see `busy waiting`.

Exercise 7 - One, two, many.

Let us execute the following program :

```
volatile int count = 0;
void handler(int sig) { count++; }

int main()
{
    int i;
    signal(SIGUSR1, handler);

    if (fork() == 0) {
        for (i = 0; i < 256; i++)
            kill(getppid(), SIGUSR1);
        exit(0);
    }
    wait(NULL);
    printf("Final: %d\n", count);
    return 0;
}
```

What is the output of this program? Explain the answer.

`man 7 signal` gives us the following definition for `SIGUSR1` :

Signal	Norm	Action	Note
SIGUSR1	P1990	Term	User signal 1

The `action` attribute is `term`, meaning XYZ...